Interaction Graphs Derived from Activation Functions and their Application to Gene Regulation

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INTERACTION GRAPHS DERIVED FROM ACTIVATION FUNCTIONS
AND THEIR APPLICATION TO GENE REGULATION

BY

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BS, National University of Ireland Maynooth, 2006
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DISSertation

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Abstract

Interaction graphs are graphic representations of complex networks of mutually interacting components. Their main application is in the field of gene regulatory networks, where they are used to visualize how the expression levels of genes activate or inhibit the expression levels of other genes.

First we develop a natural transformation of activation functions and their derived interaction graphs, called conjugation, that is related to a natural transformation of signed digraphs called switching isomorphism. This is a useful tool for the analysis of interaction graphs used throughout the rest of the dissertation.

We then discuss the question of what restrictions, if any, apply to interaction graphs derived from activation functions. Within these restrictions, we then construct activation functions with any desired interaction graph. The specific case of threshold activation functions, a commonly used kind of activation function, is also considered.

We then conclude with some discussion, and new proofs of the conjectures of René Thomas, using the theory of conjugate activation functions. These conjectures relate feedback in the interaction graph to dynamic properties of multi-stationarity and periodic stability. We prove a more general form of Richard and Comet’s version of René Thomas’ first conjecture. Included is a new counterexample to a local version of Thomas’ second conjecture, on only eight components. This is the smallest counterexample that I am currently aware of.
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Introduction

Signed digraphs are directed graphs where each arc is signed, + or -. They are often used to represent complex systems of multiple interacting components, such as gene regulatory networks. Positive or negative arcs between nodes denote positive or negative influence of one component on another. This dissertation is primarily concerned with interaction graphs, a kind of signed directed graph that is derived from an activation function. Interaction graphs, together with their activation function intend to model these kinds of complex systems of multiple interacting components.

One of the main purposes of this dissertation is to determine what, if any, restrictions there are on the structure of interaction graphs. More specifically, for a given signed digraph $\Delta$, when can we find an activation function whose interaction graph is $\Delta$? What can we say about $\Delta$ if no such activation function exists?

We also discuss the conjectures of René Thomas. These conjectures, loosely stated, say that positive and negative feedback in the interaction graph are necessary for multi-stationary and stable-periodic dynamics respectively. These phenomenon are of great interest in biology since multi-stationary and stable-periodic dynamics of gene expression levels correspond to processes of cell differentiation and homeostasis respectively. These kinds of dynamics are also of interest in the theory of chemical reaction networks and population models. The strongest versions of René Thomas’ conjectures that I am aware of have been proved in [4] and [5]. Using my new techniques developed herein, we will prove the same, or slightly more general versions of these conjectures. We also present a smaller counter example to a local version of René Thomas’ second conjecture than the first known counter example presented in [6].

The dissertation is arranged as follows.

Chapter 1 is a short development of relevant background theory on digraphs and signed digraphs that will be used throughout this dissertation.

In chapter 2, we begin with the development of my take on the theory of activation functions. We develop some theory on their associated interaction graphs and present new methods of transforming activation functions and their interaction graphs which aid greatly in their analysis. Then
we tackle what, if any, restrictions there are on the structure of interaction graphs as well as how to construct activation functions that have a given interaction graph. We do this first for the general case, and then for the specific case of threshold activation functions.

In chapter 3 we introduce the state transition graph, a graph that represents how the states of an activation function can change over time. Then, using new techniques developed in this dissertation, we explore René Thomas’ first and second conjectures, on multi-stationary and stable-periodic dynamics in the state transition graph respectively. Finally, we present a new counter-example to a local version of René Thomas’ second conjecture. This counter-example has fewer components than previously known counterexamples.
Chapter 1

Signed Digraphs

1.1 Digraphs

A directed graph, or digraph is a pair $D = (V, A)$ where $V$ the finite, non-empty vertex set of $D$ and $A$ is the finite arc set of $D$, a multiset whose elements are from $V \times V$. For an arc $a = (v, w)$, we say that $a$ is incident with $v$ and $w$. For two subsets of vertices, $A$ and $B$, we say that there is an arc between $A$ and $B$ if there is an arc incident with a vertex in $A$ and a vertex in $B$.

A digraph $D$ is connected if $D$ contains a single vertex, or for all non-empty $W \subset V$, there is an arc between $W$ and $V \setminus W$. For us, $\subset$ means strict-subset, i.e., $W \neq V$.

The in-degree (out-degree) of $v$ in $D$ is $|\{(w \mid (w, v) \in D)\}|$ ($|\{(w \mid (v, w) \in D)\}|$). The degree of $v$ is the sum of the in-degree and out-degree of $v$.

The degree of the vertices and the number of arcs in a digraph are related in a nice way.

**Proposition 1.1.1.** Let $d_v$ be the degree of the vertex $v \in D$. Then $\sum_{v \in D} d_v = 2|A|$.

**Proof.** Any arc $(a, b)$ will contribute 1 to the sum in $d_a$ and 1 to the sum in $d_b$. So this arc is counted twice in the sum $\sum_{v \in D} d_v$. Therefore adding the degree of every vertex counts each arc twice.  

Let the bijection $\phi : V \rightarrow \phi V$. The digraph $\phi D = (\phi V, \phi A)$ where $\phi A = \{(\phi v, \phi w) \mid (v, w) \in D\}$. Two digraphs $D$ and $D'$ are called isomorphic if there exists a bijection $\phi$ such that $\phi D = D'$.

### 1.1.1 Subgraphs, Paths, Trees, Circles and Cycles

For the remainder of this section, let the digraphs $D = (V, A)$, $D_1 = (V_1, A_1)$ and $D_2 = (V_2, A_2)$. $D_1$ is a sub-digraph of $D_2$ if $V_1 \subseteq V_2$ and $A_1 \subseteq A_2$. Often we will call sub-digraphs just subgraphs.

The union of two digraphs is just the union of their vertex and arc sets respectively, or $D_1 \cup D_2 = (V_1 \cup V_2, A_1 \cup A_2)$.

If $V_1 \subseteq V_2$ and $A_1 = \{(i, j) \in A_2 \mid i, j \in V_1\}$, then we call $D_1$ the subgraph of $D_2$ induced by $V_1$. Subgraphs of this form are called induced subgraphs.
If \( D \) is connected and \(|A| = |V| - 1\), then \( D \) is called a tree. Combining this definition with proposition [1.1.1]

\[
\sum_{v \in D} d_v = 2|A| = 2|V| - 2.
\] (1.1)

I have much use for trees, so their theory demands development.

**Proposition 1.1.2.** A tree \( T \) contains a vertex of degree zero if and only if the tree consists of a single vertex.

**Proof.** Suppose that \( T \) contains a vertex \( v \) of degree zero. If \(|V| > 1\), then there is an arc to or from \( v \) to \( V \setminus v \) since \( T \) is connected. So the degree of \( v \) is at least 1, a contradiction. Therefore \(|V| = 1\). Using equation (1.1), \( \sum_{v \in V} d_v = 2|V| - 2 = 0 = 2|A| \). Therefore \( |A| = 0 \), so \( T \) contains no arcs and so consists of a single vertex. Conversely, if \( T \) consists of a single vertex \( v \), then clearly the degree of \( v \) is zero. \( \square \)

**Proposition 1.1.3.** A tree \( T \) that contains more than one vertex has at least two vertices of degree-1.

**Proof.** By proposition [1.1.2], \( T \) has no degree-0 vertex, since \( T \) contains more than one vertex. Suppose all but one vertex in \( T \) has degree at least two. Then \( \sum_{v \in T} d_v \geq 2|V| - 1 \). Combining this with equation (1.1), we see that \( 2|V| - 2 \geq 2|V| - 1 \), a contradiction. Therefore \( T \) contains at least two degree-1 vertices. \( \square \)

If the tree \( T \) is a subgraph of the digraph \( D \) and \( T \) contains every vertex of \( D \), then \( T \) is called a spanning-tree of \( D \).

**Proposition 1.1.4.** Every connected digraph \( D = (V, A) \) contains a spanning-tree.

**Proof.** The proof is an iterative construction.

Any subgraph of \( D \) consisting of a single vertex is a tree, so any digraph \( D \) contains tree-subgraphs.

So let \( T_k = (V_k, A_k) \) be a tree-subgraph of \( D \). If \( T_k \) is a spanning-tree, we are done. If not, then there is an arc \( a_k \) between \( V_k \) and \( V \setminus V_k \) since \( D \) is connected. Let \( v_k \) be the vertex in \( V \setminus V_k \) incident with \( a_k \). Let \( V_{k+1} = V_k \cup v_k \) and \( A_{k+1} = A_k \cup a_k \). I claim the subgraph \( T_{k+1} = (V_{k+1}, A_{k+1}) \) is also a tree. To justify this, we have to show \( T_{k+1} \) is connected and has one fewer arcs than vertices.

Let us first show that \( T_k \) one fewer arcs than vertices. Since \( T_k \) is a tree, \(|A_k| = |V_k| - 1\). Therefore \(|A_k \cup a_k| = |V_k \cup v_k| - 1\), i.e., \(|A_{k+1}| = |V_{k+1}| - 1\).
To see $T_{k+1}$ is connected, let $W \subset V_{k+1}$ be non-empty. We may assume that $W$ contains $v_k$ since the following argument can be just as easily applied to the compliment of $W$. If $W$ contains only $v_k$, then $a_k$ is an arc between $W$ and $V_{k+1} \setminus v_k = V_k$. If $W$ contains vertices other than $v_k$, then there is an arc between $W \setminus v_k$ and $V_k \setminus W$ in $T_k$ since $T_k$ is connected. Therefore the same arc is between $W$ and $V_{k+1} \setminus W$ in $T_{k+1}$. Therefore $T_{k+1}$ is connected.

Now if $T_{k+1}$ is a spanning tree, we are done. If not, then we can iterate this process $|V| - |V_k|$ times to get a tree contained in $D$ whose vertex set is $V$. This tree is then by definition a spanning tree of $D$.

A tree that contains no vertex of degree 3 or more is called a path.

**Proposition 1.1.5.** A path is either a single vertex, or it contains exactly two vertices of degree-1 and all other vertices have degree 2.

**Proof.** Let $P = (V, A)$ be the path. If $|V| = 1$, then it $P$ consists of a single vertex by proposition 1.1.2.

What if $P$ contains more than one vertex? By proposition 1.1.3, $P$ contains at least two degree-1 vertices. Since the degree of all other vertices is less than three, $\sum_{v \in V} d_v = |V| - k$, where $k$ is the number of degree-1 vertices. Combining this with equation (1.1), $|V| - 2 = |V| - k$, i.e., $k = 2$. So $P$ contains exactly two degree-1 vertices. □

If a digraph $D$ contains a path such that the degree-1 vertices of the path are $v$ and $w$, then we say there is a path between $v$ and $w$ in $D$.

**Proposition 1.1.6.** In a tree, there is a path between every pair of vertices.

**Proof.** The proof is by induction on the number of vertices.

Clearly the proposition is true if the tree consists of a single vertex. If the tree contains exactly two vertices, then each vertex has degree 1 by proposition 1.1.3. So the tree contains a single arc between the two vertices, so it is a path.

Now suppose $T = (V, A)$ is a tree on $k > 2$ vertices and that the proposition is true for trees with less than $k$ vertices. We need the following lemma to complete the induction step.

**Lemma 1.1.7.** If $T = (V, A)$ is a tree and $v$ is a degree-1 vertex in $T$ incident with the arc $a$, then the subgraph $S = (V \setminus v, A \setminus a)$ is a tree.
Proof. To prove this, we have to show that $S$ is connected and that $|A \setminus a| = |V \setminus a| - 1$.

Since $T$ is a tree, $|A| = |V| - 1$. Therefore $|A \setminus a| = |V \setminus a| - 1$.

Is $S$ connected? Since $T$ has a degree-1 vertex, $V$ contains at least two vertices since it is not just a single vertex by proposition 1.1.2.

If $|V| = 2$, then $T$ contains a single arc and $S$ consists of a single vertex. So the result is true in this case.

If $|V| > 2$, let $W \subset V \setminus v$ be non-empty. This is possible since $V \setminus v$ contains at least two vertices. We just have to show that there is an arc between $W$ and $(V \setminus v) \setminus W$. Since $T$ is connected, there is an arc $a$ between $W$ and $V \setminus W$. Similarly there is an arc $b$ between $W \cup v$ and $(V \setminus v) \setminus W$.

If $a$ and $b$ are both incident with $v$, then $a$ and $b$ are distinct arcs since the other end of $a$ is in $W$ and the other end of $b$ is in $(V \setminus v) \setminus W$. But then $v$ would have degree 2. Since the degree of $v$ is one, at most one of $a$ or $b$ is incident with $v$. So whichever arc is not incident with $v$ is an arc between $W$ and $(V \setminus v) \setminus W$. Therefore, $S$ is connected, so $S$ satisfies both tree conditions.

Let us continue the proof of proposition 1.1.6. By proposition 1.1.3, $T$ contains a degree-1 vertex $v$ incident with an arc $a$. Let $S = (V \setminus v, A \setminus a)$. By lemma 1.1.7, $S$ is a tree on less than $k$ vertices. Therefore it satisfies the proposition by the induction hypothesis. So there is a path between every pair of vertices in $S$. This mean that in $T$, there is a path between every pair of vertices excluding $v$. We just need to show that in addition, there is a path between $v$ and any other vertex $w \in T$. Let $x$ be the other vertex incident with $a$. Since $x, w \in S$, there is a path between $x$ and $w$ in $S$. This path is also contained in $T$. Combining this path with arc between $v$ and $x$ in $T$ yields a path between $w$ and $v$ in $T$.

A directed path, or dipath is a path that consists of a single vertex, or a path that contains a single vertex of out-degree 1, a single vertex of in-degree 1 and all other vertices have in-degree and out-degree 1. In other words, a dipath is a path such that the direction of every arc in the path is consistent. If a digraph $D$ contains a dipath such that the vertex with in-degree 1 is $w$ and the vertex with out-degree 1 is $v$, then we say there is a dipath in $D$ from $v$ to $w$.

A digraph $D$ is strongly connected if for any pair of vertices $v$ and $w$ in $D$, there is dipath from $v$ to $w$ and a dipath from $w$ to $v$. A subgraph of $D$ is called a strong component of $D$ if it is a maximal strongly connected subgraph of $D$. A strong component $C$ is terminal there are no arcs $(v, w)$ in $D$ such that $v \in C$ and $w \notin C$. A strong component $C$ is initial there are no arcs $(v, w)$ in $D$ such that $v \notin C$ and $w \in C$. 

6
Proposition 1.1.8. The strong components of a digraph partition the vertex set of the digraph.

Proof. Consider the relation ∼ on the vertex set where \( v \sim w \) if there is a dipath from \( v \) to \( w \) and a dipath from \( w \) to \( v \). If ∼ is an equivalence relation, then the equivalence classes of ∼ are strong components of the digraph. So to prove the proposition, we will show that ∼ is an equivalence relation.

There is always a path from \( v \) to itself, so the relation is reflexive. The relation is clearly symmetric. If \( v \sim w \) and \( w \sim x \) then there is a dipath from \( v \) to \( w \) and from \( w \) to \( x \). The union of two dipaths is not necessarily itself a dipath, but we can use the following lemma.

Lemma 1.1.9. If \( P \) is a dipath from \( v \) to \( w \) and \( Q \) is a dipath from \( w \) to \( x \), then \( P \cup Q \) contains a dipath from \( v \) to \( x \).

Proof. Starting at \( v \), we follow the dipath \( P \) until we encounter the first vertex \( y \) contained in \( Q \). This is guaranteed to happen since \( P \) contains at least one vertex in \( Q \), namely \( x \). Let \( P' \subseteq P \) be the part of \( P \) that goes from \( v \) to \( y \). Let \( Q' \subseteq Q \) be the part of \( Q \) that goes from \( y \) to \( x \). Both \( P' \) and \( Q' \) are dipaths that only intersect at \( y \) since \( y \) is the earliest vertex in \( P \) that is also in \( Q \). Therefore \( P' \cup Q' \) is a dipath from \( v \) to \( x \). □

This lemma shows that ∼ is transitive. Therefore ∼ is an equivalence relation. □

A circle in \( D \) is a connected subgraph \( C \subseteq D \) such that the degree of every vertex in \( C \) is 2. If every vertex of \( C \) has in-degree one and out-degree one, then we call \( C \) a cycle.

1.2 Signed Digraphs

A signed digraph is a pair \( \Delta = (V, A) \) where \( V \) is the non-empty vertex set of \( \Delta \) and \( A \subseteq V \times V \times \{\pm\} \) is the arc set of \( \Delta \). The vertex set of \( \Delta \) will also be denoted by \( V(\Delta) \). If \( (v, w, \sigma_{vw}) \in A \), we say there is an arc from \( v \) to \( w \) with sign \( \sigma_{vw} \). An illustration of a signed digraph is given in figure 1.1

Signed digraphs are themselves digraphs if you ignore the arc signs. So every digraph concept applies equally well to signed digraphs.

If there is at most one arc between any two vertices in the same direction, and at most one loop on any vertex, then the signed digraph can be represented with a matrix. If \( \Delta = (V, A) \) is a signed digraph with this property, then the adjacency matrix of \( \Delta \), \( \text{Adj}(\Delta) \) is the \( |V| \times |V| \) matrix whose \( ij \) entry is \( \sigma_{ij} \) if the arc \((i, j, \sigma_{ij})\) is contained in \( \Delta \) and 0 otherwise.
Figure 1.1: A signed digraph with arc set
\{ (1, 2, +), (1, 2, -), (2, 3, +), (3, 2, -), (3, 1, +), (4, 4, +), (4, 4, -) \}.

Figure 1.2: Signed digraph \( \Delta \) that has an adjacency matrix.

For example, the signed digraph \( \Delta \) in figure 1.2 has the following adjacency matrix.

\[
\text{Adj}(\Delta) = \begin{pmatrix}
0 & - & 0 & 0 \\
0 & 0 & + & 0 \\
+ & - & 0 & 0 \\
0 & 0 & 0 & +
\end{pmatrix}
\]

However, the signed digraph in figure 1.1 does not have an adjacency matrix because it has a pair of parallel arcs from vertex 1 to vertex 2.

1.2.1 Balance

The \textit{sign of a signed digraph} is the product of the signs of all of its arcs. The sign of \( \Delta \) is written as \( \text{sgn} \, \Delta = \prod_{(i,j,\sigma_{ij}) \in A} \sigma_{ij} \). So a signed digraph is positive (negative) if and only if it contains an even (odd) number of negative arcs.

A signed digraph \( \Delta \) is \textit{balanced} if every circle contained in \( \Delta \) is positive. A signed digraph is \textit{cycle-balanced} if every cycle it contains is positive. Examples of balanced and cycle-balanced digraphs are given in figure 1.3. The signed digraph in figure 1.1 is not balanced.
1.3 Switching Isomorphism

Given a signed digraph $\Delta$ with vertex set $V$. Let the bijection $\phi : V \to W$. We may write $W = \phi V$.

Let $\zeta \in \{\pm\}^V$. We call the pair $\lambda = (\phi, \zeta)$ a switching isomorphism of $\Delta$. If $\phi$ is a permutation, then we may also call $\lambda$ a signed permutation of $V$. The switching isomorphism $\lambda$ transforms $\Delta$ in the following way. The signed digraph $\lambda \Delta = (W, \{(\phi i, \phi j, \zeta_i \sigma_{ij} \zeta_j) | (i, j, \sigma_{ij}) \in \Delta\})$. If we write $\zeta \Delta$, this means that $\zeta \Delta = \lambda \Delta$ where the bijection $\phi$ in $\lambda$ is the identity map. In this case we can call the map $\zeta$ a switching of $\Delta$. Similarly $\phi \Delta = \lambda \Delta$ where $\zeta = \{+\}^V$. In this case we call $\phi$ an isomorphism of $\Delta$. Let $\Delta_1$ and $\Delta_2$ be signed digraphs. If there is a switching isomorphism $\lambda$ of $\Delta_1$ such that $\lambda \Delta_1 = \Delta_2$, then we say that $\Delta_1$ and $\Delta_2$ are switching isomorphic. An example of a switching isomorphism is given in figure 1.3.

It is an elementary observation that switching isomorphisms preserve the signs of circles.

**Proposition 1.3.1.** Let $\lambda$ be a switching isomorphism of the signed digraph $\Delta$. If $C \subseteq \Delta$ is a circle, then $\text{sgn } C = \text{sgn } \lambda C$.

**Proof.** Note that $\text{sgn } \lambda C = \prod_{(i, j, \sigma_{ij}) \in \Delta} \zeta_i \sigma_{ij} \zeta_j = \prod_{(i, j, \sigma_{ij}) \in C} \zeta_i \sigma_{ij} \zeta_j$. Since the degree of each vertex in $C$ is two, $\zeta_i$ occurs twice in this product for any vertex $i$ of $C$. Therefore $\prod_{(i, j, \sigma_{ij}) \in C} \zeta_i \sigma_{ij} \zeta_j = \prod_{(i, j, \sigma_{ij}) \in C} \sigma_{ij} = \text{sgn } C$. \qed
Here is arguably the most important result on balance and switching. This is corollary 3.3 in [8].

**Theorem 1.3.2.** A signed digraph $\Delta$ is balanced if and only if there is $\zeta \in \{\pm\}^V$ such that every arc in $\zeta\Delta$ is positive.

**Proof.** If there is $\zeta \in \{\pm\}^V$ such that every arc in $\zeta\Delta$ is positive, then every circle in $\Delta$ is positive by proposition 1.3.1.

Conversely, suppose $\Delta$ is balanced. We will construct $\zeta$, a switching of $\Delta$ such that $\zeta\Delta$ is positive.

First we will prove this result specifically for trees.

**Lemma 1.3.3.** If $T$ is a tree, then there exists $\zeta$, a switching of $T$, such that every arc of $\zeta T$ is positive.

**Proof.** The proof is by induction.

The proposition is trivially true if the tree consists of a single vertex since then there are no arcs.

Suppose the tree has $k > 1$ vertices and that the result is true for trees with less than $k$ vertices. By proposition 1.1.3, $T$ has a degree-1 vertex $v$ incident with arc $a$ between $v$ and $w$ with sign $\sigma$. By lemma 1.1.7, $S = (V \setminus v, A \setminus a)$ is a tree on fewer than $k$ vertices. By assumption, there is $\zeta'$ a switching of $S$ such that every arc in $\zeta'S$ is positive. Let $\zeta$ be the switching of $T$ such that $\zeta_i = \zeta'_i$ if $i \neq v$ and $\zeta_v = \zeta'_w \sigma$. In $\zeta T$, the sign of $a$ is $\zeta_w \zeta_v \sigma = \zeta'_w (\zeta'_w \sigma) \sigma = +$. Any other arc $(i, j, \zeta_i \zeta_j \sigma_{ij})$ in $\zeta T$ is also positive because $\zeta_i \zeta_j \sigma_{ij} = \zeta'_i \zeta'_j \sigma_{ij} = +$ since $(i, j, \zeta'_i \zeta'_j \sigma_{ij}) \in S$, a positive tree by induction. Therefore every arc in $\zeta T$ is positive.\[\Box\]

Back to the proof of the theorem. Each connected component of $\Delta$ can be treated separately. So for the purposes of this proof we will assume that $\Delta$ is connected.

Since $\Delta$ is connected it contains a spanning-tree $T$ by proposition 1.1.4. By lemma 1.3.3, there is $\zeta$, a switching of $T$, such that $\zeta T$ is positive. Since $T$ and $\Delta$ have the same vertices, $\zeta$ is also a switching of $\Delta$. Every arc in $T$ will be positive in $\zeta \Delta$ so we just need to show that the arcs outside of $T$ are also positive in $\zeta \Delta$.

Let $(i, j, \sigma_{ij})$ be an arc in $\zeta \Delta$ not contained in $T$. By proposition 1.1.6, there is a path in $\zeta T$ between $i$ and $j$. This path together with $a$ forms a circle $C$ in $\zeta \Delta$. Since $\Delta$ is balanced, $C$ is positive in $\zeta \Delta$ by proposition 1.3.1. But every arc in the path is positive since $\zeta T$ is positive. Therefore the
sign of \( C \) is the sign of \( a \). Therefore \( \sigma_{ij} = + \). This shows that every arc in \( \zeta \Delta \) is positive, whether in \( T \) or not.
Chapter 2

Interaction Graphs

2.1 Introduction

Interaction graphs are signed digraphs derived from a special kind of function, called an activation function. They are used to represent a group of interacting components, such as a gene regulatory network.

In section 2.2, activation functions and their interaction graphs are defined.

Section 2.3 deals with a useful way of transforming activation functions and their interaction graphs and establishes the relationship between the interaction graphs of an activation function and its conjugate activation functions.

Next are the main results of this chapter. Section 2.4 deals with the question of when a given signed directed graph $\Delta$ is a global interaction graph. A general method to construct an activation function whose interaction graph is $\Delta$ is then detailed.

Section 2.5 deals with the same question for local interaction graphs and provides a general method to construct an activation function whose local interaction graph is a given signed digraph.

Finally, section 2.6 covers the question of when a given signed directed graph is an interaction graph of a threshold activation function.

2.2 Interaction Graphs

2.2.1 Activation Functions

Let $C_f$ be a finite set. For each $i \in C_f$, let $S_i$ be a non-empty finite set of integers and let $S_f = \prod_{i \in C_f} S_i$. An activation function is a function $f : S_f \to S_f$. The set $C_f$ is called the set of components of $f$. The set $S_f$ is the state space of $f$. We may also refer to $S_f$ as the states of $f$ and call elements of $S_f$ states. Each component function $f_j : S_f \to S_j$ is called a component activation function. For $x \in S_f$, $S_i$ will normally denote the set of possible values of $x_i$. 

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If \(|S_i| \leq 2\) for all \(i \in C_f\), then \(f\) is called a Boolean activation function.

The states of an activation function \(f\) will be ordered using the product ordering. That is for two states \(x, y \in S_f\), \(x \geq y\) if \(x_i \geq y_i\) for all \(i \in C_f\).

If \(f(x) = x\), then we call \(x\) a steady state of \(f\).

We use Hamming metric to measure distance between states \(x, y \in S_f\). That is, the distance between \(x\) and \(y\) is \(|x, y| := |\{i \in C_f \mid x_i \neq y_i\}|\).

### 2.2.2 Interaction Graphs of an Activation Function

For this section we are interested in how changing a single input of an activation function \(f\) affects the outputs of \(f\).

First we will establish some notation to aid in the analysis. Let \(f\) be an activation function, \(x \in S_f\), \(i \in C_f\), and \(a \in S_i\). Define \(x^{i \rightarrow a}\) to be the state in \(S_f\) such that

\[
(x^{i \rightarrow a})_j := x_j^{i \rightarrow a} = \begin{cases} 
  x_j & j \neq i \\
  a & j = i.
\end{cases}
\]  

(2.1)

The idea here is we are only changing the value of the \(i\)th component of \(x\).

We shall be comparing the outputs of \(f(x)\) and \(f(x^{i \rightarrow a})\) for different values of \(i\) and \(a\). We will also want to see if the outputs of \(f\) increase or decrease when varying the \(i\)th input.

First some notation. Recall that for any real number \(x\),

\[
\text{sgn } x = \begin{cases} 
  + & x > 0 \\
  0 & x = 0 \\
  - & x < 0.
\end{cases}
\]

For any \(x \in S_f\) and \(a \in S_i\), let

\[
\partial^{i \rightarrow a} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow a}) - f_j(x)][a - x_i].
\]  

(2.2)

Let us examine what this definition means. The value of \(\partial^{i \rightarrow a} f_j(x)\) signifies what effect changing the \(i\)th input of \(f\) has on the \(j\)th output of \(f\). Specifically, suppose \(\partial^{i \rightarrow a} f_j(x) = +\). Then \(\text{sgn}(f_j(x^{i \rightarrow a}) - f_j(x)) = \text{sgn}(a - x_i) \neq 0\). So either \(f_j(x^{i \rightarrow a}) < f_j(x)\) and \(a < x_i\), or \(f_j(x^{i \rightarrow a}) > f_j(x)\) and \(a > x_i\). Therefore increasing the \(i\)th input of \(f\) increased the \(j\)th output, or decreasing the \(i\)th input decreased the \(j\)th output. So the \(j\)th output of \(f\) changes in the same way as the \(i\)th input was changed. Similarly, if \(\partial^{i \rightarrow a} f_j(x) = -\), then \(\text{sgn}(f_j(x^{i \rightarrow y_i}) - f_j(x)) = - \text{sgn}(y_i - x_i) \neq 0\).
So in this case increasing the \(i\)th input of \(f\) decreased the \(j\)th output, or decreasing the \(i\)th input increased the \(j\)th output. So the value of \(f_j\) changed contrary to the change of \(x_i\). If \(\partial^{i\to a} f_j(x) = 0\), then \(f_j(x^{i\to a}) = f_j(x)\) regardless or what the value of \(a\) is, so changing \(x_i\) has no effect on \(f_j\).

Here is one fact about \(\partial^{i\to a} f_j(x)\) that will prove useful. In \(\partial^{i\to a} f_j(x)\) we start with \(x\) and change the \(i\)th input to \(a\). But what if we start with \(x^{i\to a}\) and change the \(i\)th input to \(x_i\)? It turns out we get the same sign.

**Proposition 2.2.1.** Given an activation function \(f\), \(x \in S_f\) and \(a \in S_i\), then
\[
\partial^{i\to x_i} f_j(x^{i\to a}) = \partial^{i\to a} f_j(x).
\]

**Proof.** This is a straightforward calculation based on equation (2.2).

\[
\partial^{i\to x_i} f_j(x^{i\to a}) = \text{sgn}[f_j([x^{i\to a}]^i\to x_i) - f_j(x^{i\to a})][x_i - a]
= \text{sgn}[f_j(x) - f_j(x^{i\to a})][x_i - a]
= \text{sgn}[-(f_j(x^{i\to a}) - f_j(x))][- (a - x_i)]
= \text{sgn}[f_j(x^{i\to a}) - f_j(x)][a - x_i]
= \partial^{i\to a} f_j(x).
\]

Now we will represent graphically how changing the inputs of an activation function affects the outputs. Let \(f\) be an activation function and \(x, y \in S_f\). The *local interaction graph of \(f\) at \(x\) in the direction of \(y\) is the signed digraph \(I_f(x, y)\) with vertex set \(C_f\) and with arc set \(\{(i, j, \sigma_{ij}) \mid \sigma_{ij} = \partial^{i\to y_j} f_j(x) \neq 0\}\).

For example, let the activation function \(f(x_1, x_2) = (1 - 2x_2^2, x_2(x_1 - 1)/2)\) where \(S_1 = \{\pm 1\}\) and \(S_2 = \{0, \pm 1\}\). Let us find the local interaction graph \(I_f((-1, -1), (1, 0))\). We just have to calculate \(\partial^{i\to y_j} f_j((-1, -1))\) for each possible \(i\) and \(j\) where \(y = (1, 0)\).

\[
\partial^{1\to 1} f_1((-1, -1)) = \text{sgn}[f_1((1, -1)) - f_1((-1, -1))][1 - (-1)] = \text{sgn}[1 - (-1)] = 0
\]
\[
\partial^{1\to 0} f_1((-1, -1)) = \text{sgn}[f_1((-0, 1)) - f_1((-1, -1))][0 - (-1)] = \text{sgn}[1 - (-1)] = +
\]
\[
\partial^{1\to -1} f_2((-1, -1)) = \text{sgn}[f_2((1, -1)) - f_2((-1, -1))][-1 - (-1)] = \text{sgn}[0 - 1] = -
\]
\[
\partial^{2\to 0} f_2((-1, -1)) = \text{sgn}[f_2((-0, 1)) - f_2((-1, -1))][0 - (-1)] = \text{sgn}[0 - 1] = -
\]

Therefore the arc set of \(I_f((-1, -1), (1, 0))\) is \(\{(2, 1, +), (1, 2, -), (2, 2, -)\}\). An illustration of \(I_f((-1, -1), (1, 0))\) is given in figure 2.1. In the same figure we also have \(I_f((1, 1), (-1, 0))\). You can easily check that \(\partial^{1\to -1} f_1((1, 1)) = \partial^{2\to 0} f_2((1, 1)) = 0, \partial^{1\to -1} f_2((1, 1)) = +\) and \(\partial^{2\to 0} f_1((1, 1)) = -\).
Figure 2.1: Local interaction graphs $\mathcal{I}_f((-1, -1), (1, 0))$ and $\mathcal{I}_f((1, 1), (-1, 0))$.

Figure 2.2: Global interaction graph $\mathcal{I}_f$

Since local interaction graphs have at most one arc from any vertex to another, they can be represented by an adjacency matrix. The $ij$th entry of $\text{Adj}(\mathcal{I}_f(x, y))$ is $\partial_i \rightarrow y \sigma_{ij}(x)$. We will abbreviate the $i$th row of this adjacency matrix by $\partial_i \rightarrow y := \text{sgn}[f(x_i \rightarrow y_i) - f(x)][y_i - x_i]$. So for the activation function $f(x_1, x_2) = (1 - 2x_2^2, x_2(x_1 - 1)/2)$,

$$\text{Adj}(\mathcal{I}_f((-1, -1), (1, 0))) = \begin{pmatrix} \partial_{-1}^{-1}f((-1, -1)) \\ \partial_{-1}^{0}f((-1, -1)) \end{pmatrix} = \begin{pmatrix} 0 & - \\ + & - \end{pmatrix},$$

$$\text{Adj}(\mathcal{I}_f((1, 1), (-1, 0))) = \begin{pmatrix} \partial_{1}^{-1}f((1, 1)) \\ \partial_{1}^{0}f((1, 1)) \end{pmatrix} = \begin{pmatrix} 0 & + \\ - & 0 \end{pmatrix}.$$

The global interaction graph of an activation function $f$, or just the interaction graph of $f$, is the union of all the local interaction graphs of $f$. Put another way, $\mathcal{I}_f$ is the signed digraph with vertex set $C_f$ and with arc set $\{(i, j, \sigma_{ij}) \mid \exists x \in S_f, a \in S_i \text{ such that } \sigma_{ij} = \partial_i \rightarrow a f_j(x) \neq 0\}$.

An illustration of the global interaction graph of the activation function $f(x_1, x_2) = (1 - 2x_2^2, x_2(x_1 - 1)/2)$ is given in figure 2.2. Notice that it includes all of the arcs from the two local interaction graphs above. Since $f_1$ does not depend on $x_1$ there will be no loops on vertex 1 in the graph. You can also check that $\partial_{2}^{-y_i}f_2(x) \leq 0$, so there is no positive loop on vertex 2.

2.2.3 Boolean Interaction Graphs

Given an activation function $f$, if $|S_i| \leq 2$ for all $i \in C_f$, then $f$ is called a Boolean activation function since each $x_i$ has exactly two possible values for each $i \in C_f$.

For Boolean activation functions, $\mathcal{I}_f(x)$ will denote the local interaction graph $\mathcal{I}_f(x, y)$ where $y$ is the state in $S_f$ such that $y_i \neq x_i$ for each $i \in C_f$. This is well defined since $|S_i| = 2$ for each
\(i \in C_f\). Also \(x^{i \rightarrow a}\) where \(a \in S_i\) and \(a \neq x_i\) will be denoted by \(x^{i \rightarrow}\) since the \(i\)th coordinate of \(x\) can only be changed to a single value that differs from \(x_i\).

Here is proposition 2.2.1 for Boolean activation functions with our new shorthand notation. The proof is practically identical.

**Proposition 2.2.2.** Given a Boolean activation function \(f\), if \(x \in \mathbb{S}_f\), then

\[
\partial^{i \rightarrow} f_j(x) = \partial^{i \rightarrow} f_j(x).
\]

### 2.2.4 A Restriction on Interaction Graphs

My definition of a local interaction graph is a bit looser than standard. Others, such as the definition in [4] and [5], compare the outputs of the activation function when only changing the value of a component by the smallest amount possible. That is, they only consider \(\partial^{i \rightarrow a} f_j(x)\) when \(a < x_i\) or \(a > x_i\). The relation \(a < b\) is the cover relation: \(b\) covers \(a\), i.e., \(a < x_i\) means that \(a < x_i\), and there is no \(b \in S_i\) such that \(a < b < x_i\). Similarly, \(a > x_i\) means that \(a > x_i\), and there is no \(b \in S_i\) such that \(a > b > x_i\).

With this in mind, we will define a restricted version of the interaction graphs. If \(a > x_i\), let 
\[
\partial^{i \rightarrow} f_j(x) := \partial^{i \rightarrow b} f_j(x) \quad \text{where } b > x_i.
\]
Similarly, if \(a < x_i\), let 
\[
\partial^{i \rightarrow} f_j(x) := \partial^{i \rightarrow b} f_j(x) \quad \text{where } b < x_i.
\]

The restricted local interaction graph of \(f\) at \(x\) in the direction of \(y\) is the signed digraph \(I^*_f(x,y)\) with vertex set \(C_f\) and with arc set \(\{(i,j,\sigma_{ij}) \mid \sigma_{ij} = \partial^{i \rightarrow y_i} f_j(x) \neq 0\}\).

For any states \(x\) and \(y\), let \(x^{\rightarrow y}\) be the state whose \(i\)th coordinate satisfies \((x^{\rightarrow y})_i < x_i\) if \(y_i < x_i\), \((x^{\rightarrow y})_i > x_i\) if \(y_i > x_i\), and \((x^{\rightarrow y})_i = x_i\) otherwise. This way, if \(z = x^{\rightarrow y}\), then

\[
x^{i \rightarrow y_i} = x^{i \rightarrow z_i}.
\]  

So \(\partial^{i \rightarrow y_i} f_j(x) = \partial^{i \rightarrow z_i} f_j(x)\). From this we get the following result immediately.

**Proposition 2.2.3.** \(I^*_f(x,y) = I_f(x, x^{\rightarrow y})\).

Many of our future results on interaction graphs also apply to restricted interaction graphs because of this fact.

Since Boolean interaction graphs can take at most two values on each of their components, there is no distinction between local interaction graphs and restricted ones in this case.

**Proposition 2.2.4.** If \(f\) is a Boolean activation function, then \(I^*_f(x,y) = I_f(x,y)\).
So we do not need to concern ourselves with restricted interaction graphs when working with Boolean activation functions.

The restricted global interaction graph, $\mathcal{I}_f^*$ is the union of all the restricted local interaction graphs of $f$. Some results about local interaction graphs, such as those in section 3.3.1 can be proved for the first kind of local interaction graphs, but are stronger if you prove them for restricted interaction graphs. You may suspect that the same is true for global interaction graphs. Fortunately, we have the following result.

**Proposition 2.2.5.** For an activation function $f$, $\mathcal{I}_f^* = \mathcal{I}_f$.

**Proof.** By proposition 2.2.3, $\mathcal{I}_f^* \subseteq \mathcal{I}_f$. So we just need to show the reverse containment. Since both interaction graphs have the same vertex set, we only need to show that any arc contained in $\mathcal{I}_f$ is also contained in $\mathcal{I}_f^*$. To do this we will use the following lemma.

**Lemma 2.2.6.** Given an activation function $f$.

1. If the arc $(i, j, \rightarrow) \notin \mathcal{I}_f$, then $f_j$ is weakly decreasing in the $i$th component. I.e., if $a > x_i$, then $f_j(x) \geq f_j(x^{i \mapsto a})$.

2. If the arc $(i, j, \leftarrow) \notin \mathcal{I}_f$, then $f_j$ is weakly increasing in the $i$th component. I.e., if $a > x_i$, then $f_j(x) \leq f_j(x^{i \mapsto a})$.

**Proof.** For the first item, suppose the arc $(i, j, \rightarrow) \notin \mathcal{I}_f$. Let $x = x_1$. So $\partial^{i \mapsto +} f_j(x_1) \neq +$. Using equation (2.2), since $a > x_i$,

$$\partial^{i \mapsto +} f_j(x) = \sgn[f_j(x^{i \mapsto a}) - f_j(x)][a - x_i] = \sgn[f_j(x^{i \mapsto a}) - f_j(x)].$$

So $f_j(x^{i \mapsto a}) \leq f_j(x_1)$.

Now let $x_2 = x^{i \mapsto a}$. Again, since $(i, j, \rightarrow) \notin \mathcal{I}_f$, $\partial^{i \mapsto +} f_j(x_2) \neq +$. Since $a > x_i$, $\partial^{i \mapsto +} f_j(x) = \sgn[f_j(x_2^{i \mapsto a}) - f_j(x_2)] \neq +$. So $f_j(x_2^{i \mapsto a}) \leq f_j(x_2)$ again.

Now let $x_3 = x_2^{i \mapsto a}$ and repeat this process. Eventually $x_n = x^{i \mapsto a}$ after some number of steps and $f_j(x_{k+1}) \leq f_j(x_k)$ for all $k \in [n - 1]$. Therefore $f_j(x_n) \leq f_j(x_1)$, i.e $f_j(x^{i \mapsto a}) \leq f_j(x)$.

The second item follows by a similar argument. If $(i, j, \leftarrow) \notin \mathcal{I}_f$, then $\partial^{i \mapsto -} f_j(x_1) \neq -$. Since $a > x_i$, $\partial^{i \mapsto -} f_j(x) = \sgn[f_j(x^{i \mapsto a}) - f_j(x_1)]$. So $f_j(x^{i \mapsto a}) \geq f_j(x_1)$.

Now let $x_2 = x^{i \mapsto a}$. Again, since $(i, j, \leftarrow) \notin \mathcal{I}_f$, $\partial^{i \mapsto -} f_j(x_2) \neq -$. Since $a > x_i$, $\partial^{i \mapsto -} f_j(x) = \sgn[f_j(x_2^{i \mapsto a}) - f_j(x_2)]$. So $f_j(x_2^{i \mapsto a}) \geq f_j(x_2)$ again.

Now let $x_3 = x_2^{i \mapsto a}$ and repeat this process. Eventually $x_n = x^{i \mapsto a}$ after some number of steps and $f_j(x_{k+1}) \geq f_j(x_k)$ for all $k \in [n - 1]$. Therefore $f_j(x_n) \geq f_j(x_1)$, i.e $f_j(x^{i \mapsto a}) \leq f_j(x)$. □
Back to the proof of proposition 2.2.5 Suppose the arc \((i, j, +) \in \mathcal{I}_f\). Then \(\partial^{i \rightarrow y_i} f_j(x) = +\) for some \(x, y \in \mathbb{S}_f\). If \(y_i > x_i\), using equation 2.2,

\[
\partial^{i \rightarrow y_i} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow y_i}) - f_j(x)][y_i - x_i] = \text{sgn}[f_j(x^{i \rightarrow y_i}) - f_j(x)].
\]

Since \(\partial^{i \rightarrow y_i} f_j(x) = +\), \(f_j(x^{i \rightarrow y_i}) > f_j(x)\). By lemma 2.2.6 if \((i, j, +) \notin \mathcal{I}_f^*\) then \(f_j(x^{i \rightarrow y_i}) \leq f_j(x)\), a contradiction. Therefore \((i, j, +) \in \mathcal{I}_f^*\). If instead \(y_i < x_i\), then

\[
\partial^{i \rightarrow y_i} f_j(x) = -\text{sgn}[f_j(x^{i \rightarrow y_i}) - f_j(x)] = +. 
\]

So \(f_j(x^{i \rightarrow y_i}) < f_j(x)\). But if \((i, j, +) \notin \mathcal{I}_f^*\) then \(f_j(x^{i \rightarrow y_i}) \geq f_j(x)\) by lemma 2.2.6 Again, this is a contradiction. Therefore \((i, j, +) \in \mathcal{I}_f^*\).

If the arc \((i, j, -) \in \mathcal{I}_f\), we can show that \((i, j, -) \in \mathcal{I}_f^*\) using essentially the same argument. Since \((i, j, -) \in \mathcal{I}_f\), \(\partial^{i \rightarrow y_i} f_j(x) = +\) for some \(x, y \in \mathbb{S}_f\). If \(y_i > x_i\), using equation 2.2,

\[
\partial^{i \rightarrow y_i} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow y_i}) - f_j(x)].
\]

Since \(\partial^{i \rightarrow y_i} f_j(x) = -\), \(f_j(x^{i \rightarrow y_i}) < f_j(x)\). By lemma 2.2.6 if \((i, j, -) \notin \mathcal{I}_f^*\) then \(f_j(x^{i \rightarrow y_i}) \geq f_j(x)\), a contradiction. Therefore \((i, j, -) \in \mathcal{I}_f^*\). If instead \(y_i < x_i\), then \(\partial^{i \rightarrow y_i} f_j(x) = -\text{sgn}[f_j(x^{i \rightarrow y_i}) - f_j(x)] = +.\) So \(f_j(x^{i \rightarrow y_i}) > f_j(x)\). But if \((i, j, -) \notin \mathcal{I}_f^*\) then \(f_j(x^{i \rightarrow y_i}) \leq f_j(x)\) by lemma 2.2.6 Again, this is a contradiction. Therefore \((i, j, -) \in \mathcal{I}_f^*\).

So every arc in \(\mathcal{I}_f\) is also contained in \(\mathcal{I}_f^*\). Therefore \(\mathcal{I}_f = \mathcal{I}_f^*\).

So when working with global interaction graphs, we can use their more convenient standard definition and all results we prove will also be true for the restricted global interaction graph.

### 2.3 Conjugate Activation Functions

Given an activation function \(f\), let \(\phi : C_f \rightarrow W\) be a bijection and \(\zeta \in \{\pm\}^{C_f}\). In particular, \(\lambda = (\phi, \zeta)\) is a switching isomorphism of \(\mathcal{I}_f\) and of \(\mathcal{I}_f(x, y)\) for any \(x, y \in \mathbb{S}_f\). We will also call \(\lambda\) a switching isomorphism of \(f\) for reasons which will become clear in the next section. Also note that \(W = \phi C_f\), since it is the codomain of \(\phi\).

There is a natural way for \(\lambda\) to transform states in \(\mathbb{S}_f\). First we will establish how the bijection \(\phi\) transforms states in \(\mathbb{S}_f\). If \(x \in \mathbb{S}_f\), \(\phi x\) is the state whose components are in \(\phi C_f\) and are such that \(\phi_j x := (\phi x)_j = x_{\phi^{-1} - j}\), or \(\phi_{\phi} x = x_i\).

Let \(\circ\) be the Hadamard component-wise product. Define \(\lambda x = (\phi, \zeta) x := \phi(\zeta \circ x)\). Since \(\circ\) is the only kind of product between states and sign-vectors, we normally suppress the \(\circ\) notation.
Note that even though \( \zeta x \) may not be in \( S_f \), its components are \( C_f \), so it can still be transformed by \( \phi \) as defined above. So if \( j \in \phi C_f \), then

\[
\lambda_j x := (\lambda x)_j = \phi_j(\zeta x) = (\zeta x)_{\phi^{-1}j} = \zeta_{\phi^{-1}j}x_{\phi^{-1}j}.
\]

Or more cleanly,

\[
\lambda_{\phi i} x = \zeta_i x_i. \quad (2.4)
\]

So when the switching isomorphism \( \lambda = (\phi, \zeta) \) transforms states in \( S_f \) this way, what is the codomain? Let \( \zeta_i S_i := \{ \zeta_i x \mid x \in X_i \} \). For \( j \in \phi C_f \), define \( S_j := \zeta_{\phi^{-1}j}S_{\phi^{-1}j} \). More cleanly, \( S_{\phi i} := \zeta_i S_i \). Let \( \lambda S_f := \prod_{j \in W} S_j \). This way, \( \lambda S_f \) is the codomain of \( \lambda \), i.e. \( \lambda : S_f \to \lambda S_f \).

Elements of \( \{ \pm \} C_f \) can be transformed by \( \phi \) the same way states are, since its components are also indexed by \( C_f \). That is \( \phi \zeta \in \{ \pm \} \phi C_f \) where \( \phi_j \zeta := (\phi \zeta)_j = \zeta_{\phi^{-1}j} \), or \( \phi_{\phi i} \zeta = \zeta_i \). It follows that

\[
\lambda_j x = \zeta_{\phi^{-1}j}x_{\phi^{-1}j} = (\phi_j \zeta)(\phi_j x).
\]

Combining this with the previous definition of \( \lambda x \) we get

\[
\lambda x = \phi(\zeta x) = (\phi \zeta)(\phi x). \quad (2.5)
\]

By this definition, the transformation \( \lambda : S_f \to \lambda S_f \) is a bijection with inverse transformation \( \lambda^{-1} = (\phi^{-1}, \phi \zeta) \) since

\[
(\phi^{-1}, \phi \zeta)[(\phi, \zeta) x] = (\phi^{-1}, \phi \zeta)[(\phi \zeta)(\phi x)] = \phi^{-1}[(\phi \zeta)(\phi \zeta)(\phi x)] = \phi^{-1}[\phi x] = x.
\]

If \( \lambda \) is a switching isomorphism of an activation function \( f \), the \( \lambda \)-conjugate activation function of \( f \) is the activation function \( f^\lambda : \lambda S_f \to \lambda S_f \) with the property that

\[
f^\lambda(\lambda x) = \lambda f(x). \quad (2.6)
\]

The conjugate activation function is defined this way so that figure 2.3 is a commutative diagram. Note that \( C_{f^\lambda} = \phi C_f \) and \( S_{f^\lambda} = \lambda S_f \). Also

\[
f_{\phi i}^\lambda(\lambda x) = [f^\lambda(\lambda x)]_{\phi i} = [\lambda f(x)]_{\phi i} = \zeta_i f_i(x). \quad (2.7)
\]

If \( \zeta \in \{ \pm \} C_f \), then \( f^\zeta = f^\lambda \) where \( \lambda = (\phi, \zeta) \) and \( \phi \) is the identity function on \( C_f \). Similarly, if \( \phi : C_f \to \phi C_f \), then \( f^\phi = f^\lambda \) where \( \lambda = (\{ \pm \} C_f, \phi) \). Two activation functions \( f \) and \( g \) are called conjugate if \( g = f^\lambda \) for some switching isomorphism \( \lambda \).
2.3.1 Interaction Graphs of Conjugate Activation Functions

It turns out that the interaction graphs of an activation function and its conjugate activation functions are related in a nice way. This will justify why we are also calling a switching isomorphism of $I_f$ a switching isomorphism of the activation function $f$ itself.

**Theorem 2.3.1.** Let $f$ be an activation function, let $x, y \in S_f$ and let $\lambda$ be a switching isomorphism of $f$. Then:

1. $I_{f\lambda}(\lambda x, \lambda y) = \lambda[I_f(x, y)]$, i.e., $I_f(x, y)$ and $I_{f\lambda}(\lambda x, \lambda y)$ are switching isomorphic.

2. $I^\ast_{f\lambda}(\lambda x, \lambda y) = \lambda[I^\ast_f(x, y)]$, i.e., $I^\ast_f(x, y)$ and $I^\ast_{f\lambda}(\lambda x, \lambda y)$ are switching isomorphic.

3. $I_{f\lambda} = \lambda I_f$, i.e., $I_f$ and $I_{f\lambda}$ are switching isomorphic.

4. $I^\ast_{f\lambda} = \lambda I^\ast_f$, i.e., $I^\ast_f$ and $I^\ast_{f\lambda}$ are switching isomorphic.

**Proof.** To prove the first item, we will show that $\varphi^{\phi i \rightarrow \lambda \phi j} f^\lambda_{\phi j}(\lambda x) = \zeta^{\phi i \rightarrow \lambda \phi j} f_j(x)\zeta_j$, so that $(i, j, \sigma_{ij}) \in I_f(x, y)$ if and only if $(\phi i, \phi j, \zeta_i \sigma_{ij} \zeta_j) \in I_{f\lambda}(\lambda x, \lambda y)$. Note that

$$\varphi^{\phi i \rightarrow \lambda \phi j} f^\lambda_{\phi j}(\lambda x) = \text{sgn}[f^\lambda_{\phi j}(\lambda x^{\phi i \rightarrow \lambda \phi j} y) - f^\lambda_{\phi j}(\lambda x)][\lambda \phi y - \lambda \phi x].$$

To complete a calculation, we need the following lemma.

**Lemma 2.3.2.** If $x, y \in S_f$, then $\lambda(x^{i \rightarrow y}) = (\lambda x)^{\phi i \rightarrow \lambda \phi j} y$. 
Proof. We just apply equation (2.4).

\[
\lambda_{\phi j}(x^{i \rightarrow y_i}) = \zeta_j x_j^{i \rightarrow y_i}
\]

= \begin{cases} 
\zeta_j x_j & j \neq i \\
\zeta_i y_i & j = i 
\end{cases}

= \begin{cases} 
\lambda_{\phi j} x_{\phi j} & \phi j \neq \phi i \\
\lambda_{\phi i} y_{\phi j} & \phi j = \phi i 
\end{cases}

= (\lambda x)^{\phi i \rightarrow \lambda_{\phi i} y}

We are now ready to complete the proof. Using this lemma, equation (2.7) and equation (2.4),

\[
\partial^{\phi i \rightarrow \lambda_{\phi i} y} f_{\phi j}(\lambda x) = \text{sgn}[f_{\phi j}^\lambda(\lambda x^{i \rightarrow y_i}) - f_{\phi j}^\lambda(\lambda x)][\lambda_{\phi i} y - \lambda_{\phi i} x]
\]

\[
= \text{sgn}[f_{\phi j}^\lambda(\lambda x^{i \rightarrow y_i}) - f_{\phi j}^\lambda(\lambda x)][\lambda_{\phi i} y - \lambda_{\phi i} x]
\]

\[
= \text{sgn}[\zeta_j f_j(x^{i \rightarrow y_i}) - \zeta_j f_j(x)][\zeta_i y_i - \zeta_i x_i]
\]

\[
= \text{sgn}[\zeta_j f_j(x^{i \rightarrow y_i}) - \zeta_j f_j(x)][\zeta_i y_i - \zeta_i x_i]
\]

\[
= \text{sgn}[\zeta_j f_j(x^{i \rightarrow y_i}) - \zeta_j f_j(x)][\zeta_i y_i - \zeta_i x_i]
\]

\[
= \text{sgn}[\zeta_j f_j(x^{i \rightarrow y_i}) - \zeta_j f_j(x)][\zeta_i y_i - \zeta_i x_i]
\]

\[
= \text{sgn}[\zeta_j f_j(x^{i \rightarrow y_i}) - \zeta_j f_j(x)][\zeta_i y_i - \zeta_i x_i]
\]

\[
= \zeta_i \partial^{\phi i \rightarrow \lambda_{\phi i} y} f_j(x) \zeta_j
\]

Now if \((i, j, \sigma_{ij}) \in \mathcal{I}_f(x, y)\), then \(\partial^{i \rightarrow y_i} f_j(x) = \sigma_{ij}\). Therefore \((\phi i, \phi j, \zeta_i \sigma_{ij} \zeta_j) \in \mathcal{I}_{f^\lambda}\) if and only if \((i, j, \sigma_{ij}) \in \mathcal{I}_f(x, y)\).

For the second item, we start with the left hand side. By proposition 2.2.3, \(\mathcal{I}_{f^\lambda}(\lambda x, \lambda y) = \mathcal{I}_{f^\lambda}(\lambda x, (\lambda x)^{\rightarrow \lambda y})\). Using the same proposition, the right hand side \(\lambda[\mathcal{I}_f(x, y)] = \lambda[\mathcal{I}_f(x, (x^{\rightarrow y}))]\).

By the first item in this theorem, \(\lambda[\mathcal{I}_f(x, (x^{\rightarrow y}))] = \mathcal{I}_{f^\lambda}(\lambda x, (\lambda x)^{\rightarrow \lambda y})\). So to prove the second item, we will use the following lemma.

Lemma 2.3.3. Given states \(x, y\) and a switching isomorphism \(\lambda\) of an activation function \(f\). Then \((\lambda x)^{\rightarrow \lambda y} = \lambda(x^{\rightarrow y})\).

Proof. Let \(z = x^{\rightarrow y}\) and \(z' = (\lambda x)^{\rightarrow \lambda y}\). We want to show that \(\lambda z = z'\). We will show this component-wise, i.e., that \(\lambda_{\phi i} z = z'_{\phi i}\) for all \(i \in C_f\).
First suppose that \( y_i > x_i \). Then \( z_i \geq x_i \), by definition. If \( \zeta_i = + \), then \( \zeta_i y_i > \zeta_i x_i \). By equation (2.4), \( \lambda_{\phi_i} y > \lambda_{\phi_i} x \). So \( z'_i > \lambda_{\phi_i} x \), by definition. But since \( z_i \geq x_i \), \( \zeta_i z_i > \zeta_i x_i \). So by equation (2.4), \( \lambda_{\phi_i} z > \lambda_{\phi_i} x \). Therefore \( \lambda_{\phi_i} z > z'_i \).

Similarly, if \( y_i < x_i \), then \( z_i \leq x_i \) by definition. If \( \zeta_i = + \), then \( \zeta_i y_i < \zeta_i x_i \). By equation (2.4), \( \lambda_{\phi_i} y < \lambda_{\phi_i} x \). So \( z'_i < \lambda_{\phi_i} x \), by definition. But since \( z_i \leq x_i \), \( \zeta_i z_i < \zeta_i x_i \). So by equation (2.4), \( \lambda_{\phi_i} z < \lambda_{\phi_i} x \). Therefore \( \lambda_{\phi_i} z = z'_i \).

We have shown that \( \lambda_{\phi_i} z = z'_i \) for all \( i \in C_f \). Therefore \( \lambda z = z' \), or \( (\lambda x) \rightarrow^\lambda y = \lambda (x \rightarrow^\lambda y) \).}

By lemma 2.3.3, \( I_{f \lambda}(\lambda x, (\lambda x) \rightarrow^\lambda y) = I_{f \lambda}(\lambda x, \lambda (x \rightarrow^\lambda y)) \). By the first item of this theorem, \( I_{f \lambda}(\lambda x, \lambda (x \rightarrow^\lambda y)) = \lambda[I_f(x, x \rightarrow^\lambda y)] \). Therefore, \( I_{f \lambda}(\lambda x, \lambda y) = \lambda[I_f(x, y)] \) by proposition 2.2.3 since \( I_{f \lambda}(\lambda x, (\lambda x) \rightarrow^\lambda y) = \lambda[I_f(x, x \rightarrow^\lambda y)] \). This proves the second part of the theorem.

Finally the same results hold for the global interaction graphs since they are the unions of local interaction graphs. This proves the third and fourth items.

### 2.4 Forbidden Interaction Graphs

Can every signed digraph be an interaction graph? More specifically, given any signed digraph \( \Delta \), is there an activation function \( f \) such that \( I_f = \Delta \)? Is there an activation function \( g \) such that \( \Delta = I_g(x, y) \) for some \( x, y \in S_g \)? If there is no such activation function, what restrictions are there on the structure of \( \Delta \) to insure that it is an interaction graph? These are the questions that motivate the next few sections.

It turns out that every signed digraph is indeed the interaction graph of some activation function. However, there are restrictions on what kind of signed directed graphs can be the interaction graph of a Boolean activation function, as the following result shows.

**Theorem 2.4.1.** Let \( f \) be a Boolean activation function. If both \( (i, j, +) \), \( (i, j, -) \) \( \in I_f \), then the in-degree of \( j \) is at least four.

**Proof.** Suppose that \( (i, j, +) \), \( (i, j, -) \) \( \in I_f \) and the in-degree of \( j \) is less than four. Then all arcs into \( j \) come only from at most two vertices, \( i \) and \( k \). That means \( \partial^i \rightarrow f_j(x) = 0 \) for \( i' \neq i, k \), i.e.,
function graph of a ternary function. That is an activation function at least four. It turns out that this condition is also sufficient, which will be proved shortly. 

There is a state \( x \) such that \( \partial^i \to f_j(x) = + \) because \( (i, j,+) \in \mathcal{I}_f \). Let \( S_i = \{x, x'_i\} \). Since \( \partial^i \to f_j(x_i, x_k) = \partial^i \to f_j(x'_i, x_k) \) by proposition 2.2.2, we may assume for the sake of argument that \( x_i < x'_i \). Putting this assumption into equation (2.2),

\[
\partial^i \to f_j(x_i, x_k) = \text{sgn}[f_j(x'_i, x_k) - f_j(x_i, x_k)](x'_i - x_i) = \text{sgn}[f_j(x'_i, x_k) - f_j(x_i, x_k)].
\]

Since \( \partial^i \to f_j(x_i, x_k) = + \), \( f_j(x'_i, x_k) > f_j(x_i, x_k) \). Also \( S_j = \{z, z'\} \) where \( z < z' \) since \( f \) is Boolean. So \( f_j(x'_i, x_k) = z' \) and \( f_j(x_i, x_k) = z \).

Similarly, since \( (i, j,-) \in \mathcal{I}_f \), there is a state \( y = (y_i, y_k) \) such that \( \partial^i \to f_j(y) = - \). Now \( y_k \neq x_k \) since otherwise \( \partial^i \to f_j(y) \) would be the same as \( \partial^i \to f_j(x_i, x_k) \) or \( \partial^i \to f_j(x'_i, x_k) \) which are both positive. So let \( y_k = x'_k \). Now by a similar calculation to that above, it follows that \( f_j(x_i, x_k') > f_j(x'_i, x'_k) \), so \( f_j(x_i, x'_k) = z' \) and \( f_j(x'_i, x'_k) = z \). Based on this,

\[
\partial^{k-} f_j(x_i, x_k) = \text{sgn}[f_j(x_i, x'_k) - f_j(x_i, x_k)](x'_k - x_k) = \text{sgn}[x'_k - x_k].
\]

Similarly \( \partial^{k-} f_j(x'_i, x_k) = -\text{sgn}[x'_k - x_k] \). Therefore \( \partial^{k-} f_j(x_i, x_k) = -\partial^{k-} f_j(x'_i, x_k) \), and so both \( (k, j,+) \) and \( (k, j,-) \) are in \( \mathcal{I}_f \), so the in-degree of \( j \) is four.

This theorem tells us that for a signed digraph \( \Delta \) to be an interaction graph of a Boolean activation function, it is necessary that if both \( (i, j,+), (i, j,-) \in \Delta \), then the in-degree of \( j \) is at least four. It turns out that this condition is also sufficient, which will be proved shortly.

What if the Boolean restriction is removed? It turns out that every signed digraph is an interaction graph of a ternary function. That is an activation function \( f: \{0, \pm 1\}^C_j \to \{0, \pm 1\}^C_j \). In fact we can be even more specific.

**Theorem 2.4.2.** Given a signed digraph \( \Delta \). There is an activation function \( f \) such that \( \mathcal{I}_f = \Delta \), and:

1. each component activation function \( f_i \) is a polynomial,

2. \( S_i = \{0, \pm 1\} \) if there is \( j \in C_f \) such that \( (i, j,+) \), \( (i, j,-) \in \Delta \) and the in-degree of \( j \) is two or three,

3. \( S_i = \{\pm 1\} \) otherwise.
As an immediate corollary to this theorem, the “in-degree four” property is sufficient for a signed digraph $\Delta$ to be an interaction graph of a Boolean activation function.

**Corollary 2.4.3.** Let $\Delta$ be a signed digraph with the property that if both $(i, j, +), (i, j, -) \in \Delta$, then the in-degree of $j$ is at least four. Then there is a Boolean activation function $f$ such that $\Delta = \mathcal{I}_f$, and each $f_j$ is a polynomial.

**Proof.** If there are no vertices $i \in C_f$ with the property that $(i, j, +), (i, j, -) \in \Delta$ and the in-degree of $j$ is two or three, then theorem 2.4.2 says there is an activation function $f$ such that $\mathcal{I}_f = \Delta$, each $f_i$ is a polynomial, and $S_i = \{\pm 1\}$ for all $i \in C_f$, i.e., $f$ is a Boolean activation function.

**Proof of theorem 2.4.2.** For each $j \in V(\Delta)$, we will provide a method to construct a component activation polynomial $f_j$ of an activation function $f$ such that $\mathcal{I}_f = \Delta$. We will do this by showing that $j$ has the same in-star in $\mathcal{I}_f$ as in $\Delta$, where the in-star of $j$ is the set of all arcs into $j$. For brevity, we will call the in-star of $j$ in $\mathcal{I}_f$ the in-star of $f_j$.

Every component activation polynomial $f_j$ we will use in the proof maps from $\{0, \pm 1\}^{C_f}$ to $\{\pm 1\}$. This way $f_j$ can be used when $S_j = \{\pm 1\}$. We also need to insure that if $(i, j, \sigma_{ij}) \in \Delta$, then there is a state $x \in \mathcal{S}_f$ such that $x_i = -1$ and $\partial^{i \rightarrow 1} f_j(x) = \sigma_{ij}$ so that if $0 \notin S_i$, $\partial^{i \rightarrow 1} f_j(x)$ is well defined for $f_j$.

Consider the in-star of $j$ in $\Delta$. The in-star contains either no parallel arcs, a single pair of parallel arcs or, multiple pairs of parallel arcs. We will break down the proof by dealing with these three cases separately.

First let us deal with the case when there are no parallel arcs in the in-star of $j$ in $\Delta$. There may be no arcs into $j$ in $\Delta$. Then let $f_j$ be the constant polynomial $f_j(x) = 1$. This does not depend on any of its inputs so the in-star of $f_j$ will be empty.

Now we provide a way to add an additional non-parallel arc into $j$ in the in-star of a seed component activation function $g_j$ in the interaction graph. We will construct a component activation function $f_j$ whose in-star includes the in-star of $g_j$ and the additional arc. Also, if $g_j$ is a polynomial, then the new function $f_j$ will also be a polynomial. Moreover, the values of the initial component of the new arc need not include 0. This lemma can then be iterated to add as many additional non-parallel arcs as we wish.

**Lemma 2.4.4.** Let $g_j : \{0, \pm 1\}^{C_g} \rightarrow \{\pm 1\}$ be a component activation function that is not identically $-1$ and does not depend on input $i \in C_g$. For $\sigma \in \{\pm \}$, let $f_j : \{0, \pm 1\}^{C_g} \rightarrow \{\pm 1\}$ be the
component activation function

\[ f_j(x) := \frac{1}{2}(g_j(x) + 1)(\sigma x_i + x_i^2) - 1 = \begin{cases} 
1 & \text{if } g_j(x) = \sigma x_i = 1 \\
-1 & \text{otherwise.}
\end{cases} \quad (2.8) \]

Then:

1. the in-star of \( f_j \) is \( T \cup (i, j, \sigma) \) where \( T \) is the in-star of \( g_j \).

2. there is \( x \in S_f \) such that \( x_i = -1 \) and \( \sigma = \partial^{i \rightarrow 1} f_j(x) \).

Proof. We have to show that the in-star of \( f_j \) contains all arcs in the set \( T \cup (i, j, \sigma) \) and nothing more.

Let us start by showing that the in-star of \( f_j \) contains \((i, j, \sigma)\) and there is \( x \in S_g \) such that \( x_i = -1 \) and \( \sigma = \partial^{i \rightarrow 1} f_j(x) \). Since \( g_j \) is not identically \(-1\), we know that there is a state \( x \) such that \( g_j(x) = 1 \). Since \( g_j \) does not depend on input \( i \) we may assume that \( x_i = -1 \). This simplifies equation (2.2) slightly since

\[
\partial^{i \rightarrow 1} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow 1}) - f_j(x)][1 - x_i] \\
= \text{sgn}[f_j(x^{i \rightarrow 1}) - f_j(x)][1 - (-1)] \\
= \text{sgn}[f_j(x^{i \rightarrow 1}) - f_j(x)].
\]

If \( \sigma = + \), then \( \sigma x_i = -1 \). So by the definition of \( f_j \), \( f_j(x) = -1 \). Since \( g_j \) does not depend on input \( i \), \( g_j(x) = g_j(x^{i \rightarrow 1}) = 1 = \sigma x_i^{i \rightarrow 1} \). So by the definition of \( f_j \) again, \( f_j(x^{i \rightarrow 1}) = 1 \). Therefore

\[
\partial^{i \rightarrow 1} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow 1}) - f_j(x)] \\
= \text{sgn}[1 + (-1)] \\
= +.
\]

Similarly, if \( \sigma = - \), then \( \partial^{i \rightarrow 1} f_j(x) = \text{sgn}[1 - 1] = - \). Either way \( \partial^{i \rightarrow 1} f_j(x) = \sigma \), so \((i, j, \sigma)\) is indeed contained in the in-star of \( f_j \) and the state \( x \) satisfies our \( \{ \pm 1 \} \) input requirements.

Next, let us show that every arc in \( T \), the in-star of \( g_j \), is also contained in the in-star of \( f_j \). If \((k, j, \sigma_{kj}) \in T \), then there is a state \( x \) such that \( \partial^{k \rightarrow a} g_j(x) = \sigma_{kj} \). Since \( g_j \) does not depend on input \( i \), we may assume that \( \sigma x_i = 1 \). If \( g_j(x) = 1 \), then \( f_j(x) = 1 \) by definition. Similarly, if \( g_j(x) = -1 \), then \( f_j(x) = -1 \). Therefore \( f_j(x) = g_j(x) \) if \( \sigma x_i = 1 \). So \( \partial^{k \rightarrow a} g_j(x) = \partial^{k \rightarrow a} f_j(x) = \sigma_{kj} \).

Therefore \((k, j, \sigma_{kj})\) is also contained in the in-star of \( f_j \). It follows that in-star of \( f_j \) contains every arc in \( T \).
Next we have to show that every arc in the in-star of \( f_j \) is either \((i, j, \sigma)\), or is contained in the in-star of \( g_j \). Suppose the in-star of \( f_j \) contains \((k, j, \sigma_{kj})\), i.e., \( \partial^{k \rightarrow a} f_j(x) = \sigma_{kj} \). This means that \( f_j(x) \neq f_j(x^{k \rightarrow a}) \), so either \( f_j(x) = 1 \) or \( f_j(x^{k \rightarrow a}) = 1 \).

First suppose that \( f_j(x) = 1 \). Then \( g_j(x) = \sigma x_i = 1 \) by the definition of \( f_j \). Recall that if \( \sigma x_i = 1 \), then \( f_j(x) = g_j(x) \). So if \( k \neq i \),

\[
\partial^{k \rightarrow a} g_j(x) = \text{sgn}[g_j(x^{k \rightarrow a}) - g_j(x)][a - x_k] = \text{sgn}[f_j(x^{k \rightarrow a}) - f_j(x)][a - x_k] = \partial^{k \rightarrow a} f_j(x) = \sigma_{kj}.
\]

Therefore \((k, j, \sigma_{kj}) \in T \). Suppose \( k = i \), i.e., \((i, j, \sigma_{ij})\) is in the in-star of \( f_j \). Since \( \sigma x_i = 1 \), \( \sigma a \neq 1 \) because \( a \neq x_i \). So by the definition of \( f_j \), \( f_j(x^{i \rightarrow a}) = -1 \). Therefore

\[
\sigma_{ij} = \partial^{i \rightarrow a} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow a}) - f_j(x)][a - x_i] = \text{sgn}[-1 - 1][a - x_i] = - \text{sgn}[a - x_i].
\]

If \( \sigma = + \), then \( x_i = \sigma x_i = 1 \) and \( a = -1 \) or \( a = 0 \), so \( a - x_i < 0 \). It follows that \( \sigma_{ij} = + \). Similarly, if \( \sigma = - \), then \( x_i = -1 \). Since \( a \) is 0 or 1, \( a - x_i > 0 \). Again \( \sigma_{ij} = - \). In all cases, \( \sigma_{ij} = \sigma \). So if the in-star of \( f_j \) contains \((i, j, \sigma_{ij})\), then \((i, j, \sigma_{ij}) = (i, j, \sigma)\) and that is the only possibility for arcs from \( i \) to \( j \).

If we had instead assumed that \( f_j(x^{k \rightarrow a}) = 1 \), then \( g_j(x^{k \rightarrow a}) = \sigma x_i^{k \rightarrow a} = 1 \). If \( k \neq i \), then \( \sigma x_i^{k \rightarrow a} = \sigma x_i = 1 \), and \( g_j(x) = f_j(x) \) again. So \( \partial^{k \rightarrow a} g_j(x) = \sigma_{kj} \) as before. If \( k = i \), then \( \sigma x_i^{k \rightarrow a} = \sigma x_i^{i \rightarrow a} = \sigma a = 1 \) and \((i, j, \sigma_{ij})\) is contained in the in-star of \( f_j \). So

\[
\sigma_{ij} = \partial^{i \rightarrow a} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow a}) - f_j(x)][a - x_i] = \text{sgn}[1 - (-1)][a - x_i] = \text{sgn}[a - x_i].
\]

If \( \sigma = + \), then \( a = \sigma a = 1 \) and \( x_i = -1 \) or \( x_i = 0 \), so \( a - x_i > 0 \). Therefore \( \sigma_{ij} = + \). Similarly, if \( \sigma = - \), then \( a = -1 \), \( x_i \) is 0 or 1, so \( a - x_i < 0 \) and \( \sigma_{ij} = - \). Again \( \sigma_{ij} = \sigma \) in all cases. So if the
in-star of \( f_j \) contains \((i, j, \sigma_{ij})\), then \((i, j, \sigma_{ij}) = (i, j, \sigma)\) as before and that is the only possibility for arcs from \( i \) to \( j \). Therefore every arc in the in-star of \( f_j \) is either \((i, j, \sigma)\) or contained in \( T \), the in-star of \( g_j \). And so the in-star of \( f_j \) is \( T \cup (i, j, \sigma) \).

If we want a component activation polynomial \( f_j \) whose in-star contains a single arc, we can start with the constant polynomial \( g_j(x) = 1 \) and apply lemma 2.4.4 to get a suitable component activation polynomial \( f_j \). If we want to construct a component activation polynomial \( f_j \) whose in-star contains as many non-parallel arcs of any sign that we wish, then we can apply lemma 2.4.4 iteratively to \( g_j \) to get the desired activation function. So if the in-star of \( j \) in \( \Delta \) contains no parallel arcs, we can construct a suitable component activation polynomial \( f_j \) with the same in-star.

It is worth pointing out that when iterating lemma 2.4.4 as described previously, we get a nice function. If the in-star of \( j \) in \( \Delta \) is \( \{(i, j, +), (i, j, -)\} \) or \( \{(i, j, +), (i, j, -), (k, j, \sigma_{kj})\} \). So let us begin with ternary component activation polynomials with these as in-stars.

**Lemma 2.4.5.** Let \( f_j : \{0, \pm 1\}^C \rightarrow \{\pm 1\} \) be the component activation polynomial \( f_j(x) = 2x_i^2 - 1 \). Then the in-star of \( f_j \) is \( \{(i, j, +), (i, j, -)\} \).

**Proof.** Observe that

\[
\partial^{i \rightarrow 0} f_j(-1) = \text{sgn}[f_j(0) - f_j(-1)][0 - (-1)] = \text{sgn}[-1 - 1] = -.
\]

Also

\[
\partial^{i \rightarrow 1} f_j(0) = \text{sgn}[f_j(1) - f_j(0)][1 - 0] = \text{sgn}[1 - (-1)] = +.
\]

Therefore the arcs \((i, j, +)\) and \((i, j, -)\) are contained in the in-star of \( j \) in \( I_f \). Since \( f_j \) depends only on \( x_i \), there are no other arcs in the in-star of \( j \).

If we apply lemma 2.4.4 to the polynomial in lemma 2.4.5, we can construct a component activation polynomial \( f_j \) with \( \{(i, j, +), (i, j, -), (k, j, \sigma_{kj})\} \) as its in-star where \( S_i = \{0, \pm 1\} \). These are the only two cases where we specifically need ternary polynomials.

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Next we will show how to produce a component activation polynomial whose in-star contains a single pair of parallel arcs and has in-degree four or more. We start with a polynomial \( f_j \) whose in-star contains exactly four arcs. Its inputs are all \( \pm 1 \) to satisfy the hypotheses of theorem 2.4.2. This polynomial depends only on three inputs \( x_{i_1}, x_{i_2} \) and \( x_{i_3} \) so we will write \( f_j(x) = f_j(x_{i_1}, x_{i_2}, x_{i_3}) \) for the following lemma.

**Lemma 2.4.6.** For \( \sigma_1, \sigma_2 \in \{\pm\} \), let \( f_j : \{0, \pm 1\}^{C_f} \to \{\pm 1\} \) be the component activation polynomial

\[
f_j(x) = \frac{1}{2}(x_{i_3}(\sigma_1 x_{i_1} + 1)(x_{i_3} + 1) + \sigma_2 x_{i_2} (\sigma_2 x_{i_2} + 1)(x_{i_3} - 1)) - 1
= \frac{1}{2}[(x_{i_1}^2 + \sigma_1 x_{i_1} - 1)(x_{i_3}^2 + x_{i_3}) + (x_{i_2}^2 + \sigma_2 x_{i_2} - 1)(x_{i_3}^2 - x_{i_3})] + (x_{i_3}^2 - 1)
= \begin{cases} 
1 & \sigma_1 x_{i_1} = x_{i_3} = 1 \text{ or } \sigma_2 x_{i_2} = -x_{i_3} = 1 \\
-1 & \text{otherwise.}
\end{cases}
\]

Then:

1. the in-star of \( f_j \) is \( \{(i_1, j, \sigma_1), (i_2, j, \sigma_2), (i_3, j, +), (i_3, j, -)\} \),

2. \( \partial^{i_3 \rightarrow 1} f_j(x_{i_1}, x_{i_2}, -1) = + \text{ if } \sigma_1 x_{i_1} = 1 \text{ and } \sigma_2 x_{i_2} = -1 \),

3. \( \partial^{i_3 \rightarrow 1} f_j(x_{i_1}, x_{i_2}, -1) = - \),

4. \( \partial^{i_1 \rightarrow 1} f_j(-1, x_{i_2}, 1) = \sigma_1 \),

5. \( \partial^{i_2 \rightarrow 1} f_j(x_{i_1}, -1, -1) = \sigma_2 \).

**Proof.** There are three parts to the proof. First we will show that the in-star of \( f_j \) contains the arcs \( (i_3, j, +) \) and \( (i_3, j, -) \). Then we will show that \( (i_1, j, \sigma_1) \) is the only arc from \( i_1 \) to \( j \) in the in-star of \( f_j \) and that \( (i_2, j, \sigma_2) \) is the only arc from \( i_2 \) to \( j \) in the in-star of \( f_j \).

First some useful observations about \( f_j \). Notice that

\[
f_j(x_{i_1}, x_{i_2}, 1) = \frac{1}{2}[(x_{i_1}^2 + \sigma_1 x_{i_1} - 1)(1^2 + 1) + (x_{i_2}^2 + \sigma_2 x_{i_2} - 1)(1^2 - 1)] + (1^2 - 1)
= x_{i_1}^2 + \sigma_1 x_{i_1} - 1.
\]

(2.9)

If in addition \( x_{i_1} = \pm 1 \), then

\[
f_j(x_{i_1}, x_{i_2}, 1) = \sigma_1 x_{i_1}.
\]

(2.10)

Similarly

\[
f_j(x_{i_1}, x_{i_2}, -1) = x_{i_2}^2 + \sigma_2 x_{i_2} - 1.
\]

(2.11)
If in addition \( x_{i_2} = \pm 1 \), then
\[
 f_j(x_{i_1}, x_{i_2}, -1) = \sigma_2 x_{i_2}. \tag{2.12}
\]
So if \( x_{i_1} = \pm 1 \) and \( x_{i_2} = \pm 1 \), by equations (2.10) and (2.12),
\[
\partial^{i_3 \rightarrow 1} f_j(x_{i_1}, x_{i_2}, -1) = \text{sgn}[f_j(x_{i_1}, x_{i_2}, 1) - f_j(x_{i_1}, x_{i_2}, -1)]-[1 - (-1)]
\]
\[
= \text{sgn}[f_j(x_{i_1}, x_{i_2}, 1) - f_j(x_{i_1}, x_{i_2}, -1)]
\]
\[
= \text{sgn}[\sigma_1 x_{i_1} - \sigma_2 x_{i_2}]. \tag{2.13}
\]
Now suppose that \( \sigma_1 x_{i_1} = 1 \) and \( \sigma_2 x_{i_2} = -1 \). In particular, this means that \( x_{i_1} = \pm 1 \) and \( x_{i_2} = \pm 1 \). So by equation (2.13),
\[
\partial^{i_3 \rightarrow 1} f_j(x_{i_1}, x_{i_2}, -1) = \text{sgn}[\sigma_1 x_{i_1} - \sigma_2 x_{i_2}] = \text{sgn}[1 - (-1)] = +.
\]
Similarly if \( \sigma_1 x_{i_1} = -1 \) and \( \sigma_2 x_{i_2} = 1 \), then
\[
\partial^{i_3 \rightarrow 1} f_j(x_{i_1}, x_{i_2}, -1) = \text{sgn}[-1 - 1] = -.
\]
Therefore \((i_3, j, +)\) and \((i_3, j, -)\) are both contained in the in-star of \( f_j \).

Next we will show that \((i_1, j, \sigma_1)\) is the only arc from \( i_1 \) to \( j \) in the in-star of \( f_j \). Consider \( \partial^{i_1 \rightarrow d} f_j(a, b, c) \). Using lemma 2.2.1 we can assume that \( d > a \), so \( \partial^{i_1 \rightarrow d} f_j(a, b, c) = \text{sgn}[f_j(d, b, c) - f_j(a, b, c)] [d - a] = \text{sgn}[f_j(d, b, c) - f_j(a, b, c)] \). Notice that \( f_j(d, b, -1) = f_j(a, b, -1) \) by equation (2.11). Also \( f_j(d, b, 0) = f_j(a, b, 0) = -1 \) from the definition of \( f_j \), so \( \partial^{i_1 \rightarrow d} f_j(a, b, c) = 0 \) if \( c \neq 1 \). So we need only consider the case when \( c = 1 \). Using equation (2.9),
\[
\partial^{i_1 \rightarrow d} f_j(a, b, 1) = \text{sgn}[f_j(d, b, 1) - f_j(a, b, 1)]
\]
\[
= \text{sgn}[d^2 + \sigma_1 d - 1 - (a^2 + \sigma_1 a - 1)]
\]
\[
= \text{sgn}[(d^2 - a^2) + \sigma_1 (d - a)]
\]
\[
= \text{sgn}[(d + a)(d - a) + \sigma_1 (d - a)]
\]
\[
= \text{sgn}[(d - a)][d + a + \sigma_1 (1)].
\]
Since \( d > a \),
\[
\partial^{i_1 \rightarrow d} f_j(a, b, 1) = \text{sgn}[(d + a) + \sigma_1 (1)].
\]
Now \( a, d \in \{0 \pm 1\} \), so \(-1 \leq (d + a) \leq 1 \). Then \((d + a) + 1 \geq 0\) and \((d + a) - 1 \leq 0\).
If \( \sigma_1 = + \), then \( \partial^{i_1 \rightarrow d} f_j(a, b, 1) = \text{sgn}[(d + a) + 1] \geq 0 \). If \( \sigma_1 = - \), then \( \partial^{i_1 \rightarrow d} f_j(a, b, 1) = \text{sgn}[(d + a) - 1] \leq 0 \). In either case \( \partial^{i_1 \rightarrow d} f_j(a, b, 1) \) is \( \sigma_1 \) or 0. Specifically,

\[
\partial^{i_1 \rightarrow 1} f_j(-1, b, 1) = \text{sgn}[(1 + (-1)) + \sigma_1(1)] = \sigma_1.
\]

Therefore \((i_1, j, \sigma_1)\) is the only arc from \(i_1\) to \(j\) in the in-star of \(f_j\).

Using a similar argument, we will show that \((i_2, j, \sigma_2)\) is the only arc from \(i_2\) to \(j\) in the in-star of \(f_j\).

Consider \(\partial^{i_2 \rightarrow d} f_j(a, b, c)\). Using lemma 2.2.1 again we can assume that \(d > b\). It follows that \(\partial^{i_2 \rightarrow d} f_j(a, b, c) = \text{sgn}[f_j(a, d, c) - f_j(a, b, c)]\). Similarly to before, if \(c \neq -1\), then \(f_j(a, d, c) = f_j(a, b, c)\), so \(\partial^{i_2 \rightarrow d} f_j(a, b, c) = 0\). By equation (2.11),

\[
\partial^{i_2 \rightarrow d} f_j(a, b, -1) = \text{sgn}[d^2 + \sigma_2 d - 1 - (b + \sigma_2 b - 1)]
\]

\[
= \text{sgn}[(d - b)][d + b + \sigma_2(1)]
\]

\[
= \text{sgn}[(d + b) + \sigma_2(1)].
\]

Again \(-1 \leq (d + b) \leq 1\) since \(b, d \in \{0, \pm 1\}\). So \((d + b) + 1 \geq 0\) and \((d + b) - 1 \leq 0\).

If \(\sigma_2 = +\), then \(\partial^{i_2 \rightarrow d} f_j(a, b, -1) = \text{sgn}[(d + b) + 1] \geq 0\). If \(\sigma_2 = -\), then \(\text{sgn}[(d + b) - 1] \leq 0\).

Either way \(\partial^{i_2 \rightarrow d} f_j(a, b, -1)\) is \(\sigma_2\) or 0. Specifically,

\[
\partial^{i_2 \rightarrow 1} f_j(x_{i_1}, -1, -1) = \text{sgn}[(1 + (-1)) + \sigma_2(1)] = \sigma_2.
\]

Therefore \((i_2, j, \sigma_2)\) is the only arc from \(i_2\) to \(j\) in the in-star of \(f_j\). \(\quad \square\)

Starting with the polynomial in lemma 2.4.6 we can then iteratively apply lemma 2.4.4 to construct a component activation polynomial \(f_j\) whose in-star contains \(\{(i_1, j, \sigma_1), (i_2, j, \sigma_2), (i_3, j, +), (i_3, j, -)\}\) and as many additional non-parallel arcs as we wish and the inputs can be taken exclusively from \(\{\pm 1\}\). This covers the case when the in-star of \(j\) in \(\Delta\) contains a single pair of parallel arcs.

Now suppose there are multiple pairs of parallel arcs into \(j\) in \(\Delta\).

First we start with a polynomial \(f_j\) whose in-star contains only two pairs of parallel arcs and whose inputs can be exclusively from \(\{\pm 1\}\). The polynomial depends only on inputs \(i\) and \(k\) so we will write \(f_j(x) = f_j(x_i, x_k)\) for the purposes of the lemma. Note that \(\lor\) means exclusive-or.

**Lemma 2.4.7.** Let \(f_j: \{0, \pm 1\}^{C_j} \rightarrow \{\pm 1\}\) be a component activation function such that

\[
f_j(x) = (1 + x_i - x_i^2)(1 + x_k - x_k^2) = \begin{cases} -1 & \text{if } (x_i = -1) \lor (x_k = -1) \\ 1 & \text{otherwise.} \end{cases}
\]

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Then:

1. the in-star of $f_j$ is \{(i, j, +), (i, j, -), (k, j, +), (k, j, -)\}.

2. $\partial^{i\mapsto 1} f_j(-1, 1) = \partial^{k\mapsto 1} f_j(1, -1) = +$.

3. $\partial^{i\mapsto 1} f_j(-1, -1) = \partial^{k\mapsto 1} f_j(-1, -1) = -$.

**Proof.** Since $f_j$ only depends on $x_i$ and $x_k$, all we need to do is calculate $\partial^{i\mapsto 1} f_j(-1, 1)$, $\partial^{k\mapsto 1} f_j(1, -1)$, $\partial^{i\mapsto 1} f_j(-1, -1)$, and $\partial^{k\mapsto 1} f_j(-1, -1)$.

\[
\begin{align*}
\partial^{i\mapsto 1} f_j(-1, 1) &= \text{sgn}[f_j(1, 1) - f_j(-1, 1)][1 - (-1)] = \text{sgn}[1 - (-1)] = +. \\
\partial^{k\mapsto 1} f_j(1, -1) &= \text{sgn}[f_j(1, 1) - f_j(1, -1)][1 - (-1)] = \text{sgn}[1 - (-1)] = +. \\
\partial^{i\mapsto 1} f_j(-1, -1) &= \text{sgn}[f_j(1, -1) - f_j(-1, -1)][1 - (-1)] = \text{sgn}[-1 - 1] = -. \\
\partial^{k\mapsto 1} f_j(-1, -1) &= \text{sgn}[f_j(-1, 1) - f_j(-1, -1)][1 - (-1)] = \text{sgn}[-1 - 1] = -. 
\end{align*}
\]

This shows that the in-star of $f_j$ is \{(i, j, +), (i, j, -), (k, j, +), (k, j, -)\}. \hfill \square

Combining lemma \[\text{Lemma 2.4.7}\] with the next, we will be able to construct a component activation polynomial $f_j$ whose in-star contains more than two pairs of parallel arcs.

**Lemma 2.4.8.** Let $g_j : \{0, \pm 1\}^{C_g} \to \{\pm 1\}$ be a non-constant component activation function that does not depend on input $i \in C_g$. Let $f_j$ be the component activation function such that $C_f = C_g$ and

\[
f_j(x) = (1 + x_i - x_i^2)g_j(x) = \begin{cases} 
-1 & \text{if } (g_j(x) = -1) \lor (x_i = -1) \\
1 & \text{otherwise.}
\end{cases}
\]

Then the in-star of $f_j$ contains $T \cup \{(i, j, +), (i, j, -)\}$ where $T$ is the in-star of $g_j$ and there are states $x, y \in S_f$ such that $x_i = y_i = -1$, $\partial^{i\mapsto 1} f_j(x) = +$ and $\partial^{i\mapsto 1} f_j(y) = -$. Also, if $k \in C_f$ such that $k \neq i$ and there is no arc from $k$ to $j$ in $T$, then there are no arcs from $k$ to $j$ in the in-star of $f_j$.

**Proof.** To prove this, we have to show that the in-star of $f_j$ contains $(i, j, +), (i, j, -)$, and all the arcs in $T$.

First let us show that every arc contained in $T$, the in-star of $g_j$, is also contained in the in-star of $f_j$. If $(k, j, \sigma_{kj}) \in T$, then there is a state $x$ such that $\partial^{k\mapsto a} g_j(x) = \sigma_{kj}$. Since $g_j$ does not depend on input $i$ we may assume that $x_i = 1$. Notice that if $x_i = 1$, then $f_j(x) = (1 + 1 - 1)g_j(x) = g_j(x)$. Therefore $\partial^{k\mapsto a} f_j(x) = \partial^{k\mapsto a} g_j(x)$. So the in-star of $f_j$ also contains $(k, j, \sigma_{kj})$. 

\[31\]
Now let us show that the in-star of $f_j$ contains $(i, j, +)$ and $(i, j, -)$. There is a state $x$ such that $g_j(x) = 1$ since $g_j$ is not constant. Since $g_j$ does not depend on input $i$, $g_j(x^{i\rightarrow 1}) = g_j(x)$ and we may assume that $x_i = -1$. Putting this into the formula for $f_j$,

$$
\partial^{i\rightarrow 1} f_j(x) = \text{sgn}[f_j(x^{i\rightarrow 1}) - f_j(x)][1 - (-1)]
= \text{sgn}[f_j(x^{i\rightarrow 1}) - f_j(x)]
= \text{sgn}[(1 + 1 - 1)g_j(x^{i\rightarrow 1}) - (1 - 1 - 1)g_j(x)]
= \text{sgn}[g_j(x) + g_j(x)]
= \text{sgn}[1 + 1]
= +.
$$

By a similar argument there is a state $y$ such that $g_j(y) = -1$ since $g_j$ is not constant and we may assume that $y_i = -1$. In this case

$$
\partial^{i\rightarrow 1} f_j(y) = \text{sgn}[f_j(y^{i\rightarrow 1}) - f_j(y)]
= \text{sgn}[(1 + 1 - 1)g_j(y^{i\rightarrow 1}) - (1 - 1 - 1)g_j(y)]
= \text{sgn}[g_j(y) + g_j(y)]
= \text{sgn}[-1 - 1]
= -.
$$

Therefore the in-star of $f_j$ contains $(i, j, +)$ and $(i, j, -)$ also.

Finally, if $k \in C_f$, and there is no arc from $k$ to $j$ in $T$, then $g(x^{k\rightarrow a}) = g(x)$ always since $g$ does not depend on $k$. Therefore $f(x^{k\rightarrow a}) = f(x)$, if $k \neq i$. \qed

Notice that if $g_j$ in lemma 2.4.8 is a polynomial, then $f_j$ is also a polynomial.

Now we can apply lemma 2.4.8 iteratively to the polynomial in lemma 2.4.7 to give us a component activation polynomial $f_j$ whose in-star contains two or more pairs of parallel arcs.

Now while applying lemma 2.4.8 there is the possibility to add an additional arc from $k$ to $j$ in the in-star of $f_j$ if there is already an arc from $k$ to $j$ in the in-star of $g_j$, $T$. But since the polynomial in 2.4.7 starts with pairs of parallel arcs, this is not an issue for our purposes since all possible arcs that could be contained in the in-star of $j$ are present at each step when we iteratively apply lemma 2.4.8 to the polynomial in lemma 2.4.7.

As we have done before, we can then iteratively apply lemma 2.4.4 to construct a component activation polynomial whose in-star contains as many pairs of parallel arcs as we like and as many

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non-parallel arcs as we like. This covers the case when the in-star of \( j \) in \( \Delta \) contains multiple parallel arc pairs. And at last we have covered all possible cases for the proof of the theorem. So we can indeed construct a component activation polynomial with any desired in-star.

### 2.4.1 Restricting to Boolean Activation Functions

It is worth noting that if we restrict to Boolean activation functions, then the polynomials used in the proof of theorem 2.4.2 simplify because if \( x \in \{\pm 1\} \), then \( x^2 = 1 \). The function in lemma 2.4.4 becomes

\[
f_j(x) = \frac{1}{2}(g_j(x) + 1)(\sigma_{ij}x_i + 1) - 1 = \min(g_j(x), \sigma_{ij}x_i).
\]

So if we want a component activation function \( f_j \) whose in-star is \( \{(i, j, \sigma_{ij}) \mid i \in A\} \), then \( f_j(x) = \min_{i \in A}(\sigma_{ij}x_i) \) works in the Boolean case. The function in lemma 2.4.6 becomes

\[
f_j(x) = \frac{1}{2}[\sigma_1 x_{i_1}(1 + x_{i_3}) + \sigma_2 x_{i_2}(1 - x_{i_3})] = \begin{cases} \sigma_1 x_{i_1} & x_{i_3} = 1 \\ \sigma_2 x_{i_2} & x_{i_3} = -1. \end{cases}
\]

The function in lemma 2.4.7 becomes \( f_j(x) = x_{i_1}x_{i_2} \) and the function in lemma 2.4.8 becomes \( f_j(x) = x_ij(x) \). So when lemma 2.4.8 is iteratively applied to the function in lemma 2.4.7, we just get a product. Specifically, if we want a component activation polynomial \( f_j \) whose in-star is \( \{(i, j, +), (i, j, -) \mid i \in A\} \) where \( A \) contains at least two vertices, then \( f_j(x) = \prod_{i \in A} x_i \) will work in the Boolean case.

### 2.5 Forbidden Local Interaction Graphs

In this section we will address the same question as in the previous section for local interaction graphs. That is, for a given signed digraph \( \Delta \), when can we find an activation function \( f \) and states \( x, y \in S_f \) such that \( I_f(x, y) = \Delta \)? We will only discuss the first kind of local interaction graphs, but by proposition 2.2.3 all the results in this section apply to restricted local interaction graphs too.

The first observation is that for this to be possible, \( \Delta \) cannot contain any parallel arcs since there is a single arc from \( i \) to \( j \) in \( I_f(x, y) \) if and only if \( \partial^{i \rightarrow j} f_j(x) \) is non-zero. It turns out that if this condition is met, then \( \Delta \) is a local interaction graph. Moreover \( f \) is at worst ternary.

For a set of vertices \( W \subseteq V(\Delta) \), we can also give a condition for when it is possible to have \( |S_i| \leq 2 \) for all \( i \in W \) for this activation function. For a vertex \( j \) in a signed digraph, if every arc into \( j \) is positive, then call \( j \) positive inward-homogeneous. Similarly, if every arc into \( j \) is
negative, call \( j \) negative inward-homogeneous. Call \( j \) inward-homogeneous if it is either positive, or negative inward-homogeneous. For \( W \subseteq V \), call \( W \) inward-homogeneous if every vertex in \( W \) is inward-homogeneous. We call a signed digraph inward-homogeneous if all of its vertices are inward-homogeneous.

**Theorem 2.5.1.** Given a signed digraph \( \Delta \) containing no parallel arcs.

1. There are an activation function \( f \) and states \( x, y \in S_f \) such that \( I_f(x, y) = \Delta \).

2. For each \( j \in C_f \), \( S_j \) can be chosen to be \( \{\pm 1\} \) or \( \{0, \pm 1\} \).

3. For \( W \subseteq C_f \), there is an activation function \( f \) such that \( |S_j| \leq 2 \) for all \( j \in W \) if and only if there is \( \zeta \in \{\pm\}^{V(\Delta)} \) such that \( W \) is inward-homogeneous in \( \zeta \Delta \).

This theorem has an immediate corollary for Boolean activation functions.

**Corollary 2.5.2.** A signed digraph \( \Delta \) is a local interaction graph of a Boolean activation function if and only if \( \Delta \) contains no parallel arcs and \( \Delta \) is switching equivalent to an inward-homogeneous digraph.

**Proof.** This is simply part three of theorem 2.5.1 when \( W = C_f \)

A natural question arises from this corollary. Given a signed digraph, is there a way to quickly identify whether it is switching equivalent to an inward-homogeneous digraph or not? As of this writing, I still have not found a satisfying answer to this question.

**Proof of theorem 2.5.1** Let us begin with the first part of the theorem, that for any signed digraph we can in fact find an appropriate activation function.

We will start with a method of constructing an activation function \( f \) such that \( I_f(x, y) = \Delta \) in a simpler case that also addresses the size of \( S_i \) for each \( i \in C_f \).

**Lemma 2.5.3.** Let \( \Delta \) be a signed digraph that does not contain any parallel arcs. Let \( W \subseteq V(\Delta) \) be inward-homogeneous. Then there are an activation function \( f \) and \( x, y \in S_f \) such that \( I_f(x, y) = \Delta \) and \( S_i = \{\pm 1\} \) if \( i \in W \) and \( S_i = \{0, \pm 1\} \) otherwise.

**Proof.** We have to find an activation function \( f \) such that \( \partial^{i \rightarrow y_i} f_j(x) = \sigma_{ij} \) if \( (i, j, \sigma_{ij}) \in \Delta \) and \( \partial^{i \rightarrow y_i} f_j(x) = 0 \) otherwise. To completely determine \( I_f(x, y) \), it will suffice to know the values of \( f_j(x) \) and \( f_j(x^{i \rightarrow y_i}) \) for all \( i, j \in C_f \).
Given a signed digraph $\Delta$ that contains no parallel arcs, $\zeta \in \{\pm\}^{V(\Delta)}$ and $W \subseteq V(\Delta)$ that is inward-homogeneous in $\zeta \Delta$. Then there are an activation function $f$ and $x, y \in S_f$ such that $\Delta = I_f(x, y)$, $S_i = \{\pm 1\}$ if $i \in W$ and $S_i = \{0, \pm 1\}$ otherwise.

**Proof.** By lemma [2.5.3](#), there are an activation function $g$ and $x, y \in S_g$ such that $I_g(x, y) = \zeta \Delta$ and $S_i = \{\pm 1\}$ if $i \in W$ and $S_i = \{0, \pm 1\}$ otherwise. Using theorem [2.3.1](#),

$$\Delta = \zeta(\zeta \Delta) = \zeta I_g(x, y) = I_{g^\zeta}(\zeta x, \zeta y).$$

Now let $f = g^\zeta$. Together with the states $\zeta x$ and $\zeta y$, these states and $f$ satisfy all hypotheses of the lemma. \qed

So now we know that we can construct an activation function $f$ such that $S_j = \{\pm 1\}$ for all $j$ in $W \subseteq C_f$ if there is $\zeta \in \{\pm\}^{C_f}$ such that $W$ is inward-homogeneous in $\zeta \Delta$. 

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For the last part of the theorem, suppose that there are an activation function \(f\) and \(x, y \in \mathbb{S}_f\) such that \(I_f(x, y) = \Delta\) and \(|S_j| \leq 2\) for all \(j \in W\) for some \(W \subseteq C_f\). We will construct \(\zeta \in \{\pm\}^C_f\) such that \(W\) is inward-homogeneous in \(\zeta \Delta\).

First a useful lemma.

**Lemma 2.5.5.** Given an activation function \(f\) and states \(x, y \in \mathbb{S}_f\) such that \(x \leq y\).

1. If \(f_j(x) \geq f_j(x^{i \rightarrow y_i})\) for all \(i \in C_f\), then \(j\) is negative inward-homogeneous in \(I_f(x, y)\).

2. If \(f_j(x) \leq f_j(x^{i \rightarrow y_i})\) for all \(i \in C_f\), then \(j\) is positive inward-homogeneous in \(I_f(x, y)\).

**Proof.** If \(y_i = x_i\), then \(\partial^{i \rightarrow y_i} f_j(x) = 0\), so suppose that \(y_i > x_i\). This means that \(\partial^{i \rightarrow y_i} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow y_i}) - f_j(x)]\). So if \(f_j(x) \geq f_j(x^{i \rightarrow y_i})\), then \(\partial^{i \rightarrow y_i} f_j(x)\) is negative or zero. Therefore every arc into \(j\) is negative.

Similarly if \(f_j(x) \leq f_j(x^{i \rightarrow y_i})\), then \(\partial^{i \rightarrow y_i} f_j(x)\) is positive or zero. So every arc into \(j\) is positive in that case.

The property \(f_j(x) \geq f_j(x^{i \rightarrow y_i})\) for all \(i \in C_f\), or \(f_j(x) \leq f_j(x^{i \rightarrow y_i})\) for all \(i \in C_f\) is automatically satisfied if \(|S_j| \leq 2\). So the next lemma follows immediately from lemma 2.5.5.

**Lemma 2.5.6.** Given an activation function \(f\) and states \(x, y \in \mathbb{S}_f\) such that \(x \leq y\). If \(|S_j| \leq 2\), then \(j\) is inward-homogeneous in \(I_f(x, y)\).

From here we are ready to show that if \(|S_j| \leq 2\) for all \(j \in W \subseteq C_f\), then there is \(\zeta \in \{\pm\}^C_f\) such that \(j\) is inward-homogeneous in \(\zeta I_f(x, y) = I_{f \zeta}(\zeta x, \zeta y)\).

**Lemma 2.5.7.** Given an activation function \(f\) and states \(x, y \in \mathbb{S}_f\). Let \(W \subseteq C_f\) be the set of vertices with the property that \(|S_j| \leq 2\) if \(j \in W\). Let \(\zeta \in \{\pm\}^C_f\) where \(\zeta_i = -1\) if and only if \(x_i > y_i\). Then \(W\) is inward-homogeneous in \(\zeta I_f(x, y)\).

**Proof.** From the definition of \(\zeta\), \(\zeta_i x_i < \zeta_i y_i\) if \(\zeta_i = -1\) and \(\zeta_i x_i \leq \zeta_i y_i\) if \(\zeta_i = +1\). Therefore \(\zeta x \leq \zeta y\). Since \(|S_j| \leq 2\), for any \(j \in W\) either \(f_j^\zeta(\zeta x) \leq f_j^\zeta(\zeta x^{i \rightarrow y_i})\) or \(f_j^\zeta(\zeta x) \geq f_j^\zeta(\zeta x^{i \rightarrow y_i})\) for all \(i \in C_f\). Therefore \(j\) is inward-homogeneous in \(I_{f \zeta}(\zeta x, \zeta y)\). And \(I_{f \zeta}(\zeta x, \zeta y) = \zeta I_f(x, y)\) by lemma 2.5.6.

So now we can construct \(\zeta \in \{\pm\}^V(\Delta)\) as required and so conclude the proof of theorem 2.4.2.
What if you have a signed digraph $\Delta$ that contains no parallel arcs, but no Boolean activation function $f$ exists such that $I_f(x,y) = \Delta$. We know now that there are non-Boolean activation functions $f$ such that $I_f(x,y) = \Delta$ by 2.5.1. Among all of these activation functions, which one is the “most Boolean”? More precisely, among all such activation functions, which one has the most components such that $|S_j| = 2$? By theorem 2.5.1 there is an activation function $f$ such that $|S_j| = 2$ for all $j \in W \subseteq V(\Delta)$ if and only if there is $\zeta \in \{\pm\}^{V(\Delta)}$ such that every vertex in $W$ is inward-homogeneous in $\zeta \Delta$. So the maximal number of Boolean components an activation function $f$ such that $I_f(x,y) = \Delta$ can have is the same as the maximal size of $W \subseteq V(\Delta)$ such that, for some $\zeta \in \{\pm\}^{V(\Delta)}$, every $j \in W$ is inward-homogeneous in $\zeta \Delta$.

2.6 Interaction Graphs of Threshold Activation Functions

Continuing the theme of what kinds of signed digraphs are interaction graphs, in this section we explore the question for threshold activation functions. This is a kind of activation function has been used to model real world gene regulatory networks. See for example [3].

A Boolean activation function $f : \{\pm 1\}^{C_f} \rightarrow \{\pm 1\}^{C_f}$ is called a threshold activation function if for each $j \in C_f$, there is an $a^j \in \mathbb{R}^{C_f}$ that defines the output of $f_j$ in the following way.

We say the $j$th component of $f$ tends active when $f_j(x) = 1$ if $a^j \cdot x \geq 0$ and $f_j(x) = -1$ if $a^j \cdot x < 0$. Similarly, we say the $j$th component of $f$ tends inactive when $f_j(x) = 1$ if $a^j \cdot x > 0$ and $f_j(x) = -1$ if $a^j \cdot x \leq 0$.

Interaction graphs of threshold activation functions have a convenient condition on their arc signs.

**Proposition 2.6.1.** Given a threshold activation function $f$. If the arc $(i,j,\sigma_{ij}) \in I_f$, then $\sigma_{ij} = \text{sgn}(a^j_i)$ where $a^j$ defines $f_j$.

**Proof.** Suppose the arc $(i,j,\sigma_{ij}) \in I_f$. Then $\sigma_{ij} = \partial^{\leftarrow} f_j(x)$ for some state $x \in S_f$. I claim that $\text{sgn}[a^j \cdot x^{\leftarrow} - a^j \cdot x] = \text{sgn}[f_j(x^{\leftarrow}) - f_j(x)]$ in this case. If this is true, our result will follow because

$$\partial^{\leftarrow} f_j(x) = \text{sgn}[f_j(x^{\leftarrow}) - f_j(x)][-x_i - x_i]$$

$$= \text{sgn}[a^j \cdot x^{\leftarrow} - a^j \cdot x][-2x_i]$$

$$= \text{sgn}[a^j \cdot (x^{\leftarrow} - x)][-2x_i]$$

$$= \text{sgn}[a^j_i (-2x_i)][-2x_i]$$

$$= \text{sgn}[a^j_i].$$
To see that $\text{sgn}[a^j \cdot x^i - a^i \cdot x] = \text{sgn}[f_j(x^i) - f_j(x)]$, observe that either $f_j(x^i) - f_j(x) < 0$, or $f_j(x^i) - f_j(x) > 0$ because $\partial^i f_j(x) \neq 0$.

First suppose that $f_j(x^i) - f_j(x) < 0$. So $f_j(x^i) = -1$ and $f_j(x) = 1$. If $f_j$ tends active, then $a^j \cdot x \geq 0$ and $a^j \cdot x^i < 0$. If $f_j$ tends inactive, then $a^j \cdot x > 0$ and $a^j \cdot x^i \leq 0$. Either way, $a^j \cdot x^i - a^j \cdot x < 0$. Therefore $\text{sgn}[a^j \cdot x^i - a^j \cdot x] = \text{sgn}[f_j(x^i) - f_j(x)]$.

Now suppose that $f_j(x^i) - f_j(x) > 0$. So $f_j(x^i) = 1$ and $f_j(x) = -1$. If $f_j$ tends active, then $a^j \cdot x < 0$ and $a^j \cdot x^i \geq 0$. If $f_j$ tends inactive, then $a^j \cdot x \leq 0$ and $a^j \cdot x^i > 0$. Either way, $a^j \cdot x^i - a^j \cdot x > 0$. Therefore $\text{sgn}[a^j \cdot x^i - a^j \cdot x] = \text{sgn}[f_j(x^i) - f_j(x)]$ again.

Since $\text{sgn}[a^j \cdot x^i - a^j \cdot x] = \text{sgn}[f_j(x^i) - f_j(x)]$, the result now follows from our earlier calculation. 

\section{Global Interaction Graphs}

Proposition \ref{prop:2.6.1} has an immediate corollary.

\begin{corollary}
If $f$ is a threshold activation function, then $\mathcal{I}_f$ contains no parallel arcs.
\end{corollary}

\begin{proof}
Suppose $\mathcal{I}_f$ contains $(i, j, +)$ and $(i, j, -)$. Since $(i, j, +) \in \mathcal{I}_f$, $\text{sgn}(a^j_i) = +$ by proposition \ref{prop:2.6.1}. Similarly, since $(i, j, -) \in \mathcal{I}_f$, $\text{sgn}(a^j_i) = -$, a contradiction. Therefore $(i, j, +)$ and $(i, j, -)$ cannot both be contained in $\mathcal{I}_f$.
\end{proof}

It turns out that this necessary condition of containing no parallel arcs is also sufficient for a signed digraph to be the interaction graph of a threshold activation function.

\begin{theorem}
A signed digraph is the interaction graph of a threshold activation function $f$ if and only if it contains no parallel arcs.
\end{theorem}

\begin{proof}
For each $j \in \Delta$, we will construct an $a^j$ that defines the component threshold activation function $f_j$ such that the in-stars of $j$ in both $\Delta$ and $\mathcal{I}_f$ are the same.

For $j \in C_f$, let $P$ be the set of all vertices in $\Delta$ such that there is a positive arc from each vertex in $P$ to $j$. Similarly, let $N$ be the set of all vertices in $\Delta$ such that there is a negative arc from each vertex in $N$ to $j$. There are three cases: the in-degree of $j$ is odd, the in-degree of $j$ is even and greater than two, or the in-degree of $j$ is two. We will show how to construct $a^j$ in each of these cases and show that the in-star of $j$ in $\mathcal{I}_f$ contains every arc in the in-star of $j$ in $\Delta$.

First suppose the in-degree of $j$ is odd. Let $a^j_i = 1$ if $i \in P$, $a^j_i = -1$ if $i \in N$ and $a^j_i = 0$ otherwise. This way if $(i, j, \sigma_{ij}) \in \Delta$, then $\text{sgn}(a^j_i) = \sigma_{ij}$. Let $k \in P \cup N$. Now since $|\{P \cup N\}|$
is odd, $|(P \cup N) \setminus k|$ is even. Since $a_i^j$ is 1 or $-1$ for all $i \in P \cup N$, as are the values of $x_i$ for any state, there is a state $x \in \mathbb{S}_f$ such that

$$\sum_{i \in (P \cup N) \setminus k} a_i^j x_i = 0.$$ 

You can just alternately add and subtract 1 for example. This works since $|(P \cup N) \setminus k|$ is even. Since $a_i^j = 0$ for all $i$ outside of $P \cup N$, $a_i^j \cdot x = a_i^j x_k$ and $a_i^j \cdot x^\leftarrow = -a_i^j x_k$ and these are non-zero. Therefore $\partial^k \to f_j(x) = \text{sgn}(a_k^j) = \sigma_{ij}$ by proposition $2.6.1$. Therefore $(i, j, \sigma_{ij}) \in \mathcal{I}_f$ if

$$(i, j, \sigma_{ij}) \in \Delta.$$ 

Now suppose that the in-degree of $j$ is even and more than 2. We choose $m \in P \cup N$ and let $a_m^j = 2$ if $m \in P$, or $a_m^j = -2$ if $m \in N$. For all other $i \neq m$, let $a_i^j = 1$ if $i \in P$, $a_i^j = -1$ if $i \in N$ and $a_i^j = 0$ otherwise. Again this means that if $(i, j, \sigma_{ij}) \in \Delta$, then $\text{sgn}(a_i^j) = \sigma_{ij}$. Now let $k \in (P \cup N)$. If $k \neq m$, then similarly to the odd case, there is a state $x \in \mathbb{S}_f$ such that

$$\sum_{i \in (P \cup N) \setminus k} a_i^j x_i = 0.$$ 

For example, we can first let $a_m^j x_m = 2$ then subtract 1 twice. This is possible since $|(P \cup N) \setminus k| > 2$ because $P \cup N$ is even but greater than two. Doing this accounts for three elements of $(P \cup N) \setminus k$, leaving us with an even number of elements. Now we can just alternately add and subtract 1 as we did previously. Therefore $a_i^j \cdot x = a_k^j x_k$ and $a_i^j \cdot x^\leftarrow = -a_k^j x_k$, as before and these are non-zero. So $\partial^m \to f_j(x) = \text{sgn}(a_m^j)$ by proposition $2.6.1$. If $k = m$, then there is a state $x \in \mathbb{S}_f$ such that

$$\sum_{i \in (P \cup N) \setminus m} a_i^j x_i = 1.$$ 

Just alternately add and subtract 1 again, which will yield a final total of 1 since $|(P \cup N) \setminus m|$ is odd. Therefore $a_i^j \cdot x = 1 + a_m^j x_m$ and $a_i^j \cdot x^\leftarrow = 1 - a_m^j x_m$. Since $a_m^j = \pm 2$, one of these is 3 and the other is $-1$. Therefore $f_j(x) = -f_j(x^\leftarrow)$, so $\partial^m \to f_j(x) = \text{sgn}(a_m^j)$ by proposition $2.6.1$. Therefore $(i, j, \sigma_{ij}) \in \mathcal{I}_f$ if $(i, j, \sigma_{ij}) \in \Delta$.

If the in-degree of $j$ is 2, then let $(i_1, j, \sigma_1)$ and $(i_2, j, \sigma_2)$ be the arcs into $j$. Let $a_i^j \in \mathbb{R}^C_j$ where $a_{i_1}^j = 1$ if $\sigma_1 = +$, $a_{i_1}^j = -1$ if $\sigma_1 = -$, $a_{i_2}^j = 1$ if $\sigma_2 = +$, $a_{i_2}^j = -1$ if $\sigma_2 = -$, and all other $a_k^j = 0$. Now let $x \in \mathbb{S}_f$ such that $a_{i_1}^j x_{i_1} + a_{i_2}^j x_{i_2} = 1$. Then $a_i^j \cdot x = 2$ and $a_i^j \cdot x^\leftarrow = a_{i_1}^j x_{i_1} x_{i_2} = 0.$ So if $f_j$ is the tending inactive component threshold activation function defined by $a_i^j$, then $f_j(x) = 1$ and $f_j(x^\leftarrow) = f_j(x^\rightarrow) = -1$. Therefore $\partial^1 \to f_j(x) = \sigma_1$ and $\partial^2 \to f_j(x) = \sigma_2$ by proposition $2.6.1$. So $(i_1, j, \sigma_1)$ and $(i_2, j, \sigma_2)$ are arcs in $\mathcal{I}_f$. 

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Finally we have to show that every arc in $I_f$ is contained in $\Delta$. We will argue the contrapositive statement. If there is no arc from $i$ to $j$ in $\Delta$, then $a_{ij}^1 = 0$ by definition. By proposition 2.6.1 there is no arc from $i$ to $j$ in $I_f$. Therefore $(i, j, \sigma_{ij}) \in I_f$ if and only if $(i, j, \sigma) \in \Delta$. 

It is worth pointing out that in order to have the in-degree of a component $j$ be two, then $|a_{ij}^1| = |a_{ij}^2|$. To see this, just observe that if $|a_{ij}^1| > |a_{ij}^2|$, and $a_{ij}^1 x_{ii} + a_{ij}^2 x_{i2} > 0$ say, then $a_{ij}^1 x_{ii} - a_{ij}^2 x_{i2} > 0$ also. This means that $f_j(x) = f_j(x^{i2-})$, so $\partial^{i2-} f_j(x) = 0$ for all states. This is why we had to invoke that the threshold activation function is tending inactive only in the in-degree two case in the previous theorem whereas this was not required for any other in-degree.

### 2.6.2 Local Interaction Graphs

What restrictions are there on which signed digraphs are local interaction graphs of threshold activation functions? That is, if $\Delta$ is a signed digraph, then is there a threshold activation function $f$ and a state $x \in S_f$ such that $I_f(x) = \Delta$? It turns out that there are no restrictions on fewer than three components, save for having no parallel arcs and the condition of inward-homogeneity from theorem 2.5.1.

First consider the case where $|C_f| = 1$. Then $f(x) = x$ is a threshold activation function where $a = 1$. Also, $\partial^{1-} f(1) = +$. So $I_f(1)$ consists of a single positive loop. Similarly, $f(x) = -x$ is a threshold activation function where $a = -1$, where $\partial^{1-} f(1) = -$. So $I_f(1)$ consists of a single negative loop. And $f(x) = 0$ is a threshold activation function where $a = 0$, where $\partial^{1-} f(x) = 0$. So $I_f(1)$ contains no arcs. So every possible graph on a single vertex that does not contain parallel arcs is the local interaction graph of a threshold activation function.

Now consider the case where $|C_f| = 2$. Let us say $C_f = \{1, 2\}$. By theorem 2.5.1 there is $\zeta \in \{\pm\}^2$ such that $\zeta \Delta$ is inward-homogeneous. Let $(a, b)$ define $f_1$. By proposition 2.6.1 there is no corresponding arc in $I_f(x)$ into 1 for any state $x$ if $a$ or $b$ is zero. So this gives us a straightforward way to handle the absence of arcs in $\Delta$.

First suppose that the arcs into 1 in $\zeta \Delta$ are +. Then let $a = 1$ if there is an arc from 1 to itself in $\Delta$ and $a = 0$ otherwise. Similarly, let $b = 1$ if there is an arc from 2 to 1 in $\Delta$ and $b = 0$ otherwise. I claim that if $f_1$ is the tending inactive threshold activation function defined by $(a, b)$, then the in-star of 1 in $I_f(1, 1)$ is the same as that in $\zeta \Delta$. We already know what happens when either $a$ or $b$ is 0, no corresponding arc. If $a = 1$, then $(a, b) \cdot (1, 1) = a + b \geq 1$ and $(a, b) \cdot (-1, 1) = -a + b \leq 0$ since $b$ is 0 or 1. Therefore $f_1(1, 1) = 1$ and $f_1(-1, 1) = -1$ since $f_1$ tends inactive. Therefore $\partial^{1-} f_1(1, 1) = \text{sgn}[f_1(-1, 1) - f_1(1, 1)][-1 - 1] = \text{sgn}[1 - 1][2] = +$. Similarly, if $b = 1$ then
\( (a, b) \cdot (1, 1) = a + b \geq 1 \) and \( (a, b) \cdot (1, -1) = -a + b \leq 0 \) since \( a \) is 0 or 1. Therefore \( f_1(1, 1) = 1 \) and \( f_1(1, -1) = -1 \) again. So \( \partial^{2\rightarrow} f_1(1, 1) = \text{sgn}[f_1(1, -1) - f_1(1, 1)][-1 - 1] = + \) also.

Now suppose that the arcs into 1 in \( \zeta \Delta \) are \(-\). Let \( a = -1 \) if there is an arc from 1 to itself in \( \Delta \) and \( a = 0 \) otherwise. Similarly, let \( b = -1 \) if there is an arc from 2 to 1 in \( \Delta \) and \( b = 0 \) otherwise. I claim that if \( f_1 \) is the active tending threshold activation function defined by \( (a, b) \), then the in-star of 1 in \( \mathcal{I}_f(1, 1) \) is the same as that in \( \zeta \Delta \). Again, we already know what happens when either \( a \) or \( b \) is 0. If \( a = -1 \), then \( (a, b) \cdot (1, 1) = a + b \leq -1 \) and \( (a, b) \cdot (-1, 1) = -a + b \geq 0 \) since \( b \) is 0 or \(-1 \). Therefore \( f_1(1, 1) = -1 \) and \( f_1(-1, 1) = 1 \) since \( f_1 \) tends active. Therefore \( \partial^{1\rightarrow} f_1(1, 1) = \text{sgn}[f_1(-1, 1) - f_1(1, 1)][-1 - 1] = \text{sgn}[1 - (-1)][-2] = - \). Similarly, if \( b = -1 \), then \( (a, b) \cdot (1, 1) = a + b \leq -1 \) and \( (a, b) \cdot (1, -1) = -a + b \geq 0 \) since \( a \) is 0 or \(-1 \). Therefore \( f_1(1, 1) = -1 \) and \( f_1(-1, 1) = 1 \) again. So \( \partial^{2\rightarrow} f_1(1, 1) = \text{sgn}[f_1(1, -1) - f_1(1, 1)][-1 - 1] = - \) also.

Now \( f_2 \) can be defined in exactly the same way depending on the sign of the arcs into 2 in \( \zeta \Delta \). That is, if the arcs into 2 are positive in \( \zeta \Delta \), then let \( f_2 \) be the tending inactive component threshold activation function where \( a = 1 \) if there is an arc from 2 to 1 in \( \Delta \) and \( a = 0 \) otherwise, and \( b = 1 \) if there is an arc from 2 to itself in \( \Delta \) and \( b = 0 \) otherwise. If the sign of the arcs into 2 are negative in \( \zeta \Delta \), then let \( f_2 \) be the active tending component threshold activation function where \( a = -1 \) if there is an arc from 2 to 1 in \( \Delta \) and \( a = 0 \) otherwise, and \( b = -1 \) if there is an arc from 2 to itself in \( \Delta \) and \( b = 0 \) otherwise. Then all calculations for \( \partial^{i\rightarrow} f_2(1, 1) \) are the same as before.

So we have shown that \( \mathcal{I}_f(1, 1) = \zeta \Delta \). Therefore \( \Delta = \zeta \mathcal{I}_f(x) = \mathcal{I}_{f^\zeta}(\zeta x) \) by theorem \[2.3.1\] So we now use the following lemma to finish the argument, which we'll also need for the case when \( \Delta \) has more than two vertices.

**Lemma 2.6.4.** Let \( f \) be a threshold activation function and let \( \lambda = (\phi, \zeta) \) be a switching isomorphism of \( f \). Then \( f^\lambda \) is also a threshold activation function.

**Proof.** I claim that \( f^\lambda \) is the same function as the threshold activation function \( g \) where \( g_{\phi j} \) is defined by \( \zeta_j(\lambda a^j) \) where \( a^j \) defines \( f_j \). If \( \zeta_j = + \), then the tendency of \( g_{\phi j} \), active or inactive, is the same as the tendency of \( f_j \). If \( \zeta_j = - \), then the tendency of \( g_{\phi j} \) is the opposite of the tendency of \( f_j \). So we just have to show that \( f^\lambda = g \) which we will do by showing they always have the same output on each component.

Since we have to compare dot products in \( \mathbb{S}_f \) with dot products in \( \mathbb{S}_f^\lambda \), the following equation
will be useful. So for a state \( x \in S_f \) and \( a^j \in \mathbb{R}^{C_f} \),

\[
a^j \cdot x = \zeta a^j \cdot \zeta x = \phi(\zeta a^j) \cdot \phi(\zeta x) = \lambda a^j \cdot \lambda x
\]

(2.14)

First suppose that \( \zeta_j = + \). If \( f^\lambda_{\phi j}(\lambda x) = 1 \), then \( f_j(x) = 1 \) by equation (2.7). If \( f_j \) tends inactive, \( \lambda a^j \cdot \lambda x > 0 \) by equation (2.14). Since \( \zeta_j = + \), \( g_{\phi j} \) also tends inactive and \( \zeta_j(\lambda a^j) \cdot \lambda x > 0 \). Therefore \( g_{\phi j}(\lambda x) = 1 \) also. If \( f_j \) tends active, \( \lambda a^j \cdot \lambda x \geq 0 \) by equation (2.14). Since \( \zeta_j = + \), \( g_{\phi j} \) also tends active and \( \zeta_j(\lambda a^j) \cdot \lambda x \geq 0 \). Therefore \( g_{\phi j}(\lambda x) = 1 \) also. If \( f_j \) tends active, \( \lambda a^j \cdot \lambda x < 0 \) by equation (2.14). Since \( \zeta_j = + \), \( g_{\phi j} \) also tends active and \( \zeta_j(\lambda a^j) \cdot \lambda x < 0 \). Therefore \( g_{\phi j}(\lambda x) = -1 \) again.

Now suppose that \( \zeta_j = - \). If \( f^\lambda_{\phi j}(\lambda x) = 1 \), then \( f_j(x) = -1 \) by equation (2.7). If \( f_j \) tends inactive, \( \lambda a^j \cdot \lambda x \leq 0 \) by equation (2.14). Since \( \zeta_j = - \), \( g_{\phi j} \) tends active and \( \zeta_j(\lambda a^j) \cdot \lambda x \geq 0 \). Therefore \( g_{\phi j}(\lambda x) = 1 \) also. If \( f_j \) tends active, \( \lambda a^j \cdot \lambda x < 0 \) by equation (2.14). Since \( \zeta_j = - \), \( g_{\phi j} \) tends inactive and \( \zeta_j(\lambda a^j) \cdot \lambda x > 0 \). Therefore \( g_{\phi j}(\lambda x) = 1 \) again. If \( f^\lambda_{\phi j}(\lambda x) = -1 \), then \( f_j(x) = 1 \) by equation (2.7). If \( f_j \) tends inactive, \( \lambda a^j \cdot \lambda x > 0 \) by equation (2.14). Since \( \zeta_j = - \), \( g_{\phi j} \) tends active and \( \zeta_j(\lambda a^j) \cdot \lambda x < 0 \). Therefore \( g_{\phi j}(\lambda x) = -1 \) also. If \( f_j \) tends active, \( \lambda a^j \cdot \lambda x \geq 0 \) by equation (2.14). Since \( \zeta_j = - \), \( g_{\phi j} \) tends inactive and \( \zeta_j(\lambda a^j) \cdot \lambda x \leq 0 \). Therefore \( g_{\phi j}(\lambda x) = -1 \) again.

So we have shown that in all cases, \( g_{\phi j}(\lambda x) = f^\lambda_{\phi j}(\lambda x) \). Therefore \( f^\lambda \) is indeed a threshold activation function.

So every signed digraph on two components that contains no parallel arcs and is switching equivalent to an inward-homogeneous digraph is a local interaction graph of a threshold activation function. However, the situation changes when there are more than two components.

**Proposition 2.6.5.** Let \( f \) be a threshold activation function such that \( n = |C_f| > 2 \). Then the in-degree of any component in \( I_f(x) \) is at most \( n - 1 \).
Proof. First observe that
\[
\sum_{i \in C_f} a^j \cdot x^{i \rightarrow} = a^j \cdot \left( \sum_{i \in C_f} x^{i \rightarrow} \right) = a^j \cdot (|n - 2| x) = (n - 2) a^j \cdot x
\] (2.15)

Now suppose there is \( j \in C_f \) in \( I_f(x) \) whose in-degree is \( n \). This means that \( f_j(x) \neq f_j(x^{i \rightarrow}) \) for all \( i \in C_f \). If for all \( i \in C_f \), \( a^j \cdot x^{i \rightarrow} > 0 \), \( a^j \cdot x^{i \rightarrow} \geq 0 \), or \( a^j \cdot x^{i \rightarrow} \leq 0 \), then by equation (2.15), \( a^j \cdot x > 0 \), \( a^j \cdot x \geq 0 \), \( a^j \cdot x < 0 \), or \( a^j \cdot x \leq 0 \) respectively. But this means that \( f_j(x) = f_j(x^{i \rightarrow}) \) for all \( i \in C_f \), so there are no arcs into \( j \) in \( I_f(x) \), a contradiction. Therefore the in-degree of \( j \) cannot be \( n \).

It turns out that this is the only additional restriction on local interaction graphs of threshold activation functions.

Theorem 2.6.6. Given a signed digraph \( \Delta \).

1. If \( \Delta \) has 1 or 2 vertices, then it is the local interaction graph of a threshold activation function if and only if it contains no parallel arcs and is switching equivalent to an inward-homogeneous signed digraph.

2. If \( \Delta \) has more than 2 vertices, then it is the local interaction graph of a threshold activation function if and only if it contains no parallel arcs, is switching equivalent to an inward-homogeneous signed digraph, and its maximum in-degree is \( (n - 1) \).

Proof. There is \( \zeta \in \{ \pm 1 \}^{C_f} \) such that \( \zeta \Delta \) is inward-homogeneous. For each \( j \in C_f \), we will construct an \( a^j \in \mathbb{R}^{C_f} \) that defines the component threshold activation function \( f_j \) with the same in-star as \( j \) in \( \zeta \Delta \).

Let \( A \) be the set of all vertices with arcs going to \( j \) in \( \Delta \). Let the state \( x = \{1\}^{C_f} \). If \( A \) is empty, then let \( a^j = \{0\}^{C_f} \). This way, the in-star of \( j \) is empty in \( I_f(x) \) by proposition 2.6.1. So from now on, we will assume that \( A \) is non-empty.

First suppose every arc into \( j \) is positive in \( \zeta \Delta \). Let \( a^j_i = 1 \) for all \( i \in A \). Now since the in-degree of \( j \) is less than \( n \), there is at least one \( k \in C_f \setminus A \). Let \( a^j_k = 1 - |A| \) and \( a^j_i = 0 \) for all other \( i \in C_f \setminus (A \cup k) \). Observe that \( a^j \cdot x = 1 \) and \( a^j \cdot x^{i \rightarrow} = -1 \) for all \( i \in A \). This means that \( f_j(x) = 1 \) and \( f_j(x^{i \rightarrow}) = -1 \) for all \( i \in A \). So since \( a^j_i > 0 \) for all \( i \in A \), \( \partial^{i \rightarrow} f_j(x) = +1 \). Therefore the in-degree of \( j \) cannot be \( n \).
for all $i \in A$ by proposition [2.6.1]. Now $a^j \cdot x^{k \rightarrow} = 2|A| - 1 > 0$ since $A$ is non-empty. So $f_j(x) = f_j(x^{k \rightarrow}) = 1$. And since $a^j_i = 0$ for all other $i$ outside of $A$, $f_j(x) = f_j(x^{i \rightarrow}) = 1$ for all $i \in C_f \setminus A$. So $\partial^{i \rightarrow} f_j(x) = 0$ for all $i \notin A$. Therefore $I_f(x)$ has the same $j$ in-star as $\zeta \Delta$.

Similarly, if every arc into $j$ is negative in $\zeta \Delta$, then let $a^j_i = -1$ for all $i \in A$. Let $a^j_k = |A| - 1$ for some $k \in C_f \setminus A$ and $a^j_i = 0$ for all other $i \in C_f \setminus (A \cup k)$. Observe that $a^j \cdot x = -1$ and $a^j \cdot x^{i \rightarrow} = 1$ for all $i \in A$. Also, $a^j \cdot x^{k \rightarrow} = 1 - 2|A| < 0$. So since $a^j_i < 0$ for all $i \in A$, $\partial^{i \rightarrow} f_j(x) = -$ for all $i \in A$ by proposition [2.6.1]. Also, $\partial^{i \rightarrow} f_j(x) = 0$ for all $i \notin A$. Therefore $I_f(x)$ has the same $j$ in-star as $\zeta \Delta$ again.

So we have shown that $I_f(x) = \zeta \Delta$. Therefore $\Delta = \zeta I_f(x) = I_{f^\zeta}(\zeta x)$ by theorem [2.3.1]. So by lemma [2.6.4] $f^\zeta$ is also a threshold activation function. Since $\Delta = I_{f^\zeta}(\zeta x)$, $\Delta$ is indeed a local interaction graph of a threshold activation function.

\[\square\]
Chapter 3

Conjectures of René Thomas

3.1 Introduction

In the context of gene regulatory networks, René Thomas introduced two general conjectures in [7]. The first was that positive feedback in the gene regulatory network is a necessary condition for multistationarity in the dynamics of gene expression. The second was that negative feedback in the gene regulatory network is a necessary condition for stable periodicity. Signed directed graphs are often used as representations of gene regulatory networks. In fact, the main intended use of interaction graphs is to model gene regulatory networks. We interpret positive and negative feedback correspond to positive and negative cycles within an interaction graph. What we currently lack is an interpretation of multistationarity and stable periodicity.

A new directed graph called the state transition graph will be introduced in section 3.2. It will represent how gene expression levels or the values of components of an activation function can change over time. Then we will be able to define the features within the state transition graph that correspond to multistationarity and stable periodicity and relate these features to feedback in the interaction graph.

Versions of René Thomas’ conjectures have been proved by several authors in different contexts. The most relevant version for our discussion is Corollary 1 in [4]. Richard and Comet’s result uses a less general definition of the state transition graph then we use here however. In subsection 3.3.1 I will present and prove a new version of Thomas’ first conjecture using our more general definition of the state transition graph, but with a weaker conclusion. Then I show that that Richard and Comet’s result follows from this version of the Thomas’ first conjecture for a special class of activation functions that produce the same kind of state transition graphs used in [4].

In [5], Adrien Richard proved a version of the second conjecture. In subsection 3.3.2 I will present a new proof of Richard’s theorem.
In [5] and [2], Adrien Richard and others ask a local version of René Thomas’ second conjecture. Is there a Boolean activation function \( f \) with no steady states whose local interaction graphs contain no negative cycles? Paul Ruet gives an example of such an activation function in [6] that has twelve components. In section 3.4 I will give another example with only eight components.

### 3.1.1 State Transition Graphs

For an activation function \( f \), we will imagine the elements of \( S_f \) as the states of a discrete dynamical system where \( f \) is telling us how the states of the system can change over time. There are a few different ways to interpret this. For our sake we are interested in a dynamical system where the values of the components in \( C_f \) can only change one at a time. This is called asynchronous dynamics. More specifically while the system is in state \( x \), the system tends towards \( f(x) \). Since we only allow one component to change at a time, the next state in the system will be \( x_i \rightarrow f_i(x) \) for some \( i \in C_f \) such that \( f_i(x) \neq x_i \).

Now we give a graphical way to represent this asynchronous dynamical system. Given an activation function \( f \), the state transition graph of \( f \), denoted \( S_f \) is a directed graph derived from \( f \) whose vertex set is \( S_f \). If \( f_i(x) \neq x_i \), then the arc \((x, x_i \rightarrow f_i(x)) \in S_f \). In particular there are no arcs in \( S_f \) between states that differ in more than one coordinate and there are no loops.

For example, let the activation function \( f(x_1, x_2) = (1 - x_2^2, x_2(x_1 - 1)/2) \) where \( S_f = \{\pm 1\} \times \{0, \pm 1\} \). A picture of \( S_f \) is given in figure 3.1.

The focus of this chapter is relating features of the state transition graph \( S_f \) with features in the interaction graphs \( I_f \) and \( I_f(x, y) \).

It turns out that the state transition graphs of conjugate activation functions are related in a nice way.

**Theorem 3.1.1.** If \( \lambda = (\phi, \zeta) \) is a switching isomorphism of an activation function \( f \), then \( S_{f^\lambda} = \phi S_f \), i.e., \( S_f \) and \( S_{f^\lambda} \) are isomorphic.

**Proof.** Suppose the arc \((x, x_i \rightarrow f_i(x)) \in S_f \), i.e., \( f_i(x) \neq x_i \). All we need to do is show that the arc \((\lambda x, \lambda x_i \rightarrow f^\lambda_{\phi i}(\lambda x)) \) is also in \( S_{f^\lambda} \). This is sufficient since \( \lambda \) is an invertible map. Since \( \lambda : S_f \rightarrow S_{f^\lambda} \) is a bijection, \( \lambda_{\phi i} f(x) \neq \lambda_{\phi i} x \) if \( f_i(x) \neq x_i \). But \( \lambda_{\phi i} f(x) = f^\lambda_{\phi i}(\lambda x) \) by equation (2.7), so \( f^\lambda_{\phi i}(\lambda x) \neq \lambda x_{\phi i} \). Therefore the arc \((\lambda x, \lambda x_i \rightarrow f^\lambda_{\phi i}(\lambda x)) \in S_f \). \( \square \)
3.2 Attractors

Given a state transition graph $S_f$. We call $A \subseteq S_f$ an attractor of $S_f$ if $A$ is a terminal strong component of $S$. If $A$ contains a single state, then we call $A$ a steady state attractor, or just a steady state. Otherwise we call $A$ a periodic attractor.

For example in figure 3.1 the attractors are $\{(1, 0)\}$ and $\{(-1, -1), (-1, 1)\}$. The first is a steady state and the second is a periodic attractor.

If we are interested in how the states of the dynamical system change over time, then attractors are important. If the system is in a state contained in an attractor, then all future states will also be in the attractor. Furthermore, from any state in $S_f$ there is a path from that state into an attractor since attractors are terminal strong components of $S_f$.

**Proposition 3.2.1.** There is a dipath into an attractor from any vertex of $S_f$.

**Proof.** The proof is by induction on the number of strong components of $S_f$.

If $S_f$ contains a single connected component, then $S_f$ is strongly connected itself. So the proposition is true in this case.

Now assume that $S_f$ has $k > 1$ strong components and that the proposition is true for digraphs with fewer than $k$ strong components. Let $x \in S_f$. If $x$ is contained in an attractor, then we are done. So suppose that $x$ is not contained in an attractor. By proposition 1.1.8, $x$ is contained in some strong component $C$ of $S_f$. Since $C$ is not an attractor, there is an arc out of $C$ to another strong component $C'$ of $S_f$. Since the initial vertex of this arc is in $C$, there is a dipath from $x$ to $x'$, the terminal vertex of the arc. Therefore there is a dipath from $x$ to $x'$.
Now consider $S'$ the subgraph of $S_f$ induced by $S_f \setminus C$. Any strong component of $S_f$ other than $C$ is also a strong component of $S'$. To see this, let $y, y' \in S_f \setminus C$ be in the same strong component of $S_f$. Then there is still a dipath from $y$ to $y'$ and a dipath from $y'$ to $y$ in $S'$ since it is an induced subgraph.

This means that $S'$ has $k - 1$ strong components. So by induction, there is a dipath in $S'$ from $x'$ into an attractor $A \subseteq S'$. Combining this dipath with the dipath in $S_f$ from $x$ to $x'$ yields a dipath from $x$ to $A$. All that is left to show is that $A$ is also an attractor in $S_f$.

As we showed earlier, every strong component of $S_f$ other than $C$ is also a strong component of $S'$. Therefore $A$ is a strong component in $S_f$. Suppose that $A$ is not an attractor in $S_f$. Then there must be an arc from $A$ into $C$ since $A$ is terminal in $S'$. But then there would be a dipath from $A$ to $C$ and a dipath from $C$ to $A$. This means that $A$ and $C$ are in the same strong component of $S_f$ which is not the case. Therefore there are no arcs out of $A$ in $S_f$ either. Therefore $A$ is an attractor in $S_f$.

It turns out that some structure of the attractors can be related to cycles in the interaction graphs.

### 3.2.1 Thomas’ Conjectures

A gene regulatory network is a graphical representation of a system of interacting genes. Signed digraphs are often used as a model of gene regulatory networks. A positive arc from $i$ to $j$ in the network indicates that changing the expression level of gene $i$ can cause the expression level of gene $j$ to change in the same way. A negative arc from $i$ to $j$ in the network indicates that changing the expression level of gene $i$ can cause a contrary change in the expression level of gene $j$. Interaction graphs and state transition graphs are used as a model gene regulatory networks and the dynamics of gene expression respectively. The activation functions gives us a more concrete way to think about how gene expression levels change over time since we can think of a dynamical system in state $x$ tending towards state $f(x)$.

In the context of gene regulatory networks, the geneticist René Thomas made two conjectures.

**Conjecture 3.2.2** (Conjectures of René Thomas).

1. **Positive feedback is a necessary condition for multi-stability.**

2. **Negative feedback is a necessary condition for stable periodicity.**
3.3 Proofs of the Conjectures using Conjugate Activation Functions

For the proofs of our versions of Thomas’ conjectures, we will need restricted versions of a given activation function. If \( f \) is an activation function and the states \( s, t \in S_f \), let

\[
[s, t] = \{ x \in S_f | \min(s_i, t_i) \leq x_i \leq \max(s_i, t_i) \}.
\]

Also, \([s_i, t_i] = \{ a \in S_i | \min(s_i, t_i) \leq a \leq \max(s_i, t_i) \}\) for and \( i \in C_f \). Define \([s, t]|f\) to be the activation function with the same components as \( f \) whose state space is \([s, t]\), where

\[
[s, t]|f(x) = \begin{cases} 
  f_i(x) & \text{if } f_i(x) \in [s_i, t_i] \\
  \max(s_i, t_i) & \text{if } f_i(x) > \max(s_i, t_i) \\
  \min(s_i, t_i) & \text{if } f_i(x) < \min(s_i, t_i).
\end{cases}
\]

(3.1)

How are the interaction graphs of \([s, t]|f\) and \( f \) related?

**Lemma 3.3.1.** Given an activation function \( f \). Let \( g = [s, t]|f \). Then

1. \( I_g(x, y) \) is a subgraph of \( I_f(x, y) \).

2. \( I'_g(x, y) \) is a subgraph of \( I'_f(x, y) \).

**Proof.** For the first item, all we need to show is that if \((i, j, \sigma_{ij}) \in I_g(x, y)\), then \((i, j, \sigma_{ij}) \in I_f(x, y)\).

If \((i, j, \sigma_{ij}) \in I_g(x, y)\), then \( \sigma_{ij} = \partial^{i \rightarrow y} g_j(x) = \text{sgn}[g_j(x^{i \rightarrow y}) - g_j(x)][y_i - x_i] \). Also, \( \partial^{i \rightarrow y} f_j(x) = \text{sgn}[f_j(x^{i \rightarrow y}) - f_j(x)][y_i - x_i] \). So if \( g_j(x^{i \rightarrow y}) < f_j(x) \), we need only compare the quantities \( f_j(x^{i \rightarrow y}) - f_j(x) \) and \( g_j(x^{i \rightarrow y}) - g_j(x) \).

Since \( \sigma_{ij} \neq 0 \), \( g_j(x^{i \rightarrow y}) \neq g_j(x) \). Suppose \( g_j(x) < g_j(x^{i \rightarrow y}) \). Since the maximum value of \( g_j \) is \( \max(x_j, y_j) \), \( g_j(x^{i \rightarrow y}) \leq \max(x_j, y_j) \). So \( g_j(x) < \max(x_j, y_j) \). By equation (3.1), either \( f_j(x) = g_j(x) \), or \( f_j(x) = \min(x_j, y_j) > f_j(x) \). In either case \( g_j(x) \geq f_j(x) \).

Similarly, \( g(x^{i \rightarrow y}) = f(x^{i \rightarrow y}) \) or \( g(x^{i \rightarrow y}) = \max(x_i, y_i) < f(x^{i \rightarrow y}) \) by equation (3.1). So \( g(x^{i \rightarrow y}) \leq f_j(x^{i \rightarrow y}) \). Therefore \( f_j(x^{i \rightarrow y}) \geq g_j(x^{i \rightarrow y}) > g_j(x) \geq f_j(x) \), so \( \text{sgn}[f_j(x^{i \rightarrow y}) - \]
Given Lemma 3.3.2. If $g_j(x^{i\rightarrow y_i}) < g_j(x)$, the same argument with all inequalities reversed shows that $\partial^{i\rightarrow y_i} f_j(x) = \partial^{i\rightarrow y_i} g_j(x) = \sigma_{ij}$ again. So in either case, $(i, j, \sigma_{ij}) \in I_f(x, y)$.

For the second item, $I^*_g(x, y) = I_g(x, x^{\rightarrow y})$ and $I^*_f(x, y) = I_f(x, x^{\rightarrow y})$ by proposition 2.2.3. By the first item of this lemma, $I^*_g(x, x^{\rightarrow y})$ is a subgraph of $I^*_f(x, x^{\rightarrow y})$. Therefore $I^*_g(x, y)$ is a subgraph of $I^*_f(x, y)$.

Conjugating a restricted activation function behaves as you would expect.

**Lemma 3.3.2.** Given $\lambda = (\phi, \zeta)$, a switching isomorphism of an activation function $f$ and $s, t \in S_f$. If $g = [s, t]|f$, then $g^\lambda = [\lambda s, \lambda t]|f^\lambda$.

**Proof.** A straightforward, if tedious calculation is all that is needed to prove this. We will show the functions are equal component-wise. Let us start with the right hand side of the equation $g^\lambda = [\lambda s, \lambda t]|f^\lambda$. By equation (3.1),

$$[\lambda x, \lambda y]|f^\lambda_{\phi\delta}(\lambda x) = \begin{cases} f^\lambda_{\phi\delta}(\lambda x) & \text{if } f^\lambda_{\phi\delta}(\lambda x) \in [\lambda s, \lambda t] \\ \max(\lambda s, \lambda t) & \text{if } f^\lambda_{\phi\delta}(\lambda x) > \max(\lambda s, \lambda t) \\ \min(\lambda s, \lambda t) & \text{if } f^\lambda_{\phi\delta}(\lambda x) < \min(\lambda s, \lambda t) \end{cases}$$

Using equation (2.4),

$$[\lambda x, \lambda y]|f^\lambda_{\phi\delta}(\lambda x) = \begin{cases} f^\lambda_{\phi\delta}(\lambda x) & \text{if } f^\lambda_{\phi\delta}(\lambda x) \in [\zeta s, \zeta t] \\ \max(\zeta s, \zeta t) & \text{if } f^\lambda_{\phi\delta}(\lambda x) > \max(\zeta s, \zeta t) \\ \min(\zeta s, \zeta t) & \text{if } f^\lambda_{\phi\delta}(\lambda x) < \min(\zeta s, \zeta t) \end{cases}$$

And by equation (2.7),

$$[\lambda x, \lambda y]|f^\lambda_{\phi\delta}(\lambda x) = \begin{cases} \zeta f_i(x) & \text{if } \zeta f_i(x) \in [s_i, t_i] \\ \max(\zeta s, \zeta t) & \text{if } \zeta f_i(x) > \max(\zeta s, \zeta t) \\ \min(\zeta s, \zeta t) & \text{if } \zeta f_i(x) < \min(\zeta s, \zeta t) \end{cases}$$

For the left hand side of $g^\lambda = [\lambda s, \lambda t]|f^\lambda$, $g^\lambda_{\phi\delta}(\lambda x) = \zeta_i g_i(x)$ by equation (2.7). Using equation (3.1),

$$\zeta_i g_i(x) = \begin{cases} \zeta_i f_i(x) & \text{if } f_i(x) \in [s_i, t_i] \\ \zeta_i \max(s_i, t_i) & \text{if } f_i(x) > \max(s_i, t_i) \\ \zeta_i \min(s_i, t_i) & \text{if } f_i(x) < \min(s_i, t_i). \end{cases}$$
If $\zeta_i = +$, then

$$[\lambda x, \lambda y] f_{\phi i}^{\lambda}(\lambda x) = \begin{cases} 
\zeta_i f_i(x) & \text{if } \zeta_i f_i(x) \in [\zeta_i s_i, \zeta_i t_i] \\
\max(\zeta_i s_i, \zeta_i t_i) & \text{if } \zeta_i f_i(x) > \max(\zeta_i s_i, \zeta_i t_i) \\
\min(\zeta_i s_i, \zeta_i t_i) & \text{if } \zeta_i f_i(x) < \min(\zeta_i s_i, \zeta_i t_i) 
\end{cases}$$

$$= \begin{cases} 
- f_i(x) & \text{if } f_i(x) \in [-s_i, -t_i] \\
\max(-s_i, -t_i) & \text{if } f_i(x) > \max(-s_i, -t_i) \\
\min(-s_i, -t_i) & \text{if } f_i(x) < \min(-s_i, -t_i) 
\end{cases}$$

$$= g_i(x) \\
= \zeta_i g_i(x) \\
= g_{\phi i}^{\lambda}(\lambda x).$$

If $\zeta_i = -$, then

$$[\lambda x, \lambda y] f_{\phi i}^{\lambda}(\lambda x) = \begin{cases} 
\zeta_i f_i(x) & \text{if } \zeta_i f_i(x) \in [\zeta_i s_i, \zeta_i t_i] \\
\max(\zeta_i s_i, \zeta_i t_i) & \text{if } \zeta_i f_i(x) > \max(\zeta_i s_i, \zeta_i t_i) \\
\min(\zeta_i s_i, \zeta_i t_i) & \text{if } \zeta_i f_i(x) < \min(\zeta_i s_i, \zeta_i t_i) 
\end{cases}$$

$$= \begin{cases} 
- f_i(x) & \text{if } f_i(x) \in [s_i, t_i] \\
\max(s_i, t_i) & \text{if } f_i(x) > \max(s_i, t_i) \\
\min(s_i, t_i) & \text{if } f_i(x) < \min(s_i, t_i) 
\end{cases}$$

$$= g_i(x) \\
= \zeta_i g_i(x) \\
= g_{\phi i}^{\lambda}(\lambda x).$$
In either case, we have shown that $g^\lambda = [\lambda s, \lambda t] / f^\lambda$.

### 3.3.1 The First Conjecture

Now for the theorem that will lead to a new version of Thomas’ first conjecture.

**Theorem 3.3.3.** Let $f$ be an activation function. Suppose there are distinct states $x, y \in S_f$ such that $x$ is a steady state of $g = [x, y] / f$ and $g(x^{i\to y_i}) \neq x$ for all $i \in W$ where $W = \{i \in C_g \mid x_i \neq y_i\}$. Then $I_f^*(x, y)$ contains a positive cycle.

**Proof.** First we will prove a simpler version of the theorem.

**Lemma 3.3.4.** Let $f$ be an activation function. Suppose that $x, y \in S_f$, $x < y$, $x$ is a steady state of $g = [x, y] / f$ and $g(x^{i\to y_i}) > x$ for all $i \in W$ where $W = \{i \in C_g \mid x_i \neq y_i\}$. Then $I_f^*(x, y)$ contains a positive cycle whose vertices are in $W$.

**Proof.** For $i \in W$, $f$ will first show that there is a positive arc out of $i$ to some other vertex in $W$.

Since $g(x^{i\to y_i}) > x$, there is some $j \in C_g$ such that $g_j(x^{i\to y_i}) > x_j$. If $j \notin W$, then $x_j = y_j$. So $g_j(x^{i\to y_i}) = g_j(x) = x_j$ since $x$ is a steady state of $g$. Therefore $j \in W$.

Since $y_i > x_i$,

$$\partial^{i\to y_i} g_j(x) = \text{sgn}[g_j(x^{i\to y_i}) - g_j(x)](y_i - x_i) = \text{sgn}[g_j(x^{i\to y_i}) - g_j(x)].$$

Since $x$ is a steady state of $g$, $\partial^{i\to y_i} g_j(x) = \text{sgn}[g_j(x^{i\to y_i}) - x_j]$. And finally, since $g_j(x^{i\to y_i}) > x_j$, $\partial^{i\to y_i} g_j(x) = +$. Therefore, there is a positive arc from $i$ to $j$ in $I_f^*(x, y)$.

Finally, $I_g^*(x, y)$ is a subgraph of $I_f^*(x, y)$ by lemma 3.3.1. Therefore $I_f^*(x, y)$ contains a positive arc from $i$ to $j$ since $I_g^*(x, y)$ does.

Finally, Pick a vertex $i_1 \in W$. There is a positive arc out of $i_1$ to a new vertex $i_2$ in $W$. Since $i_2 \in W$, there is a positive arc from $i_2$ to $i_3 \in W$, and so on. Repeating this creates a sequence of vertices. Since $W$ is finite, there is a first repeated vertex $i_k$ in this sequence. Following the vertices from the first occurrence of $i_k$ in the sequence until it repeats closes a positive cycle in $I_f^*(x, y)$.

Back to the proof of theorem 3.3.3. The idea is to apply lemma 3.3.4 to $f^\zeta$, a conjugate activation function of $f$. Since $I_f^*(x, y)$ and $I_{f^\zeta}^*(\zeta x, \zeta y)$ are switching isomorphic by theorem 2.3.1, $I_f^*(x, y)$ will contain a positive cycle because $I_{f^\zeta}^*(\zeta x, \zeta y)$ contains a positive cycle by lemma 3.3.4.

Let $\zeta \in \{\pm\}^{C_f}$ such that $\zeta_i = +$ if $x_i \leq y_i$ and $\zeta_i = -$ otherwise. We have to show that $f^\zeta$ together with the states $\zeta x$ and $\zeta y$ satisfy all three hypotheses of lemma 3.3.4.
By the definition of $\zeta$, $\zeta_i x_i \leq \zeta_i y_i$ for all $i \in C_f$. Therefore $x < \zeta y$ since $x \neq y$. So the first hypothesis of lemma 3.3.4 is satisfied.

For the next hypothesis, we have to show that $\zeta x$ is a steady state of $[\zeta x, \zeta y] f^\zeta$. By lemma 3.3.2, $[\zeta x, \zeta y] f^\zeta = g^\zeta$. By equation 2.6, $g^\zeta(\zeta x) = \zeta g(x)$. Since $x$ is a steady state of $g$, $g^\zeta(\zeta x) = \zeta x$. Therefore $\zeta x$ is a steady state of $g^\zeta$.

Finally, we have to show that $g^\zeta(\zeta x^i \to^{\zeta \ast y_i}) > \zeta x$ for all $i \in W$. We know that $g(x^i \to^{\ast y_i}) \neq x$ for all $i \in W$ by assumption. So $\zeta x \neq \zeta g(x^i \to^{\ast y_i})$. Using equation (2.3), $\zeta g(x^i \to^{\ast y_i}) = \zeta g(x^i \to^{z_i})$ where $z = x \to^{\ast y_i}$. Using equation (2.6), $\zeta g(x^i \to^{z_i}) = g^\zeta(\zeta x^i \to^{z_i})$. Then $g^\zeta(\zeta x^i \to^{z_i}) = g^\zeta(\zeta x^i \to^{\zeta \ast y_i})$ by lemma 2.3.2. Since $\zeta z = z'$ where $z' = [\zeta x] \to^{\ast y_i} \zeta y$ by lemma 2.3.3, $z' = \zeta z_i$. So $g^\zeta(\zeta x^i \to^{\zeta \ast y_i}) = g^\zeta(\zeta x^i \to^{\ast y_i}) = g^\zeta(\zeta x^i \to^{\zeta \ast y_i})$ by equation (2.3). Therefore $g^\zeta(\zeta x^i \to^{\ast y_i}) = \zeta x$. Now by the definition of $\zeta$, $\zeta x$ is the minimum state in $[\zeta x, \zeta y]$. Therefore $g^\zeta(\zeta x^i \to^{\ast y_i}) > \zeta x$ since $g^\zeta(\zeta x^i \to^{\ast y_i}) \neq \zeta x$, and $g^\zeta(\zeta x^i \to^{\ast y_i}) \in [\zeta x, \zeta y]$.

We can now apply lemma 3.3.4 to $f^\zeta$ and the states $\zeta x$ and $\zeta y$. Therefore $I_f^\zeta(\zeta x, \zeta y)$ contains a positive cycle. Since $I_f^\zeta(\zeta x, \zeta y)$ and $I_f^\zeta(x, y)$ are switching isomorphic by theorem 2.3.1, $I_f(x, y)$ must also contain a positive cycle.

With this theorem, we can now prove a new version of Thomas’ first conjecture.

**Theorem 3.3.5 (Thomas’ First conjecture).** If an activation function $f$ has at least two steady states, then there exist states $x, y \in S_f$ such that $I_f^\zeta(x, y)$ contains a positive cycle.

**Proof.** If $z$ and $y$ are distinct steady states of $f$, then $z$ and $y$ are also steady states of $[z, y] f$. So pairs of distinct states $x$ and $y$ such that $x$ is a steady state of $[x, y] f$ certainly exist.

Among all such pairs, there is a pair of states that are closest. So let us assume that among all pairs of distinct states $(x, y)$ such that $x$ is a steady state of $g = [x, y] f$, $[x, y]$ is minimum. Let $W = \{i \in C_y | x_i \neq y_i\}$. It turns out that with these assumptions, $g(x^i \to^{y_i}) \neq x$ for all $i \in W$. So we can apply theorem 3.3.3 to get our result.

Let $z = x^i \to^{y_i}$ and $g' = [z, y] f$. It is possible that $z = y$. In this case, $g(x^i \to^{y_i}) = g(y) = y$ since $y$ is a steady state of $g$. So $g(x^i \to^{y_i}) \neq x$ for all $i \in W$ since $x$ and $y$ only differ in the $i$th component. So theorem 3.3.3 applies in this case.

Now suppose that $z \neq y$, i.e., $z$ and $y$ are distinct. If $x_i < y_i$, then by equation (3.1), $f(x^i \to^{y_i}) \leq x_i$ since $g_i(z) = x_i$. So $f_i(z) < z_i$. Similarly if $x_i > y_i$, then $f_i(z) > x_i$. In
Proof. Let \( x \) be as close as possible, i.e., \( |x, y| \) is minimal. Therefore \( g(x^{i \rightarrow y}) \neq x \) for all \( i \in W \). Now we can apply theorem 3.3.3 to \( f, x \) and \( y \) to get our result.

Call an activation function \( f \) unitary if for all states \( x \in S_f \) and \( i \in C_f \), either \( f_i(x) > x_i \), \( f_i(x) < x_i \), or \( f_i(x) = x_i \). So if \( f \) is unitary and \( (x, x^{i \rightarrow a}) \in S_f \), then \( a > x_i \), or \( a < x_i \) since \( a = f_i(x) \) and \( f_i(x) > x_i \) or \( f_i(x) < x_i \). So the state transition graphs of unitary activation functions are the same as those used in [4]. If we restrict to unitary activation functions, then corollary 1 in [4] follows from corollary 3.3.5.

**Corollary 3.3.6.** Given a unitary activation function \( f \). If \( S_f \) contains at least two attractors, then there are states \( x, y \in S_f \) such that \( I_f^*(x, y) \) contains a positive cycle.

**Proof.** Let \( A \) and \( B \) be attractors in \( S_f \). Let \( s \in A \) and \( t \in B \) be as close as possible, i.e., \( |s, t| \) is minimum. Let \( g = [s, t] f \). I claim that \( s \) and \( t \) are steady states of \( g \). If this is true, then we can apply corollary 3.3.5 to \( g \). Our result will then follow since \( I_f^*(x, y) \) is a subgraph of \( I_f^*(x, y) \) by lemma 3.3.1.

First, we will show that \( s \) is a steady state of \( g \). The situation of \( s \) and \( t \) is symmetric, so the same argument shows \( t \) is also a steady state of \( g \).

Suppose that \( g_i(s) \neq s_i \) for some \( i \in C_g \). This means that \( s_i \neq t_i \) since \( g_i(s) = s_i \) by equation (3.1) otherwise. Now \( g_i(s) \in [s_i, t_i] \) by equation (3.1) also. Suppose that \( s_i < t_i \). Since \( f \) is unitary, \( f_i(s) > s_i \) since \( g_i(s) = s_i \) otherwise. Similarly, if \( s_i > t_i \), then \( f_i(s) < s_i \). In either
Figure 3.2: $S_{-x}$ has two attractors, but $I_{-x}$ contains no positive cycle.

case, $f_i(s) \in [s_i, t_i]$. Therefore $f_i(s) = g_i(s)$ by equation (3.1). Since $g_i(s) \neq s_i$, $f_i(s) \neq s_i$ also. Therefore the arc $(s, s^{i\rightarrow f_i(s)}) \in S_f$. So the state $s^{i\rightarrow f_i(s)} \in A$. But $|s^{i\rightarrow f_i(s)}, t| < |s, t|$, a contradiction since $s$ and $t$ were chosen to be closest. Therefore $g_i(s) = s_i$ for all $i \in C_g$, i.e., $s$ is a steady state of $g$.

By symmetry, the same argument shows that $t$ is a steady state of $g$. So $g$ has at least two steady states. Therefore, there exist states $x, y \in S_g$ such that $I^*_g(x, y)$ contains a positive cycle by corollary 3.3.5. Finally, $I^*_g(x, y)$ is a subgraph of $I^*_f(x, y)$ by lemma 3.3.1 so $I^*_f(x, y)$ also contains a positive cycle.

You may ask whether unitary activation functions are required for corollary 3.3.6 to be true. It turns out that this version of Thomas’ first conjecture is not true in general. To see this, let $f(x) = -x$ where $S_f = \{0, \pm 1\}$. Then $S_f$ contains the two attractors $\{-1, 1\}$ and $\{0\}$. This can be seen in figure 3.2. But

$$\partial^{i\rightarrow a} f(x) = \text{sgn}[f(a) - f(x)][a - x] = \text{sgn}[a + x][a - x] = -\text{sgn}[a - x]^2 = -.$$ 

So $I_f$ consists of a single vertex and a single negative loop. So no local interaction graph of $f$ can contain a positive cycle. However, this does not contradict corollary 3.3.5 since $S_f$ does not contain multiple steady states.

### 3.3.2 The Second Conjecture

Theorem 1 in [5] proves that a negative cycle in the interaction graph $I_f$ is a necessary condition for $S_f$ to contain a periodic attractor. What follows is my own new proof of the theorem using techniques I have developed.

**Theorem 3.3.7.** If $S_f$ contains a periodic attractor, then $I_f$ contains a negative cycle.

**Proof.** I will prove the contrapositive statement.

**Lemma 3.3.8.** If $I_f$ is cycle-balanced, then every attractor in $S_f$ is a steady state.
\textbf{Proof.} We will prove several weaker intermediate versions of the lemma on our way to the final version. We begin with a result about positive interaction graphs.

\textbf{Lemma 3.3.9.} \emph{Given an activation function \( f \) such that every arc in \( I_f \) is positive. If \( x, y \in S_f \) and \( x \leq y \), then \( f(x) \leq f(y) \), i.e., \( f \) is isotone.}

\textbf{Proof.} We proceed by induction on \( |x, y| \).

If \( |x, y| = 1 \) then \( y = x^{i \rightarrow a} \) for some \( i \in C_f \). Since every arc in \( I_f \) is positive, \( \partial_i \rightarrow a f_j(x) = \text{sgn}(f_j(x^{i \rightarrow a}) - f_j(x))[a - x_i] \geq 0 \) for any \( j \in C_f \). Now \( y > x \), so \( a > x_i \). Therefore \( \partial_i \rightarrow a f_j(x) = \text{sgn}(f_j(x^{i \rightarrow a}) - f_j(x)) \), so \( f_j(x^{i \rightarrow a}) \geq f_j(x) \). This is true for all \( j \in C_f \), so \( f(x^{i \rightarrow a}) \geq f(x) \) i.e., \( f(y) \geq f(x) \).

Now suppose \( |x, y| = k \) and that the lemma is true when the distance is less than \( k \). Say \( x \) and \( y \) differ in the \( i \)th coordinate. Since \( y > x \), \( y_i > x_i \). But \( |y, x^{i \rightarrow y_i}| = k - 1 \) since \( |x, y| = k \). By the induction hypothesis, \( f(y) \geq f(x^{i \rightarrow y_i}) \). Since \( f(x^{i \rightarrow y_i}) \geq f(x) \) by the first step of the induction, it follows that \( f(y) \geq f(x) \). \( \square \)

We are now ready to prove the weakest version of lemma 3.3.8.

\textbf{Lemma 3.3.10.} \emph{If every arc in \( I_f \) is positive, then every attractor in \( S_f \) is a steady state.}

\textbf{Proof.} Let \( \mathcal{A} \) be an attractor in \( S_f \) and let \( x \) be a maximal state in \( \mathcal{A} \). I claim that \( f(x) \leq x \). To see this, suppose that \( f_j(x) > x_j \) for some \( j \in C_f \). So \( (x, x^{i \rightarrow f_j(x)}) \in S_f \). This means that \( x^{i \rightarrow f_j(x)} \in \mathcal{A} \). But \( x^{i \rightarrow f_j(x)} > x \). Since \( x \) is a maximal state in \( \mathcal{A} \), \( x^{i \rightarrow f_j(x)} \) cannot be contained in \( \mathcal{A} \). Therefore \( f_j(x) \leq x_j \). Since this is true for each component of \( f \), it follows that \( f(x) \leq x \).

Now suppose that \( x \) is not a steady state. So \( f(x) < x \). In particular, \( f_i(x) < x_i \) for some \( i \in C_f \). This means that there is an arc \( (x, x') \in S_f \) where \( x' = x^{i \rightarrow f_i(x)} \) and \( x' \in \mathcal{A} \).

I claim that there is no directed path from \( x' \) back to \( x \), a contradiction since \( x \in \mathcal{A} \). To prove the claim, I will show that if \( y \) is any state such that there is a directed path from \( x' \) to \( y \) in \( S_f \), then \( y_i < x_i \).

Let \( y \) and \( z \) be consecutive states in a dipath in \( S_f \) that starts at \( x' \) where \( y < x \) and \( y_i < x_i \) in particular. I claim that this implies that \( z < x \) and \( z_i < x_i \) also. Since \( y < x \), \( f(y) \leq f(x) \) by lemma 3.3.9. This means \( f_i(y) \leq f_i(x) < x_i \).

The arc \( (y, z) \in S_f \), so \( z = y^{j \rightarrow f_j(y)} \) for some \( j \in C_f \). If \( j = i \), then \( z_i = f_i(y) < x_i \). If \( j \neq i \), then \( z_i = y_i < x_i \). Either way \( z_i < x_i \). Also, since \( y < x \) and \( f_j(y) \leq f_j(x) \leq x_j \), \( z = y^{j \rightarrow f_j(y)} \leq x \). So \( z < x \) and \( z_i < x_i \) in particular, establishing the claim.

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Now for the proof of this lemma, we are assuming that \( x' < x \) and \( x'_i < x_i \). So for any state \( y \) such that there is a directed path from \( x' \) to \( y \) in \( S_f \), it follows by the argument above that \( y_i < x_i \) since this property holds for consecutive states in such a dipath. Therefore there is no directed path from \( x' \) to \( x \). This cannot happen since \( x \in A \). Therefore \( f(x) = x \), so \( A = \{x\} \) is a steady state.

This lemma can be immediately improved.

**Lemma 3.3.11.** If \( \mathcal{I}_f \) is balanced, then every attractor in \( S_f \) is a steady state.

**Proof.** By theorem 1.3.2, there is \( \zeta \) such that every arc in \( \zeta \mathcal{I}_f \) is positive. By corollary 2.3.1 \( \zeta \mathcal{I}_f = \mathcal{I}_{f^\zeta} \). Therefore every attractor in \( S_{f^\zeta} \) is a steady state by lemma 3.3.10. By theorem 3.1.1 \( S_{f^\zeta} \) and \( S_f \) are isomorphic, so every attractor in \( S_f \) is also a steady state.

We are now ready to tackle the final version of lemma 3.3.8. There are two cases, \( \mathcal{I}_f \) is strongly connected, or not. In the former case, we can use the following result.

**Lemma 3.3.12** ([Corollary 13.11a]). A strongly connected signed digraph is cycle-balanced if and only if it is balanced.

So if \( \mathcal{I}_f \) is strongly connected, it is balanced by lemma 3.3.12 since it is cycle-balanced. Then all attractors in \( S_f \) are steady states by lemma 3.3.11.

Suppose \( \mathcal{I}_f \) is not strongly connected. Let \( f \) be an activation function whose interaction graph is cycle-balanced and not strongly connected, and such that \( S_f \) contains periodic attractors. In addition, suppose that among all activation functions with these properties, \( |S_f| \) is minimum. We will show that from any state in an attractor \( A \) of \( f \), there is a dipath to a steady state of \( f \). This will mean that the attractor must have been a steady state to begin with.

We will use a special kind of restricted activation function for the remainder of the proof. Given an activation function \( f \), \( x \in S_f \) and \( W \subseteq C_f \). Let \( s \in S_f \) such that \( s_i = \min S_i \) for \( i \in W \) and \( s_i = x_i \) otherwise. Let \( t \in S_f \) such that \( t_i = \max S_i \) for \( i \in W \) and \( t_i = x_i \) otherwise. Similar to equation (3.1), define

\[
[W, x]|f := [s, t]|f.
\]

**Lemma 3.3.13.** Given an activation function \( f \), \( x \in S_f \) and \( W \subseteq C_f \). If \( g = [W, x]|f \), then \( S_g \subseteq S_f \).
Proof. Suppose the arc \((y, y^{i\rightarrow a})\) ∈ \(\mathcal{S}_g\), i.e., \(g_i(y) = a\). Note that \(i \notin W\), then \(a = x_i = y_i\) since there is only one possible output of \(g_i\) in this case. So there would be no arc if this were the case. Therefore \(i \in W\). By equation (3.1), \(g_i(y) = f_i(y)\) since \(s_i\) and \(t_i\) are the minimum and maximum values of \(S_i\) respectively. Therefore \(f_i(y) = a\), so \((y, y^{i\rightarrow a}) \in \mathcal{S}_f\) also.

Now let \(\mathcal{A}\) be an attractor of \(f, x\) be a state in \(\mathcal{A}\), and let \(W\) be an initial strong component of \(\mathcal{I}_f\). Let \(g = [W, x]|f\). By lemma 3.3.1, \(\mathcal{I}_g\) is a subgraph of \(\mathcal{I}_f\). For this particular activation function \(g\), we can actually say more than this.

First, observe that for all \(j \in W\), \(f_j\) depends only on inputs in \(W\). To see this, let \(j \in W\). If \(k \notin W\), there is no arc from \(k\) to \(j\) in \(\mathcal{I}_f\) since \(W\) is an initial strong component. Therefore

\[
\partial^{k\rightarrow a} f_j(x) = \text{sgn}[f_j(x^{k\rightarrow a}) - f_j(x)][a - x_k] = 0.
\]

If \(x_k \neq a\), \(f_j(x^{k\rightarrow a}) - f_j(x) = 0\), so \(f_j(x^{k\rightarrow a}) = f_j(x)\). So \(f_j\) does not depend on inputs outside of \(W\).

Now suppose \((i, j, \sigma_{ij}) \in \mathcal{I}_f\) where \(i, j \in W\). So there is \(y \in \mathcal{S}_f\) such that \(\partial^{i\rightarrow a} f_j(y) = \sigma_{ij}\). Let \(z \in \mathcal{S}_f\) such that \(z_k = y_k\) for all \(k \in W\) and \(z_k = x_k\) otherwise. Since \(f_j\) only depends on inputs from \(W\),

\[
\partial^{i\rightarrow a} f_j(y) = \text{sgn}[f_j(y^{i\rightarrow a}) - f_j(y)][y_i - a]
= \text{sgn}[f_j(z^{i\rightarrow a}) - f_j(z)][z_i - a]
= \text{sgn}[g_j(z^{i\rightarrow a}) - g_j(z)][z_i - a]
= \partial^{i\rightarrow a} g_j(z).
\]

Therefore \(\mathcal{I}_g\) contains every arc between vertices in \(W\) that is in \(\mathcal{I}_f\). So the subgraph of \(\mathcal{I}_g\) on \(W\) is actually the induced subgraph of \(\mathcal{I}_f\) on \(W\).

Since \(W\) is a strong component of \(\mathcal{I}_f\), the subgraph of \(\mathcal{I}_g\) on \(W\) is strongly connected. Since \(\mathcal{I}_g\) is cycle-balanced, this strong component of \(\mathcal{I}_g\) is balanced by 3.3.12. Since the components in \(C_g \setminus W\) can take only a single value, there are no arcs to or from these vertices in \(\mathcal{I}_g\), i.e., they are all isolated vertices. Therefore \(\mathcal{I}_g\) is balanced. So all attractors in \(\mathcal{S}_g\) are steady states by lemma 3.3.11. Therefore, there is a dipath \(P_1\) from \(x\) to the steady state \(x'\) in \(\mathcal{S}_g\) by lemma 3.2.1.

Now let \(g' = [C_f \setminus W, x']|f\). By lemma 3.3.1 again, \(\mathcal{I}_{g'}\) is a subgraph of \(\mathcal{I}_f\). So \(\mathcal{I}_{g'}\) is cycle-balanced. Now if \(\mathcal{I}_{g'}\) is strongly connected, then it is balanced by lemma 3.3.12. So all attractors of \(g'\) are steady states by lemma 3.3.11. If \(\mathcal{I}_{g'}\) is not strongly connected, then \(g'\) is an activation function whose interaction graph is cycle-balanced and not strongly connected such that \(|\mathcal{S}_{g'}| < 58\).
Therefore all attractors of $g'$ are steady states by the minimality of $f$. Either way, there is a dipath $P_2$ in $S_{g'}$ from $x'$ to $x''$, a steady state of $g'$.

I claim that $x''$ is a steady state of $f$. To see this, let $k \in W$. Since there is only one possible value for the components of $g'$ in $W$, $x''_k = x'_k$. Now $f_k$ depends only on inputs in $W$, so $f_k(x'') = f_k(x')$. By the definition of $g$, $f_k(x') = g_k(x')$. Since $x'$ is a steady state of $g$, $g_k(x') = x'_k = x''_k$. Therefore $f_k(x'') = x''_k$.

If instead $k \notin W$, then $g'_k(x'') = x''_k$ since $x''$ is a steady state of $g'$. But $g'_k(x'') = f_k(x)$ by the definition of $g$, so $f_k(x'') = x''_k$ in this case also. Therefore $x''$ is indeed a steady state of $f$.

Finally, by lemma 3.3.13, $P_1$ and $P_2$ are also contained in $S_f$. Therefore $P_1 \cup P_2 \subset S_f$ is a dipath from $x$ to $x''$, a steady state of $f$. Remember that we assumed $x$ is contained in the attractor $A \subseteq S_f$. The only way it is possible for there to be a path from a state in an attractor to a steady state is if $A$ is a steady state to begin with. Therefore all attractors in $S_f$ are steady states.

This completes the proof of the contrapositive statement of theorem 3.3.7.

### 3.4 The Local Version of the Second Conjecture

#### 3.4.1 A Boolean Counterexample

In [5], Adrian Richard proves the following.

**Theorem 3.4.1 ([5, Corollary 1]).** If an activation function $f$ has no steady states, then $I_f$ contains a negative cycle.

Adrian Richard then asks a natural follow-up question. If an activation function $f$ has no steady states, are there states $x, y \in S_f$ such that the local interaction graph $I_f(x, y)$ contains a negative cycle? Put another way, if every local interaction graph of $f$ is cycle-balanced, must $f$ have a steady state? A non-Boolean counterexample to this question is given in example 6 in [5]. However the question was left open in the boolean case, i.e., when each component of $f$ can takes only two values. The boolean version of the question is also asked in [2]. In [6], Paul Ruet gives a Boolean counterexample on twelve components. I will now present another Boolean activation function that has no steady states and for which every local interaction graph is cycle-balanced, but on only eight components. I suspect that there are no counter examples on fewer components, but I have not been able to prove this yet.
For this activation function \( f \), \( C_f = [8] := \{1, 2, \ldots, 8\} \), \( S_i = \{\pm 1\} \) and \( S_f = \{\pm 1\}^8 \).

Since the activation function is boolean and has eight components, it has 256 local interaction graphs. Deriving each of these would be impractical so we will develop some tools to streamline the presentation.

Given an activation function \( f \) and \( \lambda \), a signed permutation of \( C_f \) such that \( \lambda S_f = S_f \). We say that \( \lambda \) commutes with \( f \) if \( f(\lambda x) = \lambda f(x) \) for all states \( x \in S_f \). Note that \( \lambda x \in S_f \) since \( \lambda S_f = S_f \).

So \( f(\lambda x) \) makes sense here.

**Lemma 3.4.2.** Let \( f \) be an activation function and let \( \lambda \) be a signed permutation of \( C_f \) such that \( \lambda S_f = S_f \). If \( f \) commutes with \( \lambda \), then \( I_f(\lambda x, \lambda y) = \lambda I_f(x, y) \), i.e., \( I_f(x, y) \) and \( I_f(\lambda x, \lambda y) \) are switching isomorphic.

**Proof.** We use the conjugate activation function \( f^\lambda \). Under these hypotheses, \( f = f^\lambda \). To see this we use equation (2.6)

\[
f(x) = f(\lambda \lambda^{-1} x) \\
= \lambda f(\lambda^{-1} x) \\
= f^\lambda(\lambda \lambda^{-1} x) \\
= f^\lambda(x).
\]

Note that \( f(\lambda^{-1} x) \) makes sense since \( \lambda^{-1} x \in S_f \) because \( \lambda S_f = S_f \). Since \( f = f^\lambda \), by lemma 2.3.1

\[
I_f(\lambda x, \lambda y) = I_f(\lambda x, \lambda y) \\
= \lambda I_f(x, y).
\]

So \( I_f(\lambda x, \lambda y) \) and \( I_f(x, \lambda y) \) are switching isomorphic.

So in the presentation of this counterexample, we will have a group \( \Gamma \) of signed permutations of \([8]\) such that every element of \( \Gamma \) commutes with our activation function \( f \). If \( I_f(x) \) is cycle-balanced, then the local interaction graph of any state in the orbit of \( x \) is also cycle-balanced by lemma 3.4.2. So we need only pick a set of distinct orbit representatives and check whether the local interaction graphs at each representative state is cycle-balanced. The action of \( \Gamma \) on \( \{\pm 1\}^n \) will have eight orbits. So we will only have to check eight local interaction graphs instead of 256.

Before we can put this plan into action, we have to address how multiplication works in the signed permutation group based on our definitions. Recall definition (2.5) which defines how a
switching isomorphism of an activation function $f$ transforms states of $f$. The set of signed permutations of $f$ forms a group with group product being composition. But how does this composition work? Let $\lambda_1 = (\zeta^1, \pi_1)$ and $\lambda_2 = (\zeta^2, \pi_2)$ be signed permutations of an activation function $f$. Using equation (2.5) several times,

\[
\lambda_2(\lambda_1 x) = \lambda_2(\pi_1 [\zeta^1 x]) = \pi_2(\zeta^2 [\pi_1 [\zeta^1 x]]) = \pi_2(\pi_1^{-1}(\zeta^2 [\pi_1 [\zeta^1 x]]) = \pi_2(\pi_1^{-1}(\zeta^2 \pi_1 [\zeta^1 x])) = \pi_2(\pi_1^{-1}(\zeta^2)(\pi_1^{-1} \pi_1 [\zeta^1 x])) = \pi_2(\pi_1^{-1}(\zeta^2)\zeta^1 x) = (\pi_2 \pi_1, (\pi_1^{-1}(\zeta^2)\zeta^1) x).
\]

So we have shown that for elements of the signed permutation group,

\[
\lambda_2 \lambda_1 = (\pi_2, \zeta^2)(\pi_1, \zeta^1) = (\pi_2 \pi_1, (\pi_1^{-1}(\zeta^2)\zeta^1)). \tag{3.3}
\]

The identity element of the group is $e = (\pi^0, \{+\}^C)$ where $\pi^0$ is the identity permutation. What about inverses? Using equation (3.3).

\[
(\pi^{-1}, \pi \zeta)(\pi, \zeta) = (\pi^{-1} \pi, [\pi^{-1} \pi \zeta] \zeta) = (\pi^0, \zeta \zeta) = e.
\]

Similarly $(\pi, \zeta)(\pi^{-1}, \pi \zeta) = e$, so $(\zeta, \pi)^{-1} = (\pi \zeta, \pi^{-1})$.

Here is a useful formula for powers in this group.

**Lemma 3.4.3.** If $\lambda = (\pi, \zeta)$ is a signed permutation, then

\[
\lambda^n = \left(\pi^n, \Pi_{i=1}^n (\pi^{-i} \zeta)\right).
\]
Proof. The proof is by induction. The formula is valid for \( n = 1 \) since

\[
\lambda^1 = (\pi^1, \prod_{i=1}^1 \pi^{1-i} \zeta)
\]

\[
= (\pi, \pi^0 \zeta)
\]

\[
= (\pi, \zeta)
\]

\[
= \lambda.
\]

Now suppose that the formula is true for \( n < k \). Using equation (3.3),

\[
\lambda^k = \lambda \lambda^{k-1}
\]

\[
= (\pi, \zeta) \left( \pi^{k-1}, \prod_{i=1}^{k-1} \pi^{1-i} \zeta \right)
\]

\[
= \left( \pi \pi^{k-1}, \pi^{-(k-1)} \zeta \prod_{i=1}^{k-1} \pi^{1-i} \zeta \right)
\]

\[
= \left( \pi^k, \pi^{1-k} \zeta \prod_{i=1}^{k-1} \pi^{1-i} \zeta \right)
\]

\[
= \left( \pi^k, \prod_{i=1}^k \pi^{1-i} \zeta \right).
\]

Therefore the formula is also true for \( k \). So the formula is valid for all \( n \) by induction. \( \Box \)

Let \( \Gamma = \langle \lambda, \zeta \rangle \), the group generated by \( \zeta = (-,+,+,+,+,+,+,+) \) and \( \lambda = [\pi, \zeta'] = [(1 2 3 4 5 6 7 8), (+,+,+,+,+,+,+,+)] \). Note that for \( \pi \) we are using the cycle notation for permutations. The group \( \Gamma \) will commute with our activation function. So it demands some detailed analysis. We will show that this group has order 32 and its action on \( \{\pm1\}^8 \) is free.

Let us first show that the order of \( \Gamma \) is 32. By equation (2.3), \( \pi_j x = x_{\pi^{-1}j} \). It follows that \( \pi(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8) = (x_8, x_1, x_2, x_3, x_4, x_5, x_6, x_7) \). For powers of \( \lambda \), using lemma 3.4.3

\[
\lambda^2 = [\pi^2, (\pi^{-1} \zeta') \zeta']
\]

\[
= [\pi^2, (+,+,+,+,+,+,-,-,+) \zeta']
\]

\[
= [\pi^2, (+,+,+,+,+,+,+,+)].
\]

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Similarly, for the remaining powers,

\[
\begin{align*}
\lambda^3 &= [\pi^3, (+, +, +, +, +, -)] \\
\lambda^4 &= [\pi^4, (+, +, +, +, -)] \\
\lambda^5 &= [\pi^5, (+, +, +, +, -)] \\
\lambda^6 &= [\pi^6, (+, +, +, +, -)] \\
\lambda^7 &= [\pi^7, (+, +, +, +, -)] \\
\lambda^8 &= (−, −, −, −, −, −)
\end{align*}
\]

So \(\lambda^{16}\) is the identity, i.e the order of \(\lambda\) is 16. By equation (3.3),

\[
\begin{align*}
\zeta \lambda &= [\pi, (\pi^{-1} \zeta)] \\
&= [\pi, \zeta'](\pi^{-1} \zeta) \\
&= \lambda(+, −, +, +, +, −) \\
&= \lambda(−, +, −, −, −, −)(−, −, +, −, +, −) \\
&= \lambda{\lambda^8} \zeta \\
&= \lambda^9 \zeta.
\end{align*}
\]

So each \(\gamma \in \Gamma\) can be written in the form \(\lambda^n \zeta^m\) for \(n \in [16]\) and \(m = 0, 1\) since \(\zeta\) has order 2. Therefore \(|\Gamma| = 32\).

It turns out that the action of \(\Gamma\) on \(S_f\) is free. So when \(\Gamma\) acts on \(S_f\), each orbit of \(\Gamma\) will contains thirty two states. Since \(|S_f| = 256 = 32 \times 8\), there are eight orbits in total. We will need a set of distinct representative states for the orbits of \(\Gamma\). I claim that the following states will forms a set of distinct representatives.

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\[ \begin{align*}
x_1 &= (1, 1, 1, 1, 1, 1, 1, 1) \\
x_2 &= (1, -1, 1, 1, 1, 1, 1, 1) \\
x_3 &= (1, 1, -1, 1, 1, 1, 1, 1) \\
x_4 &= (1, 1, 1, -1, 1, 1, 1, 1) \\
x_5 &= (1, 1, 1, 1, -1, 1, 1, 1) \\
x_6 &= (1, 1, 1, 1, 1, -1, 1, 1) \\
x_7 &= (1, 1, -1, -1, 1, 1, 1, 1) \\
x_8 &= (1, 1, -1, 1, -1, 1, 1, 1)
\end{align*} \]

To show that \( \{x_1, \ldots, x_8\} \) forms a set of distinct representatives, we just look at each set of states \( \{\gamma x_i \mid \gamma \in \Gamma, i \in [8]\} \) and confirm that these sets partition \( S_f \) into eight blocks, each containing 32 states. This will also show that this group action is free. The orbit of \( x_1 \) is given in table \( 3.4.1 \) It occupies half a page of space, so four pages are needed to print each orbit. So to confirm that we have a set of distinct orbit representatives and that the action of \( \Gamma \) on \( S_f \) is free, we just check that each table contains 32 distinct states and that there is no overlap between the tables. Easy! But in the interest of the reader’s eyes and the Amazon rain forest, or server space more likely, let us try to be more concise.

For \( x \in S_f \), since \( \zeta = (-, +, -, +, -, +, +) \in \Gamma \), either \( x_7 = 1 \) or \( (\zeta x)_7 = 1 \). Similarly \( \zeta^{8} = (+, -,-,+,+,-,+) \in \Gamma \), so either \( x_8 = 1 \) or \( (\zeta^{8} x)_8 = 1 \). So each state in \( S_f \) is equivalent to a state \( x \) such that \( x_7 = x_8 = 1 \) under the action of \( \Gamma \). So we do not need to include states whose seventh or eighth coordinates are negative in the table.

To save even more space, we will use the following function. For \( x \in S_f \), let \( N(x) = \{i \mid x_i = -1\} \). If we exclude states whose seventh or eighth coordinate is negative, \( N(x) \subseteq [6] \). Table \( 3.2 \) contains each orbit of \( \Gamma \) where \( N(\lambda^n x) \) is represented by \( N(\lambda^n x) \), \( N(\zeta \lambda^n x) \), \( N(\lambda^n+8 x) \) or \( N(\zeta \lambda^{n+8} x) \), depending on which one is a subset of \( [6] \). As you can see, the table contains every subset of \( [6] \) and there is no overlap between the rows. This justifies that \( \{x_1, \ldots, x_8\} \) forms a set of distinct representatives and that the action of \( \Gamma \) on \( S_f \) is free.

Now we are ready to define the activation function \( f \) itself. The first defining property of \( f \) is that it commutes with \( \Gamma \). Now every state \( y = \gamma x_i \) for some \( \gamma \in \Gamma \) and orbit representative \( x_i \).

Since \( f \) commutes with \( \Gamma \), \( f(y) = f(\gamma x_i) = \gamma f(x_i) \). So to completely define \( f \), we only need to specify the outputs of \( f \) on each orbit representative. This will determine the output of \( f \) on every
\begin{table}[h]
\centering
\begin{tabular}{|c|c|c|}
\hline
$n$ & $\lambda^n x_1$ & $\lambda^n \zeta x_1$ \\
\hline
0 & \( (1, 1, 1, 1, 1, 1, 1) \) & \( (-1, 1, -1, 1, -1, 1, -1) \) \\
1 & \( (-1, 1, 1, 1, 1, 1, 1) \) & \( (-1, -1, 1, -1, 1, -1, 1) \) \\
2 & \( (-1, -1, 1, 1, 1, 1, 1) \) & \( (1, -1, -1, 1, -1, 1, -1) \) \\
3 & \( (-1, -1, -1, 1, 1, 1, 1) \) & \( (-1, 1, -1, -1, 1, -1, 1) \) \\
4 & \( (-1, -1, -1, -1, 1, 1, 1) \) & \( (1, -1, -1, -1, 1, -1, 1) \) \\
5 & \( (-1, -1, -1, -1, -1, 1, 1) \) & \( (-1, -1, 1, -1, -1, 1, -1) \) \\
6 & \( (-1, -1, -1, 1, -1, 1, 1) \) & \( (1, -1, -1, 1, -1, 1, -1) \) \\
7 & \( (-1, -1, 1, -1, 1, -1, 1) \) & \( (1, -1, 1, -1, 1, -1, 1) \) \\
8 & \( (-1, 1, -1, 1, -1, 1, -1) \) & \( (1, -1, -1, 1, -1, 1, -1) \) \\
9 & \( (1, -1, -1, 1, -1, 1, -1) \) & \( (-1, -1, 1, -1, 1, -1, 1) \) \\
10 & \( (1, 1, -1, 1, -1, 1, -1) \) & \( (1, 1, 1, -1, 1, -1, 1) \) \\
11 & \( (1, 1, -1, -1, 1, -1, 1) \) & \( (1, -1, 1, -1, 1, -1, 1) \) \\
12 & \( (1, 1, 1, -1, 1, -1, 1) \) & \( (-1, 1, -1, 1, -1, 1, -1) \) \\
13 & \( (1, 1, 1, 1, -1, 1, -1) \) & \( (1, -1, -1, 1, -1, 1, -1) \) \\
14 & \( (1, 1, 1, 1, 1, -1, 1) \) & \( (-1, 1, -1, -1, 1, -1, 1) \) \\
15 & \( (1, 1, 1, 1, 1, 1, 1, -1) \) & \( (1, -1, -1, 1, -1, 1, 1, -1) \) \\
\hline
\end{tabular}
\caption{Table 3.1: Orbit of $x_1$}
\end{table}

\begin{table}[h]
\centering
\begin{tabular}{|c|c|c|c|c|c|c|c|c|}
\hline
$i$ & \( N(x_i) \) & \( N(\lambda x_i) \) & \( N(\lambda^2 x_i) \) & \( N(\lambda^3 x_i) \) & \( N(\lambda^4 x_i) \) & \( N(\lambda^5 x_i) \) & \( N(\lambda^6 x_i) \) & \( N(\lambda^7 x_i) \) \\
\hline
1 & \{ \} & \{1\} & \{1, 2\} & \{1, 2, 3\} & \{1, 2, 3, 4\} & \{1, 2, 3, 4, 5\} & \{1, 2, 3, 4, 5, 6\} & \{1, 2, 4, 6\} \\
2 & \{2\} & \{1, 3\} & \{1, 2, 4\} & \{1, 2, 3, 5\} & \{1, 2, 3, 4, 6\} & \{2\} & \{1, 3, 4\} & \{1, 2, 4, 6\} \\
3 & \{3\} & \{1, 4\} & \{1, 2, 5\} & \{1, 2, 3, 6\} & \{2, 4\} & \{1, 3, 5\} & \{1, 2, 4, 5\} & \{1, 2, 4, 5, 6\} \\
4 & \{4\} & \{1, 5\} & \{1, 2, 6\} & \{2\} & \{1, 3, 6\} & \{2, 3, 4\} & \{1, 3, 4, 5\} & \{1, 2, 4, 6\} \\
5 & \{5\} & \{1, 6\} & \{2, 3, 5\} & \{1, 3, 4, 6\} & \{2, 3, 4\} & \{1, 3, 4\} & \{1, 2, 4\} & \{1, 2, 3, 6, 5\} \\
6 & \{6\} & \{3, 5\} & \{1, 4, 6\} & \{2, 3\} & \{1, 3, 4\} & \{1, 2, 4\} & \{1, 2, 3, 6\} & \{2, 4, 5, 6\} \\
7 & \{3, 4\} & \{1, 4, 5\} & \{1, 2, 5, 6\} & \{2, 5, 6\} & \{5, 6\} & \{3, 5, 6\} & \{3, 4, 5, 6\} & \{3, 4, 6\} \\
8 & \{3, 6\} & \{3, 4, 5\} & \{1, 4, 5, 6\} & \{2, 3, 6\} & \{4, 5\} & \{1, 5, 6\} & \{2, 3, 5, 6\} & \{4, 5, 6\} \\
\hline
\end{tabular}
\caption{Table 3.2: Orbits of $\Gamma$}
\end{table}
state via the group action.

Here are the outputs of $f$ for each orbit representative.

$$f(x_1) = f(x_2) = (-1, 1, 1, 1, 1, 1, 1)$$
$$f(x_3) = f(x_4) = f(x_5) = x_1$$
$$f(x_6) = (1, 1, 1, 1, 1, 1, -1)$$
$$f(x_7) = x_4$$
$$f(x_8) = (-1, -1, 1, 1, -1, 1, 1).$$

For any $\gamma \in \Gamma$, $\mathcal{I}_f(x)$ and $\mathcal{I}_f[\gamma x]$ are switching isomorphic by lemma 3.4.2. Therefore, if $x$ and $y$ are in the same orbit of $\Gamma$, $\mathcal{I}_f(x)$ and $\mathcal{I}_f(y)$ are switching isomorphic. Also, since $f(x) \neq x$ for each orbit representative, $f[\gamma x] = \gamma f(x) \neq \gamma(x)$ since the action of $\Gamma$ is free. Therefore $f$ has no steady states, as required.

Next we need to show that $\mathcal{I}_f(x_i)$ is cycle-balanced for each $i \in [8]$. I will generate $\mathcal{I}_f(x_1)$ in detail so we can see how the process goes. We will need to know $f(x_1^{i \rightarrow})$ for each $i \in [8]$. First, notice that $x_1^{i \rightarrow} = x_i$ for $i \in [2, 6]$. So we know the output of $f$ for each of these states from above. Now $x_1^{1 \rightarrow} = \lambda x_1$, so $f(x_1^{1 \rightarrow}) = \lambda f(x_1) = (-1, -1, 1, 1, 1, 1, 1)$. Next, $x_1^{7 \rightarrow} = \lambda^{-2}x_2$, so $f(x_1^{7 \rightarrow}) = \lambda^{-2}f(x_2) = (1, 1, 1, 1, 1, 1, -1)$. Finally, $x_1^{8 \rightarrow} = \lambda^{-1}x_1$, so $f(x_1^{8 \rightarrow}) = \lambda^{-1}f(x_1) = x_1$. This is all the data needed to generate the adjacency matrix of $\mathcal{I}_f(x_1)$. A picture of $\mathcal{I}_f(x_1)$ is given in figure 3.3. It contains no cycles, so it is cycle-balanced.

For the local interaction graph of each other orbit representative, we proceed in the same way. To find $f(x_i^{j \rightarrow})$, we just need to express $x_i^{j \rightarrow}$ as $\gamma x_k$. Then $f(x_i^{j \rightarrow}) = \gamma f(x_k)$. 

![Figure 3.3: $\mathcal{I}_f(x_1)$](image-url)
\[
\text{Adj}[\mathcal{I}_f(x_1)] = \begin{bmatrix}
\partial^1 \rightarrow f(x_1) & \partial^2 \rightarrow f(x_1) & \partial^3 \rightarrow f(x_1) & \partial^4 \rightarrow f(x_1) & \partial^5 \rightarrow f(x_1) & \partial^6 \rightarrow f(x_1) & \partial^7 \rightarrow f(x_1) & \partial^8 \rightarrow f(x_1) \\
-sgn([f(x_1^1) - f(x_1)]) & -sgn([f(x_1^2) - f(x_1)]) & -sgn([f(x_1^3) - f(x_1)]) & -sgn([f(x_1^4) - f(x_1)]) & -sgn([f(x_1^5) - f(x_1)]) & -sgn([f(x_1^6) - f(x_1)]) & -sgn([f(x_1^7) - f(x_1)]) & -sgn([f(x_1^8) - f(x_1)])
\end{bmatrix}
\]

\[
= \begin{bmatrix}
-sgn([-1, -1, 1, 1, 1, 1, 1, 1]) & -sgn([-1, 1, 1, 1, 1, 1, 1, 1]) & -sgn([1, 1, 1, 1, 1, 1, 1, 1]) & -sgn([1, 1, 1, 1, 1, 1, 1, 1]) & -sgn([1, 1, 1, 1, 1, 1, 1, 1]) & -sgn([1, 1, 1, 1, 1, 1, 1, 1]) & -sgn([1, 1, 1, 1, 1, 1, 1, 1]) & -sgn([1, 1, 1, 1, 1, 1, 1, 1]) \\
-sgn([0, -2, 0, 0, 0, 0, 0, 0]) & -sgn([0, 0, 0, 0, 0, 0, 0, 0]) & -sgn([2, 0, 0, 0, 0, 0, 0, 0]) & -sgn([2, 0, 0, 0, 0, 0, 0, 0]) & -sgn([2, 0, 0, 0, 0, 0, 0, 0]) & -sgn([2, 0, 0, 0, 0, 0, 0, 0]) & -sgn([2, 0, 0, 0, 0, 0, 0, 0]) & -sgn([2, 0, 0, 0, 0, 0, 0, 0])
\end{bmatrix}
\]

\[
= \begin{bmatrix}
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0
\end{bmatrix}
\]
\[
\text{Adj}[I_f(x_2)] = \begin{bmatrix}
- \text{sgn}[f(x_2^2) - f(x_2)] \\
\text{sgn}[f(x_2^2) - f(x_2)] \\
- \text{sgn}[f(x_2^2) - f(x_2)] \\
- \text{sgn}[f(x_2^2) - f(x_2)] \\
- \text{sgn}[f(x_2^2) - f(x_2)] \\
- \text{sgn}[f(x_2^2) - f(x_2)] \\
- \text{sgn}[f(x_2^2) - f(x_2)] \\
- \text{sgn}[f(x_2^2) - f(x_2)]
\end{bmatrix}
\begin{bmatrix}
- \text{sgn}[\lambda x_1 - f(x_2)] \\
\text{sgn}[x_1 - f(x_2)] \\
- \text{sgn}[\lambda x_6 - f(x_2)] \\
- \text{sgn}[\lambda x_6 - f(x_2)] \\
- \text{sgn}[\lambda x_6 - f(x_2)] \\
- \text{sgn}[\lambda x_6 - f(x_2)] \\
- \text{sgn}[\lambda x_6 - f(x_2)] \\
- \text{sgn}[\lambda x_6 - f(x_2)]
\end{bmatrix}
\begin{bmatrix}
- \text{sgn}[-1, -1, -1, 1, 1, 1, 1, 1] - f(x_2)] \\
\text{sgn}[-1, -1, -1, 1, 1, 1, 1, 1] - f(x_2)] \\
- \text{sgn}[1, -1, 1, -1, 1, -1, 1, 1] - f(x_2)] \\
- \text{sgn}[1, -1, 1, -1, 1, -1, 1, 1] - f(x_2)] \\
- \text{sgn}[1, -1, 1, -1, 1, -1, 1, 1] - f(x_2)] \\
- \text{sgn}[1, -1, 1, -1, 1, -1, 1, 1] - f(x_2)] \\
- \text{sgn}[1, -1, 1, -1, 1, -1, 1, 1] - f(x_2)] \\
- \text{sgn}[1, -1, 1, -1, 1, -1, 1, 1] - f(x_2)]
\end{bmatrix}
\begin{bmatrix}
0 + + 0 0 0 0 0 0 \\
0 0 0 0 0 0 0 0 \\
0 + 0 0 0 0 0 0 \\
- + 0 + 0 + + 0 \\
- + 0 0 + 0 + 0 \\
- + 0 + 0 + 0 0 \\
- + 0 + + 0 + 0 \\
- 0 0 0 0 0 0 +
\end{bmatrix}
\]
\[ \text{Adj}[\mathcal{A}_f(x_3)] = \begin{bmatrix}
- \text{sgn}[f(x_1) - f(x_3)] \\
- \text{sgn}[f(x_2) - f(x_3)] \\
\text{sgn}[f(x_3) - f(x_3)] \\
- \text{sgn}[f(x_4) - f(x_3)] \\
- \text{sgn}[f(x_5) - f(x_3)] \\
- \text{sgn}[f(x_6) - f(x_3)] \\
- \text{sgn}[f(x_7) - f(x_3)] \\
- \text{sgn}[f(x_8) - f(x_3)] \\
\end{bmatrix} = \begin{bmatrix}
- \text{sgn}[f(x_2) - f(x_3)] \\
- \text{sgn}[f(x_6) - f(x_3)] \\
\text{sgn}[f(x_1) - f(x_3)] \\
- \text{sgn}[f(x_7) - f(x_3)] \\
- \text{sgn}[f(x_8) - f(x_3)] \\
- \text{sgn}[f(x_4) - f(x_3)] \\
- \text{sgn}[f(\zeta x_6) - f(x_3)] \\
- \text{sgn}[f(\zeta x_4) - f(x_3)] \\
\end{bmatrix} \]
\[
\text{Adj}[Z_f(x_4)] = \begin{bmatrix}
- \text{sgn}[f(x_1^4) - f(x_4)] \\
- \text{sgn}[f(x_2^4) - f(x_4)] \\
- \text{sgn}[f(x_3^4) - f(x_4)] \\
- \text{sgn}[f(x_4^4) - f(x_4)] \\
- \text{sgn}[f(x_5^4) - f(x_4)] \\
- \text{sgn}[f(x_7^4) - f(x_4)] \\
- \text{sgn}[f(x_8^4) - f(x_4)] \\
- \text{sgn}[f(x_9^4) - f(x_4)]
\end{bmatrix} = \begin{bmatrix}
- \text{sgn}[f(x_1^4) - f(x_4)] \\
- \text{sgn}[f(x_2^4) - f(x_4)] \\
- \text{sgn}[f(x_3^4) - f(x_4)] \\
- \text{sgn}[f(x_4^4) - f(x_4)] \\
- \text{sgn}[f(x_5^4) - f(x_4)] \\
- \text{sgn}[f(x_7^4) - f(x_4)] \\
- \text{sgn}[f(x_8^4) - f(x_4)] \\
- \text{sgn}[f(x_9^4) - f(x_4)]
\end{bmatrix}
\]
\[
\text{Adj}[I_f(x_5)] = \begin{bmatrix}
-\text{sgn}[f(x_5^{1\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{2\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{3\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{4\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{5\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{6\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{7\to}) - f(x_5)] \\
-\text{sgn}[f(x_5^{8\to}) - f(x_5)]
\end{bmatrix} = \begin{bmatrix}
-\text{sgn}[f(\lambda x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta^2 x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta^3 x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta^4 x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta^5 x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta^6 x_4) - f(x_5)] \\
-\text{sgn}[f(\zeta^7 x_4) - f(x_5)]
\end{bmatrix}
\]

\[
\begin{bmatrix}
-\text{sgn}[(-1, 1, 1, 1, 1, 1, 1, 1) - f(x_5)] \\
-\text{sgn}[(1, -1, 1, 1, -1, 1, -1, 1) - f(x_5)] \\
-\text{sgn}[(-1, 1, -1, 1, -1, 1, -1, 1) - f(x_5)] \\
-\text{sgn}[(1, 1, -1, 1, 1, 1, 1, 1) - f(x_5)] \\
\text{sgn}[-1, 1, 1, 1, 1, 1, 1, 1) - f(x_5)] \\
-\text{sgn}[(1, 1, 1, -1, 1, -1, 1, 1, 1) - f(x_5)] \\
-\text{sgn}[(-1, -1, 1, 1, 1, 1, 1, 1) - f(x_5)] \\
-\text{sgn}[(1, -1, -1, 1, -1, 1, -1, 1) - f(x_5)] \\
-\text{sgn}[(1, 1, 1, 1, -1, 1, -1, 1) - f(x_5)] \\
-\text{sgn}[(1, 1, 1, 1, 1, -1, 1, -1) - f(x_5)]
\end{bmatrix}
= \begin{bmatrix}
+ 0 0 0 0 0 0 0 0 \\
0 + 0 0 + 0 + 0 \\
+ 0 + 0 + 0 + 0 \\
0 0 0 + 0 0 0 0 \\
- 0 0 0 0 0 0 0 0 \\
0 0 0 0 + + + 0 \\
0 + + 0 + 0 + 0 \\
0 0 0 0 0 0 + +
\end{bmatrix}
\]

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\[
\text{Adj}[I_f(x_6)] = \\
\begin{bmatrix}
- \text{sgn}[f(x_1^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_2^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_3^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_4^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_5^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_6^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_7^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_8^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_9^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_{10}^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_{11}^{\rightarrow}) - f(x_6)] \\
- \text{sgn}[f(x_{12}^{\rightarrow}) - f(x_6)] \\
\end{bmatrix}
= \\
\begin{bmatrix}
- \text{sgn}[f(x_1) - f(x_6)] \\
- \text{sgn}[f(x_2) - f(x_6)] \\
- \text{sgn}[f(x_3) - f(x_6)] \\
- \text{sgn}[f(x_4) - f(x_6)] \\
- \text{sgn}[f(x_5) - f(x_6)] \\
- \text{sgn}[f(x_6) - f(x_6)] \\
- \text{sgn}[f(x_7) - f(x_6)] \\
- \text{sgn}[f(x_8) - f(x_6)] \\
- \text{sgn}[f(x_9) - f(x_6)] \\
- \text{sgn}[f(x_{10}) - f(x_6)] \\
- \text{sgn}[f(x_{11}) - f(x_6)] \\
- \text{sgn}[f(x_{12}) - f(x_6)] \\
\end{bmatrix}
= \\
\begin{bmatrix}
- \text{sgn}((-1, 1, 1, 1, 1, 1, 1, 1) - f(x_6)) \\
- \text{sgn}((1, -1, 1, -1, 1, -1, 1, 1) - f(x_6)) \\
- \text{sgn}((-1, 1, -1, 1, 1, -1, 1, 1) - f(x_6)) \\
- \text{sgn}((1, -1, 1, -1, 1, -1, 1, 1) - f(x_6)) \\
- \text{sgn}((1, 1, 1, 1, 1, 1, 1, 1) - f(x_6)) \\
- \text{sgn}((-1, 1, 1, 1, 1, 1, 1, 1) - f(x_6)) \\
- \text{sgn}((1, 1, 1, 1, 1, 1, 1, 1, 1) - f(x_6)) \\
- \text{sgn}((1, 1, 1, 1, 1, 1, 1, 1, 1) - f(x_6)) \\
\end{bmatrix}
= \\
\begin{bmatrix}
+ 0 0 0 0 0 0 0 0 \\
0 + 0 + 0 + 0 + 0 \\
+ 0 + 0 0 + 0 + 0 \\
0 + 0 + 0 + 0 + 0 \\
0 0 0 0++ + + + \\
- 0 0 0 0 0 0 0 + \\
0 0 0 0 0 + + 0 \\
0 0 0 0 0 + 0 + 0 \\
\end{bmatrix}
\]
\[
\text{Adj}[I_f(x_7)] =
\begin{bmatrix}
- \text{sgn}[f(x_1)] - f(x_7) \\
- \text{sgn}[f(x_2)] - f(x_7) \\
- \text{sgn}[f(x_3)] - f(x_7) \\
- \text{sgn}[f(x_4)] - f(x_7) \\
- \text{sgn}[f(x_5)] - f(x_7) \\
- \text{sgn}[f(x_6)] - f(x_7) \\
- \text{sgn}[f(x_7)] - f(x_7) \\
- \text{sgn}[f(x_8)] - f(x_7) \\
- \text{sgn}[f(x_9)] - f(x_7)
\end{bmatrix}
= 
\begin{bmatrix}
- \text{sgn}[f(x_1) - f(x_7)] \\
- \text{sgn}[f(x_2) - f(x_7)] \\
- \text{sgn}[f(x_3) - f(x_7)] \\
- \text{sgn}[f(x_4) - f(x_7)] \\
- \text{sgn}[f(x_5) - f(x_7)] \\
- \text{sgn}[f(x_6) - f(x_7)] \\
- \text{sgn}[f(x_7) - f(x_7)] \\
- \text{sgn}[f(x_8) - f(x_7)] \\
- \text{sgn}[f(x_9) - f(x_7)]
\end{bmatrix}
= 
\begin{bmatrix}
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)] \\
- \text{sgn}[(1, 1, 1, 1, 1, 1, 1, 1, 1)] - f(x_7)]
\end{bmatrix}
= 
\begin{bmatrix}
+ + + + - 0 0 0 0 \\
+ + + + 0 0 0 0 0 \\
0 0 0 + + 0 0 0 0 \\
0 0 0 + + 0 0 0 0 \\
0 + + + + 0 0 0 0 \\
0 + + + + 0 0 0 0 \\
0 0 + + - 0 0 + 0 \\
0 0 + + - 0 0 0 + \\
0 0 + + - 0 0 0 + 
\end{bmatrix}
\]
Each of these interaction graphs is cycle-balanced. Since every local interaction graph is switching isomorphic to one of these, every local interaction graph is cycle-balanced. So \( f \) is indeed an activation function whose local interaction graphs contain no negative cycles that has no steady-states.
Figure 3.4: Local interaction graphs of the first four orbit representatives.
Figure 3.5: Local interaction graphs of last four orbit representatives.
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